

# 2013 APPM 6640 Multigrid Methods

Lectures: Mondays & Wednesdays 10-10<sup>50</sup>am ECCR 1B51  
Computing Lab: Mondays 12-12<sup>50</sup> ECCR 143  
Office: Mondays 11-11<sup>50</sup> & Wednesdays 9-9<sup>50</sup> ECCR 257

**Steve McCormick**  
[stevem@colorado.edu](mailto:stevem@colorado.edu)

**Toby Jones**  
[Tobias.Jones@colorado.edu](mailto:Tobias.Jones@colorado.edu)

**Chris Leibs**  
[chris.leibs@gmail.com](mailto:chris.leibs@gmail.com)

303-492-0662 (email is better!)  
<http://amath.colorado.edu/faculty/stevem>

CU-Boulder

1 of 396

## APPM 6640 Syllabus

- Coursework
  - No exams
  - Homework exercises & computing assignments (Lab)
  - Team project: identify & pursue a target application
- Philosophy (the course is not just about multigrid!!!)
  - Guiding rule: **understanding trumps knowledge**
  - Look for **fundamentals** & **basic principles** to guide you
  - Solutions are often **straightforward consequences**
  - Think about the **scientific method** in general
  - Understand by **concrete examples** & **experience**
  - Know the **whole picture**
  - Ask **questions** & **interject comments**
  - Make sure I **explain** what you need
- Text
  - A Multigrid Tutorial, 2<sup>nd</sup> edition, 2<sup>nd</sup> printing
  - Supplemental material as needed

CU-Boulder

2 of 396

## Sources

**MGNet** Newsletter & software repository  
<http://www.mgnet.org>

**MathSciNet** Many papers electronically available  
<http://www.ams.org/mathscinet>

**Copper Mountain Conference** March 17-22, 2013  
<http://grandmaster.colorado.edu/~copper>  
We plan to support your attendance there.

**Course web site:**

<http://grandmaster.colorado.edu/appm6640/>  
These slides are there.

CU-Boulder

3 of 396

## Homework exercises

Due one week after the relevant chapter is covered in class.

- |                       |                            |
|-----------------------|----------------------------|
| 1: 1, 3, 4            | 6: 1, 2, 6                 |
| 2: 1-3, 8, 10, 13, 16 | 7: 1, 2, 7, 10, 12, 15, 21 |
| 3: 2-6                | 8: 2, 6, 8                 |
| 4: 6, 7, 11           | 9: 1, 3                    |
| 5: 1-3, 12, 13        | 10: 1, 3, 5, 8, 9          |

## Computing Assignments

You will learn about this in the lab &/or from the course website.

CU-Boulder

4 of 396

## Team project

main objective: understand applied math research

- Group into teams of possible common interest.
- Meet in lab as time permits later in the semester.
- Identify an application area of interest.
- Identify a problem in that area & learn about it.
- Learn current methods & their limitations.
- Get experience with these methods on typical cases.
- Try cases where these methods (begin to) fail.
- Brainstorm a better (multilevel?) method.
- Implement & test your idea.
- Inform the rest of the class along the way.

CU-Boulder

5 of 396

## Prolog: Multigrid in Action

The following soldier slides were created by

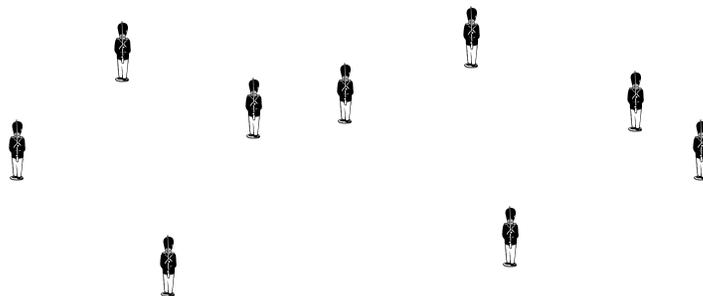
Irada Yavneh  
Department of Computer Science  
Technion - Israel Institute of Technology  
[irad@cs.technion.ac.il](mailto:irad@cs.technion.ac.il)

Basic Concepts: Local vs. Global processing.

- Imagine a large number of soldiers who need to be arranged in a straight line and at equal distances from each other.
- The two soldiers at the ends of the line are fixed. Suppose we number the soldiers  $0$  to  $N$ , and that the length of the entire line is  $L$ .

CU-Boulder

6 of 396



Initial Position



Final Position

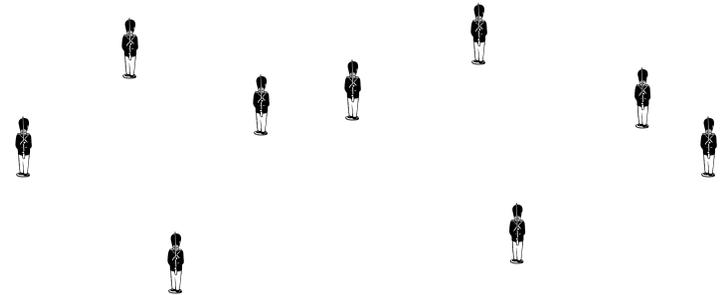
CU-Boulder

7 of 396

CU-Boulder

8 of 396

**Global processing.** Let soldier number  $j$  stand on the line connecting soldier  $0$  to soldier  $N$  at a distance  $jL/N$  from soldier number  $0$ .



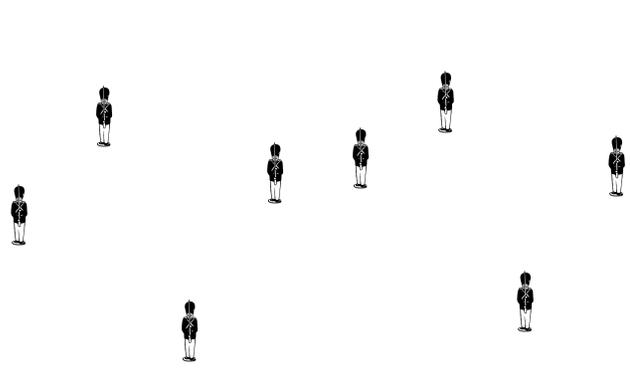
This method solves the problem **directly**, but requires substantial **sophistication**: recognition of the extreme soldiers and some pretty fancy arithmetic.

**Local processing (iterative method).** Suppose that the initial position of inner soldier  $j$  is  $x_j$ . Then if every  $j$  moves **all at once** to the point midway between the initial locations of neighboring soldiers,  $j-1$  &  $j+1$ , we get

$$x_j \leftarrow (x_{j-1} + x_{j+1})/2.$$

(Assume for simplicity that the soldiers have guides to make sure they're evenly spaced, so they only have to get in a straight line. Thus,  $x_j$  is their signed distance from that line.)

