

## Fast Matrix Algorithms for Data Analytics: Problem Set 2

1. Let  $\mathbf{A} \in \mathbb{R}^{m \times n}$ ,  $J$  be an  $n \times 1$  permutation vector, and  $J_s = J(1 : k)$  for some  $k \in \mathbb{N}$ . Now suppose

$$\begin{array}{ccccc} \mathbf{A} & = & \mathbf{E} & \mathbf{F}, & \\ m \times n & & m \times k & k \times n & \end{array} \quad (1)$$

and suppose that for some matrix  $\mathbf{T}$  of size  $k \times (n - k)$ , it holds that

$$\mathbf{F}(:, J) = \mathbf{F}(:, J_s) \begin{bmatrix} \mathbf{I}_k & \mathbf{T} \end{bmatrix},$$

where  $\mathbf{I}_k$  is the  $k \times k$  identity matrix. Prove that

$$\mathbf{A}(:, J) = \mathbf{A}(:, J_s) \begin{bmatrix} \mathbf{I}_k & \mathbf{T} \end{bmatrix}.$$

2. Suppose that  $\mathbf{A}$  is an  $m \times n$  matrix of precise rank  $k$ . Moreover, suppose that you have available a double-sided ID of  $\mathbf{A}$  of the form

$$\begin{array}{ccccc} \mathbf{A} & = & \mathbf{X} & \mathbf{A}_s & \mathbf{Z}, \\ m \times n & & m \times k & k \times k & k \times n \end{array} \quad (2)$$

where  $\mathbf{A}_s = \mathbf{A}(I_s, J_s)$  for some index vectors  $I_s$  and  $J_s$ . Recall from the definition of the double-sided ID that  $\mathbf{X}$  and  $\mathbf{Z}$  must contain the  $k \times k$  identity matrix  $\mathbf{I}_k$  as a submatrix, so that  $\mathbf{X}(I_s, :) = \mathbf{I}_k$  and  $\mathbf{Z}(:, J_s) = \mathbf{I}_k$ .

- (a) Prove that  $\mathbf{A}_s$  must be of full rank.  
 (b) Set  $\mathbf{U} = \mathbf{A}_s^{-1}$ , so that (2) can be written

$$\begin{array}{ccccc} \mathbf{A} & = & \mathbf{X}\mathbf{A}_s & \mathbf{U} & \mathbf{A}_s\mathbf{Z}. \\ m \times n & & m \times k & k \times k & k \times n \end{array} \quad (3)$$

Prove that  $\mathbf{X}\mathbf{A}_s = \mathbf{A}(:, J_s)$  and  $\mathbf{A}_s\mathbf{Z} = \mathbf{A}(I_s, :)$  so that (3) is a CUR factorization.

3. Suppose that  $\mathbf{A}$  is an  $m \times n$  matrix of approximate rank  $k$ , and that we have identified two index sets  $I_s$  and  $J_s$  such that the matrices

$$\mathbf{C} = \mathbf{A}(:, J_s), \quad \mathbf{R} = \mathbf{A}(I_s, :)$$

hold  $k$  columns/rows that span the column/row space of  $\mathbf{A}$ . Then

$$\mathbf{A} \approx \mathbf{C}\mathbf{C}^\dagger\mathbf{A}\mathbf{R}^\dagger\mathbf{R},$$

and the optimal choice for the “U” factor in the CUR decomposition is

$$\mathbf{U} = \mathbf{C}^\dagger\mathbf{A}\mathbf{R}^\dagger.$$

Set  $\mathbf{X} = \mathbf{C}\mathbf{C}^\dagger$ .

- (a) Suppose that  $\mathbf{C}$  has the SVD  $\mathbf{C} = \mathbf{W}\mathbf{D}\mathbf{V}^*$ . Prove that  $\mathbf{X} = \mathbf{W}\mathbf{W}^*$ .
- (b) Suppose that  $\mathbf{C}$  has the QR factorization  $\mathbf{C}\mathbf{P} = \mathbf{Q}\mathbf{S}$ . Prove that  $\mathbf{X} = \mathbf{Q}\mathbf{Q}^*$ .
- (c) Prove that  $\mathbf{X}$  is the orthogonal projection onto  $\text{Col}(\mathbf{C})$ .
- (d) Suppose that  $\mathbf{A}$  has precisely rank  $k$  and that  $\mathbf{C}$  and  $\mathbf{R}$  are both of rank  $k$ . Prove that then  $\mathbf{C}^\dagger\mathbf{A}\mathbf{R}^\dagger = (\mathbf{A}(I_s, J_s))^{-1}$ .

4. Let  $\mathbf{A} \in \mathbb{R}^{m \times n}$  have rank exactly  $k$ . In this problem, we will prove that  $\mathbf{A}$  admits a factorization  $\mathbf{A} = \mathbf{A}(:, J_s)\mathbf{Z}$ , where  $\mathbf{A}(:, J_s) \in \mathbb{R}^{m \times k}$  and  $\mathbf{Z} \in \mathbb{R}^{k \times n}$  such that  $\mathbf{Z}(:, J_s) = \mathbf{I}_k$  and  $\max_{i,j} |\mathbf{Z}(i, j)| \leq 1$ .

(a) **case 1:**  $m = k$ .

- i. Pick a permutation vector  $J_s$  such that  $|\det(\mathbf{A}(:, J_s))|$  is maximized, and let  $J_r$  denote the remaining indices so that  $[J_s \ J_r]$  is some permutation of the vector  $[1 \ 2 \ \cdots \ n]$ . Then we have that

$$\mathbf{A}(:, [J_s \ J_r]) = [\mathbf{A}(:, J_s) \ \mathbf{A}(:, J_r)]$$

can be written as  $\mathbf{A}\mathbf{P}$  for some permutation matrix  $\mathbf{P}$ . Find an interpolative decomposition  $\mathbf{A} = \mathbf{C}\mathbf{Z}$  of  $\mathbf{A}$ , where the columns of  $\mathbf{C}$  are some of the columns of  $\mathbf{A}$ .  $\mathbf{C}$  and  $\mathbf{Z}$  should be in terms of  $\mathbf{A}(:, J_s)$ ,  $\mathbf{A}(:, J_r)$ ,  $\mathbf{P}$ , and the identity matrix  $\mathbf{I}$ .

- ii. Consider the matrix  $\mathbf{T} = \mathbf{A}(:, J_s)^{-1}\mathbf{A}(:, J_r)$ . If we can show that

$$\max_{i,j} |\mathbf{T}(i, j)| \leq 1, \tag{4}$$

then we will be done with the case  $m = k$  (why?). Find a way to show (4) by applying Cramer's Rule to our definition of  $\mathbf{T}$ .

**Cramer's Rule:** Consider the linear system  $\mathbf{A}\mathbf{x} = \mathbf{b}$ . The  $i$ -th entry of the solution  $\mathbf{x}$  is given by

$$x_i = \frac{\det(\mathbf{A}_i)}{\det(\mathbf{A})},$$

where  $\mathbf{A}_i$  is matrix formed by replacing the  $i$ -th column of  $\mathbf{A}$  with  $\mathbf{b}$ .

(b) **case 2:**  $m \geq k$ .

Then  $\mathbf{A}$  admits a factorization  $\mathbf{A} = \mathbf{E}\mathbf{F}$ , where  $\mathbf{E}$  is  $m \times k$  and  $\mathbf{F}$  is  $k \times n$ . Apply case 1 to  $\mathbf{F}$  to show the result for this case (something we proved in a previous problem may help, too...).