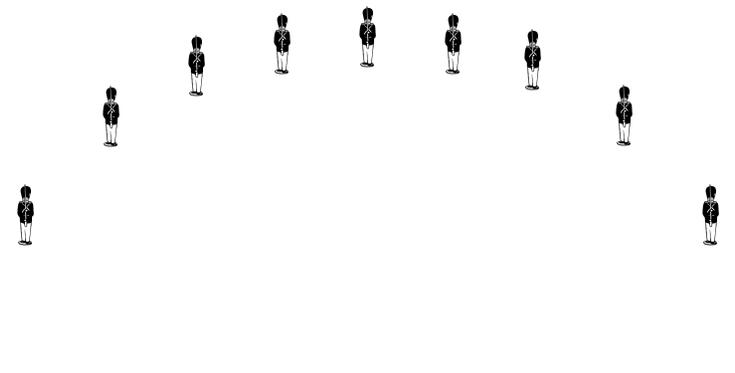
- This is an **iterative** process.
- Each step brings us closer to the solution (**convergence**).
- The arithmetic is **trivial**.
- The process is **local**.



A step in the right direction



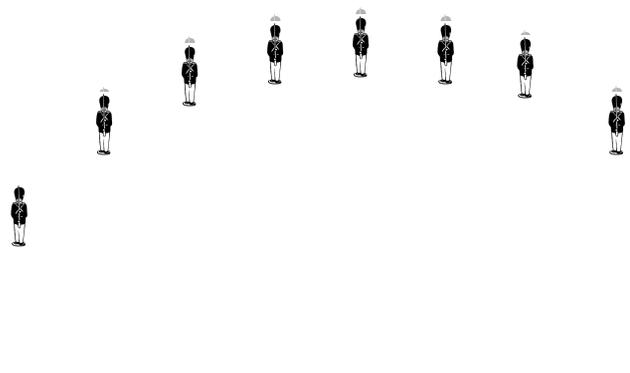
13 of 396



14 of 396

CU-Boulder

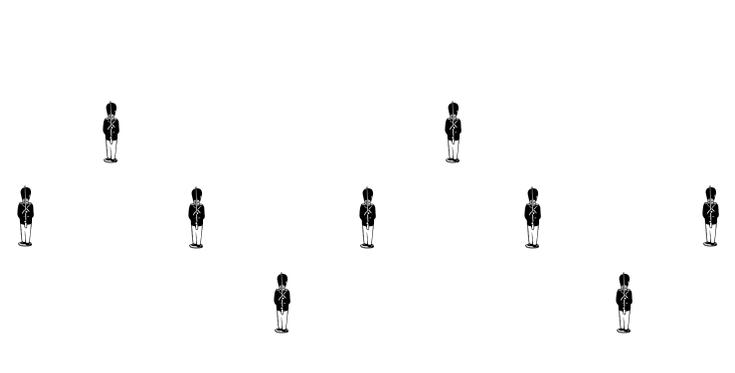
CU-Boulder



Slow convergence



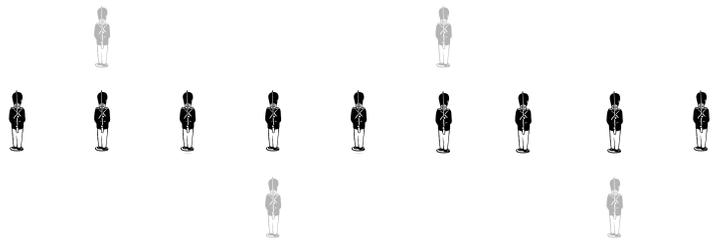
15 of 396



16 of 396

CU-Boulder

CU-Boulder

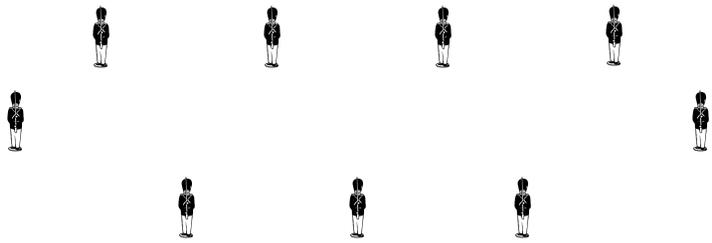


Fast convergence



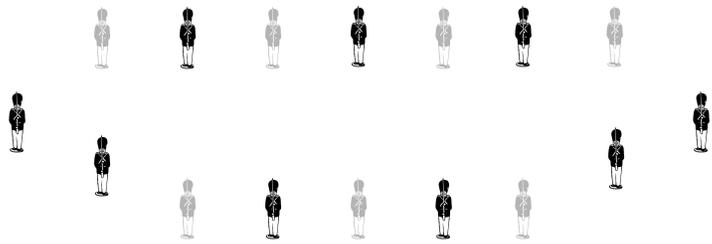
17 of 396

CU-Boulder



18 of 396

CU-Boulder

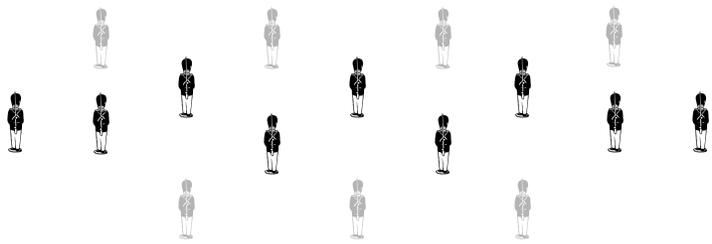


Slow convergence



19 of 396

CU-Boulder



Local solution: damping



20 of 396

CU-Boulder



Local solution: damping

CU-Boulder



21 of 396



Local solution: damping

CU-Boulder



22 of 396



Local solution: damping

CU-Boulder



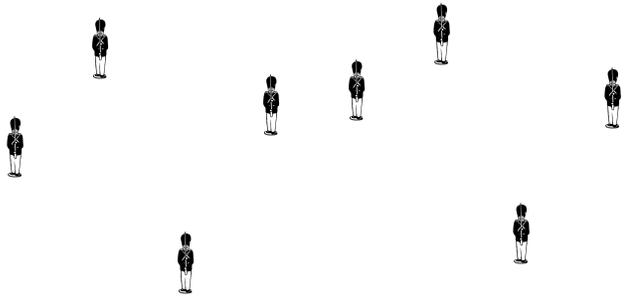
23 of 396

The multiscale idea: Employ the local processing with simple arithmetic. But do this on all the different scales.

CU-Boulder

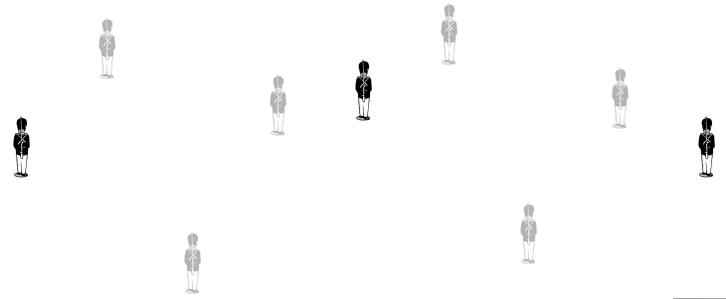
24 of 396

CU-Boulder



25 of 396

CU-Boulder

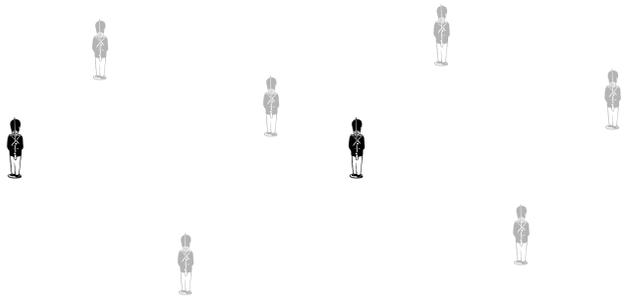


Large scale



26 of 396

CU-Boulder

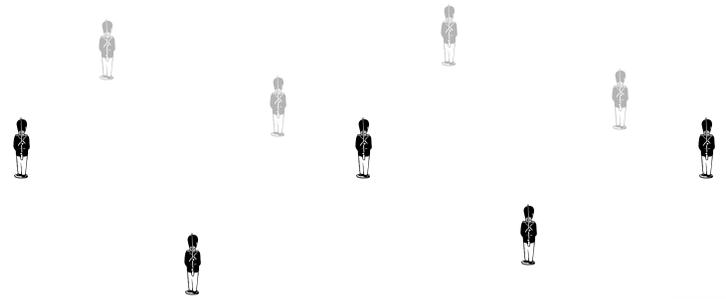


Large scale



27 of 396

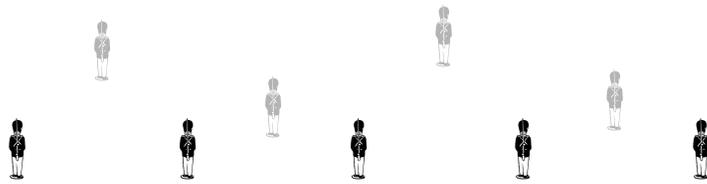
CU-Boulder



Intermediate scale



28 of 396

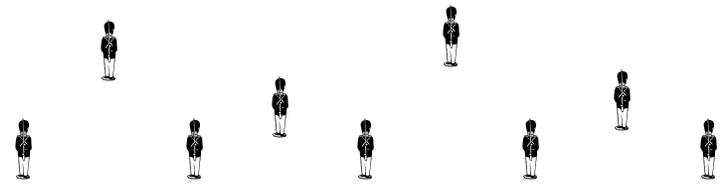


Intermediate scale



29 of 396

CU-Boulder



Small scale



30 of 396

CU-Boulder

### How much work do we save?

Jacobi needs about  $N^2$  iterations &  $N^2 \times N = N^3$  ops to improve accuracy by an order of magnitude.

Brandt solves the problem in only about  $N$  operations.

Example: for  $N = 1000$ , MG needs about

1,000 operations

instead of about

1,000,000,000 operations !!!



31 of 396

CU-Boulder

CU-Boulder

32 of 396

## How important is computational efficiency?

Suppose we have 3 different algorithms for a given problem, with different computational complexities for input size  $N$ :

Algorithm 1:  $10^6 N$  ops

Algorithm 2:  $10^3 N^2$  ops

Algorithm 3:  $N^3$  ops

Suppose  $N$  is such that algorithm 1 requires one second.

How long do the others require?

Computer Speed (ops/sec)	$N$	Algorithm 1 $O(N)$	Algorithm 2 $O(N^2)$	Algorithm 3 $O(N^3)$
1M ( $\sim 10^6$ ) (1980's)	1	1 sec	0.001 sec	0.000001 sec
1G ( $\sim 10^9$ ) (1990's)	1K	1 sec	1 sec	1 sec
1T ( $\sim 10^{12}$ ) (2000's)	1M	1 sec	17 min	12 days
1P ( $\sim 10^{15}$ ) (2010's)	1G	1 sec	12 days	31,710 years

Stronger computers  $\Rightarrow$  more gain!

## The catch

In less trivial problems, we can't construct appropriate equations on the large scales without first propagating information from the small scales.

Skill in developing efficient multigrid is needed for:

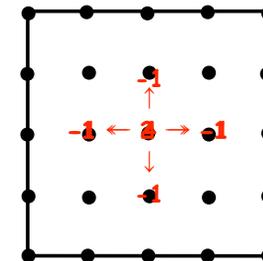
1. Choosing a good local iteration.
2. Choosing appropriate coarse-scale variables.
3. Choosing inter-scale transfer operators.
4. Constructing coarse-scale approximations to fine scales.

## What about two dimensions?

Put points midway between horizontal (or vertical) neighbors.

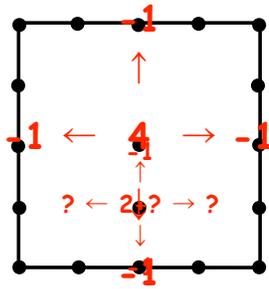
This is just imposing  $x_i = (x_{i-1} + x_{i+1})/2$  on each row  $j$  or

$$2x_i - x_{i-1} - x_{i+1} = 0.$$

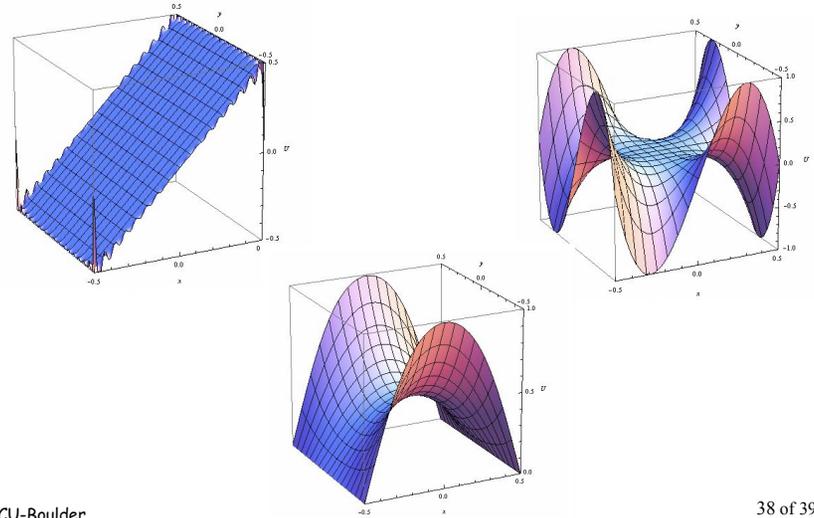


The hitch is that this is not a common physical problem. More common is to ask that some physical quantity at each point be an average of its **FOUR** neighbors (Poisson).

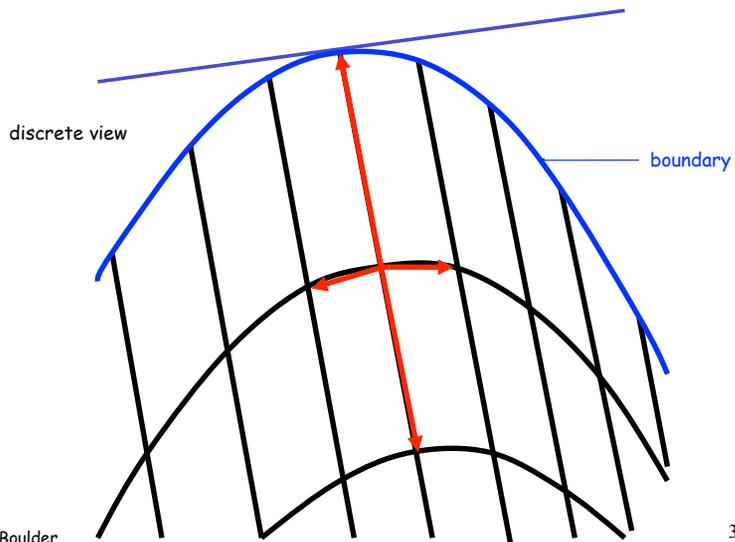
## Poisson



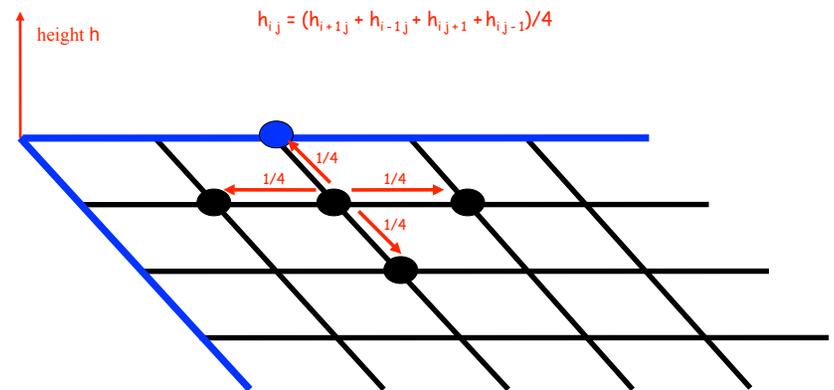
## Poisson: minimal surface given boundary values, minimize surface area

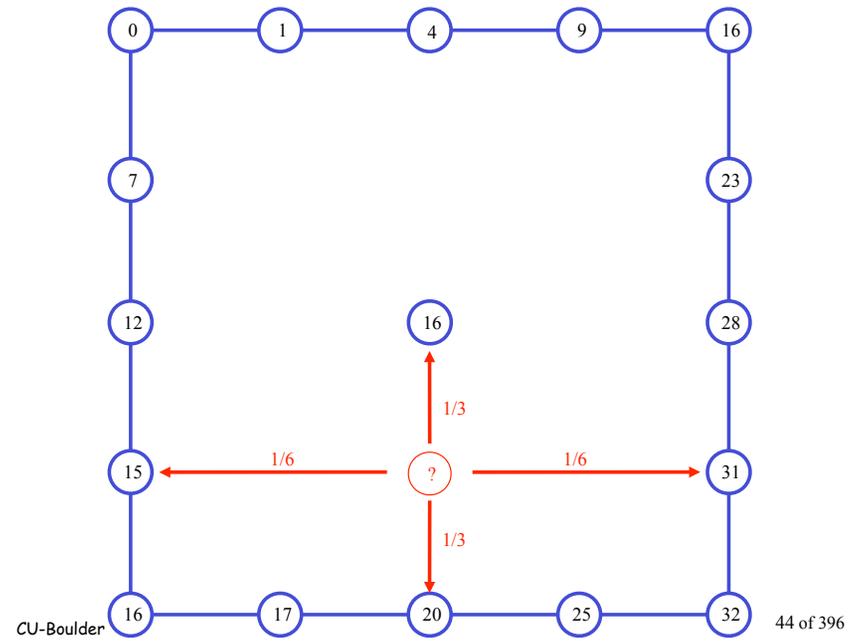
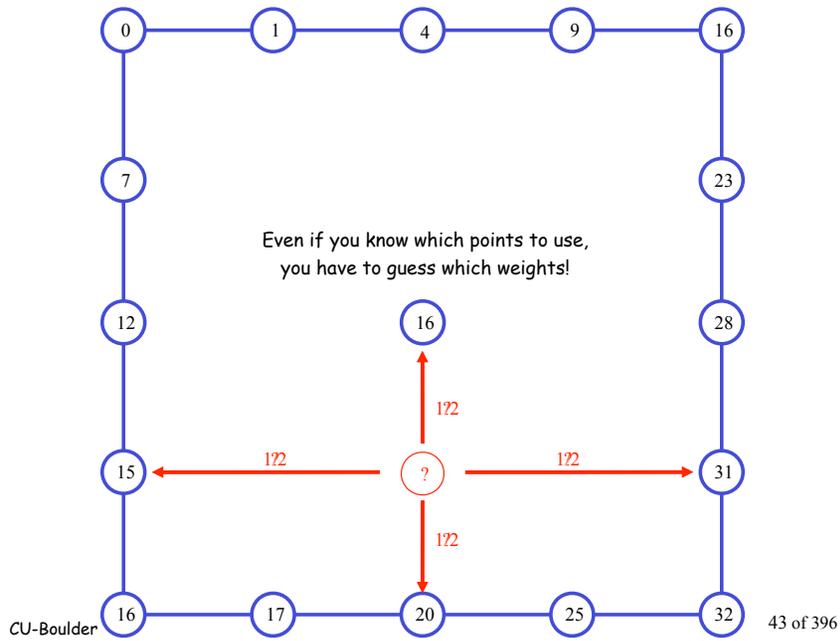
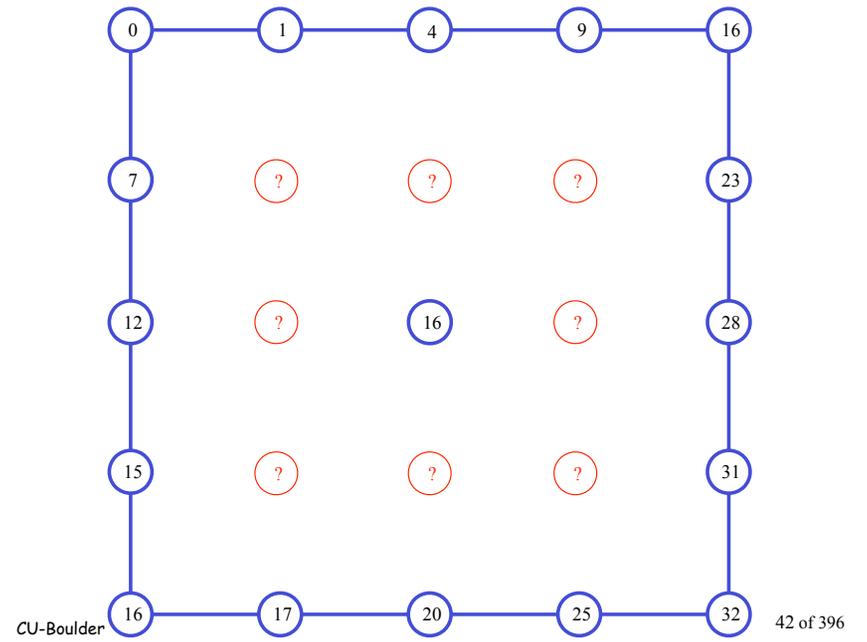
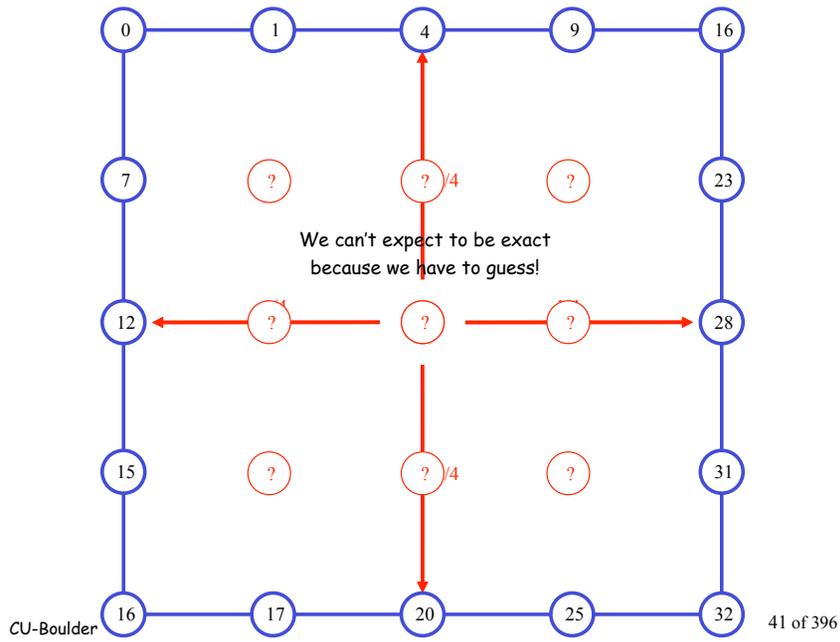


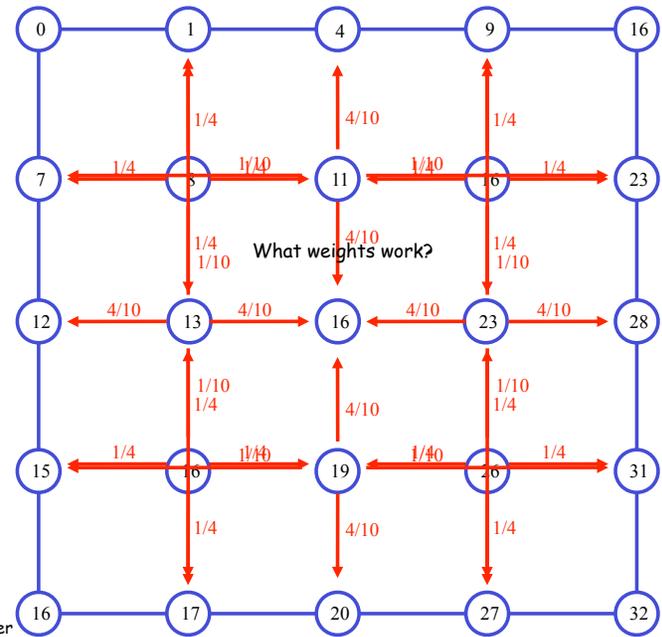
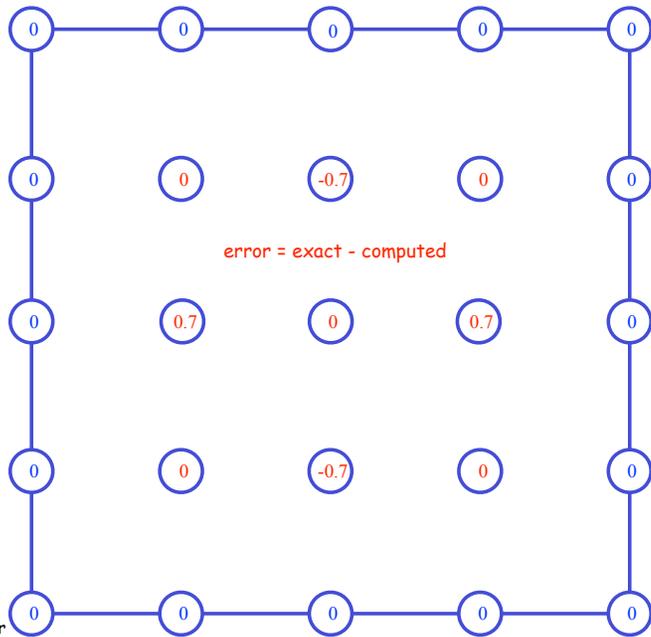
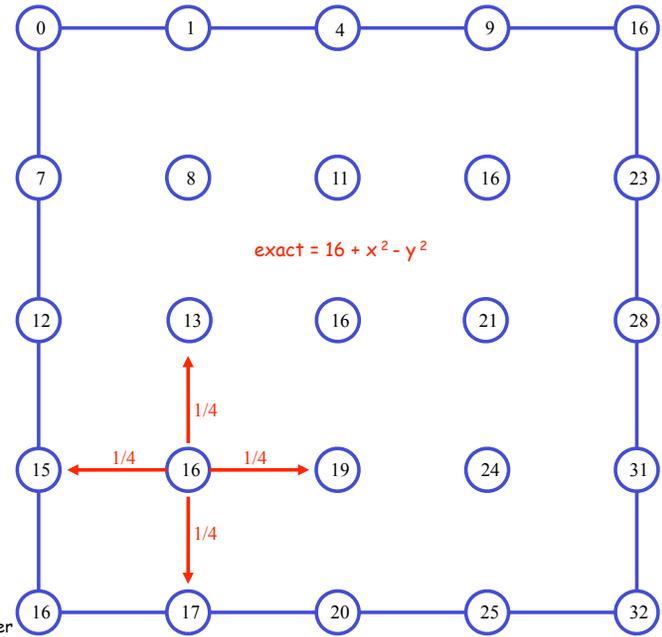
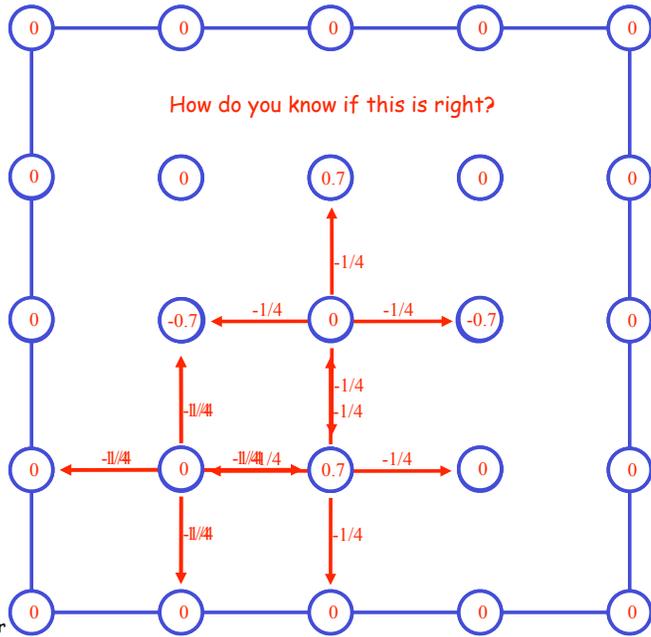
## Physical principle



## Local relationship







# A Multigrid Tutorial

2<sup>nd</sup> Edition, 2<sup>nd</sup> Printing

By

William L. Briggs

CU-Denver

Van Emden Henson

LLNL

Steve McCormick

CU-Boulder

## Outline

by chapter

1. Model Problems
2. Basic Iterative Methods
  - Convergence tests
  - Analysis
3. Elements of Multigrid
  - Relaxation
  - Coarsening
4. Implementation
  - Complexity
  - Diagnostics
5. Some Theory
  - Spectral vs. algebraic
6. Nonlinear Problems
  - Full approximation scheme
7. Selected Applications
  - Neumann boundaries
  - Anisotropic problems
  - Variable meshes
  - Variable coefficients
8. Algebraic Multigrid (AMG)
  - Matrix coarsening
9. Multilevel Adaptive Methods
  - FAC
10. Finite Elements
  - Variational methodology

## Suggested reading

CHECK THE MG LIBRARY & MGNET REPOSITORY

- A. Brandt, "Multi-level Adaptive Solutions to Boundary Value Problems," *Math Comp.*, 31, 1977, pp 333-390.
- A. Brandt, "Multigrid techniques: 1984 guide with applications to computational fluid dynamics," *GMD*, 1984.
- W. Hackbusch & U. Trottenberg, "Multigrid Methods", Springer-Verlag, 1982.
- S. McCormick, ed., "Multigrid Methods," *SIAM Frontiers in Applied Math.* III, 1987.
- U. Trottenberg, C. Oosterlee, & A. Schüller, "Multigrid," *Academic Press*, 2000.
- P. Wesseling, "An Introduction to Multigrid Methods," *Wylie*, 1992.

## Multilevel methods have been developed for...

- PDEs, CFD, porous media, elasticity, electromagnetics.
- Purely algebraic problems, with no physical grid; for example, network & geodetic survey problems.
- Image reconstruction & tomography.
- Optimization (e.g., the traveling salesman & long transportation problems).
- Statistical mechanics, Ising spin models.
- Quantum chromodynamics.
- Quadrature & generalized FFTs.
- Integral equations.
- 
- 
-

# Everyone uses multilevel methods

- Multigrid, multilevel, multiscale, multiphysics, ...  
Use local "governing rules" at the finest resolution to resolve details of the state of the system, but use coarser resolution to resolve larger scales. Continual feedback is essential because improving one scale impacts other scales.
- Common uses  
Sight, art, team sports, politics, society, thinking, scientific research, ...

# 1. Model problems

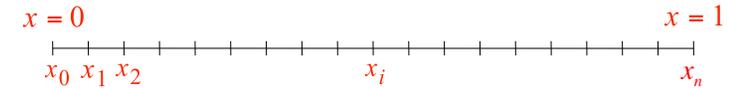
- 1-D boundary value problem:

$$-u''(x) + \sigma u(x) = f(x) \quad 0 < x < 1, \quad \sigma \geq 0$$

$$u(0) = u(1) = 0$$

- Grid:

$$h = \frac{1}{n}, \quad x_i = ih, \quad i = 0, 1, \dots, n$$



- Let  $v_i \approx u(x_i)$  &  $f_i \approx f(x_i)$  for  $i = 0, 1, \dots, n$ .

This discretizes the variables, but what about the equations?

## Approximate $u''(x)$ via Taylor series

$O(h^2)$  means a quantity bounded in norm by  $Ch^2$  for some constant  $C$ .

- Approximate 2<sup>nd</sup> derivative using Taylor series:

$$u(x_{i+1}) = u(x_i) + h u'(x_i) + \frac{h^2}{2!} u''(x_i) + \frac{h^3}{3!} u'''(x_i) + O(h^4)$$

$$+$$

$$u(x_{i-1}) = u(x_i) - h u'(x_i) + \frac{h^2}{2!} u''(x_i) - \frac{h^3}{3!} u'''(x_i) + O(h^4)$$

- Summing & solving:

$$u''(x_i) = \frac{u(x_{i+1}) - 2u(x_i) + u(x_{i-1}))}{h^2} + O(h^2)$$

## Approximate equation via finite differences

Approximate the BVP

$$-u''(x) + \sigma u(x) = f(x) \quad 0 < x < 1, \quad \sigma \geq 0$$

$$u(0) = u(1) = 0$$

by a finite difference scheme:

$$\frac{-v_{i-1} + 2v_i - v_{i+1}}{h^2} = -\frac{v_{i+1} - v_i}{h} - \frac{v_i - v_{i-1}}{h}$$

$$v_0 = v_n = 0$$



## Discretizing the 2-D problem

- Let  $v_{ij} \approx u(x_i, y_j)$  &  $f_{ij} \approx f(x_i, y_j)$ . Again, using 2<sup>nd</sup>-order finite differences to approximate  $u_{xx}$  &  $u_{yy}$ , we arrive at the approximate equation for the unknown  $u(x_i, y_j)$ , for  $i=1,2,\dots,l-1$  &  $j=1,2,\dots,m-1$ :

$$\frac{-v_{i-1,j} + 2v_{ij} - v_{i+1,j}}{h_x^2} + \frac{-v_{i,j-1} + 2v_{ij} - v_{i,j+1}}{h_y^2} + \sigma v_{ij} = f_{ij}$$

$v_{i,j} = 0: i=0, i=l, j=0, \text{ or } j=m$

- Order the unknowns (& also the vector  $f$ ) lexicographically by y-lines:

$$v = (v_{1,1}, v_{1,2}, \dots, v_{1,m-1}, v_{2,1}, v_{2,2}, \dots, v_{2,m-1}, \dots, v_{l-1,1}, v_{l-1,2}, \dots, v_{l-1,m-1})^T$$

CU-Boulder

61 of 396

## Resulting linear system

We obtain a block-tridiagonal system  $Av = f$ :

$$\begin{pmatrix} A_1 & -I_x & & & \\ -\frac{2}{h_y^2} A_2 & -I_x & & & \\ & -I_x & A_3 & & \\ & & \ddots & \ddots & \\ & & & -I_x & A_{l-2} & -I_x \\ & & & & -I_x & A_{l-1} \end{pmatrix} \begin{pmatrix} v_1 \\ v_2 \\ v_3 \\ \vdots \\ v_{l-2} \\ v_{l-1} \end{pmatrix} = \begin{pmatrix} f_1 \\ f_2 \\ f_3 \\ \vdots \\ f_{l-2} \\ f_{l-1} \end{pmatrix}$$

where  $I_x$  is the  $h_x^2$  times the identity matrix &

$$A_i = \begin{pmatrix} \frac{2}{h_x^2} + \frac{2}{h_y^2} + \sigma & -\frac{1}{h_y^2} & & & \\ -\frac{1}{h_y^2} & \frac{2}{h_x^2} + \frac{2}{h_y^2} + \sigma & -\frac{1}{h_y^2} & & \\ & -\frac{1}{h_y^2} & \frac{2}{h_x^2} + \frac{2}{h_y^2} + \sigma & \ddots & \\ & & \ddots & \ddots & -\frac{1}{h_y^2} \\ & & & & -\frac{1}{h_y^2} & \frac{2}{h_x^2} + \frac{2}{h_y^2} + \sigma & -\frac{1}{h_y^2} \\ & & & & & -\frac{1}{h_y^2} & \frac{2}{h_x^2} + \frac{2}{h_y^2} + \sigma \end{pmatrix}$$

CU-Boulder

62 of 396

## Stencils

preferred for grid issues

Stencils are much better for showing the grid picture:

$$\begin{pmatrix} & & -\frac{1}{h_y^2} & & \\ -\frac{1}{h_x^2} & \frac{2}{h_x^2} + \frac{2}{h_y^2} + \sigma & -\frac{1}{h_x^2} & & \\ & & -\frac{1}{h_y^2} & & \end{pmatrix}$$

dropping the mesh sizes &  $\sigma$

Stencils show local relationships--grid point interactions.

CU-Boulder

63 of 396

## Inhomogeneous boundary conditions

superposition

- Consider a boundary-value problem on domain  $\Omega$  with nonzero data,  $g$ , on the boundary,  $\partial\Omega$ :

$$Lu = f \text{ on } \Omega, \quad Mu = g \text{ on } \partial\Omega. \quad \boxed{\text{Dirichlet: } M = I}$$

- Find a suitable  $w$  satisfying the boundary condition:
 
$$Mw = g \text{ on } \partial\Omega.$$
- Now just find  $z$  to correct  $w$  so that  $w + z = u$ , that is,  $z = u - w$ :
 
$$Lz = f - Lw \equiv \hat{f} \text{ on } \Omega, \quad Mz = g - Mw = 0 \text{ on } \partial\Omega.$$
- Message: Don't look for  $u$ . Instead, look for  $w$  so that  $Mw = g$  on  $\partial\Omega$  & then look for  $z = u - w$  such that  $Lz = f - Lw$  on  $\Omega$  &  $Mz = 0$  on  $\partial\Omega$ .
- In the discrete Dirichlet case, set  $w^h$  to  $g^h$  on  $\partial\Omega$  and  $0$  inside  $\Omega$ .
- So we consider only the homogeneous case from now on.

CU-Boulder

64 of 396

# Outline

## Chapters 1-5:

- Model Problems
- Basic Iterative Methods
  - Convergence tests
  - Analysis
- Elements of Multigrid
  - Relaxation
  - Coarsening
- Implementation
  - Complexity
  - Diagnostics
- Some Theory
  - Spectral vs. algebraic

## Chapters 6-10:

- Nonlinear Problems
  - Full approximation scheme
- Selected Applications
  - Neumann boundaries
  - Anisotropic problems
  - Variable meshes
  - Variable coefficients
- Algebraic Multigrid (AMG)
  - Matrix coarsening
- Multilevel Adaptive Methods
  - FAC
- Finite Elements
  - Variational methodology

Homework Due!

# 2. Basic iterative methods

- Consider the matrix equation  $Au = f$  & let  $v$  be an approximation to  $u$ .
- Two important measures:

The Error:  $e = u - v$  with norms

$$\|e\|_\infty = \max |e_i| \quad \& \quad \|e\|_2 = (\sum e_i^2)^{1/2}.$$

The Residual:  $r = f - Av$  with

$$\|r\|_\infty \quad \& \quad \|r\|_2.$$

What does  $\|r\|$  measure???

Why have both  $r$  &  $e$ ???

## Residual correction

- Note:  $e = u - v \Rightarrow Ae = A(u - v) = f - Av = r$ .

- Residual Equation:

$$Ae = r.$$

What does this do for us?

- Residual Correction:

$$u = v + e.$$

Solve  $Au = f$  with guess  $v$   
or  $Ae = r$  with guess  $0$ .

## Relaxation

- Consider the 1-D model problem

$$-u_{i-1} + 2u_i - u_{i+1} = h^2 f_i \quad 1 \leq i \leq n-1 \quad u_0 = u_n = 0$$

- Jacobi (simultaneous displacement): Solve the  $i^{\text{th}}$  equation for  $v_i$  holding all other variables fixed:

$$v_i^{(new)} = \frac{1}{2} (v_{i-1}^{(old)} + v_{i+1}^{(old)} + h^2 f_i) \quad 1 \leq i \leq n-1$$

## Jacobi in matrix form

- Let  $A = D - L - U$ , where  $D$  is diagonal &  $-L$  &  $-U$  are the strictly lower & upper triangular parts of  $A$ .
- Then  $Au = f$  becomes

$$\begin{aligned} (D - L - U)u &= f \\ Du &= (L + U)u + f \\ u &= D^{-1}(L + U)u + D^{-1}f \end{aligned}$$

- Let  $R_J = D^{-1}(L + U)$ .  $R_J = D^{-1}(D - A) = I - D^{-1}A$   
"Error propagation or iteration matrix".

- Then the iteration is

$$v^{(new)} = R_J v^{(old)} + D^{-1}f$$

## Error propagation matrix & the error

From the derivation,

$$\begin{aligned} u &= D^{-1}(L + U)u + D^{-1}f \\ u &= R_J u + D^{-1}f \end{aligned}$$

the iteration is

$$v^{(new)} = R_J v^{(old)} + D^{-1}f$$

subtracting,

$$u - v^{(new)} = R_J u + D^{-1}f - (R_J v^{(old)} + D^{-1}f)$$

or

$$u - v^{(new)} = R_J u - R_J v^{(old)}$$

hence,

Error propagation!

$$e^{(new)} = R_J e^{(old)}$$

$$R_J = I - D^{-1}A$$

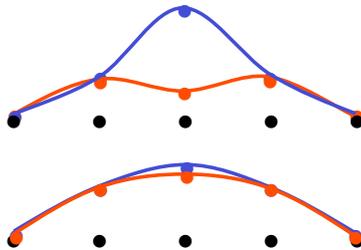
## A picture

1D Poisson

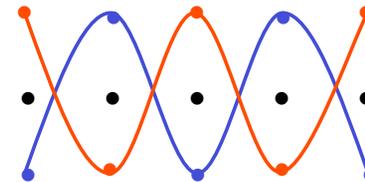
$$R_J = D^{-1}(L + U) = \begin{bmatrix} \frac{1}{2} & 0 & \frac{1}{2} \end{bmatrix}$$

so Jacobi is an error averaging process:

$$e_i^{(new)} \leftarrow (e_{i-1}^{(old)} + e_{i+1}^{(old)})/2$$



## But...



## Another matrix look at Jacobi

$$\begin{aligned} v^{(new)} &\leftarrow D^{-1}(L + U)v^{(old)} + D^{-1}f \quad (L + U = D - A) \\ &= (I - D^{-1}A)v^{(old)} + D^{-1}f \\ v^{(new)} &= v^{(old)} - D^{-1}(Av^{(old)} - f) = v^{(old)} + D^{-1}r \end{aligned}$$

- Exact:  $u = u - D^{-1}(Au - f)$
- Subtracting:  $e^{(new)} = e^{(old)} - D^{-1}Ae^{(old)} = (I - D^{-1}A)e^{(old)}$
- Exact:  $u = u - A^{-1}(Au - f) = A^{-1}f$
- General form:  $u = u - B(Au - f)$  with  $B \sim A^{-1}$
- Damped Jacobi:  $u = u - \omega D^{-1}(Au - f)$  with  $0 < \omega < 2/\rho(D^{-1}A)$
- Gauss-Seidel:  $u = u - (D - L)^{-1}(Au - f)$

Note that  $R_\omega = I - \omega D^{-1}A$  is a polynomial in  $A$  when  $D = I$ .  
 We exploit this simplicity (symmetry, etc.) in what follows.  
 !!!This special property doesn't usually hold in practice!!!

## Weighted Jacobi

safer changes:  $0 < \omega < 2/\rho(D^{-1}A) \approx 1$

- Consider the iteration

$$v_i^{(new)} \leftarrow (1 - \omega)v_i^{(old)} + \frac{\omega}{2}(v_{i-1}^{(old)} + v_{i+1}^{(old)} + h^2 f_i)$$

- Letting  $A = D - L - U$ , the matrix form is

$$\begin{aligned} v^{(new)} &= [(1 - \omega)I + \omega D^{-1}(L + U)]v^{(old)} + \omega h^2 D^{-1}f \\ &= R_\omega v^{(old)} + \omega h^2 D^{-1}f \end{aligned}$$

- Note that

$$R_\omega = [(1 - \omega)I + \omega R_J]$$

- It is easy to see that if  $e^{(approx)} = u - v^{(approx)}$ , then

$$e^{(new)} = R_\omega e^{(old)}$$

## Gauss-Seidel (1-D)

- Solve equation  $i$  for  $u_i$  & update immediately.
- Equivalently: set each component of  $r$  to zero in turn.
- Component form: for  $i = 1, 2, \dots, n-1$ , set

$$v_i \leftarrow \frac{1}{2}(v_{i-1} + v_{i+1} + h^2 f_i)$$

- Matrix form:  $A = (D - L - U)$   
 $(D - L)u = Uu + f$   
 $u = (D - L)^{-1}Uu + (D - L)^{-1}f$
- Let  $R_G = (D - L)^{-1}U$   $R_G = (D - L)^{-1}(D - L - A) = I - (D - L)^{-1}A$

- Then iterate:  $v^{(new)} \leftarrow R_G v^{(old)} + (D - L)^{-1}f$

- Error propagation:  $e^{(new)} \leftarrow R_G e^{(old)}$

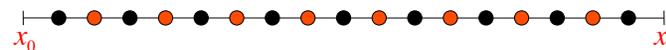
## Red-black Gauss-Seidel

- Update the EVEN points:

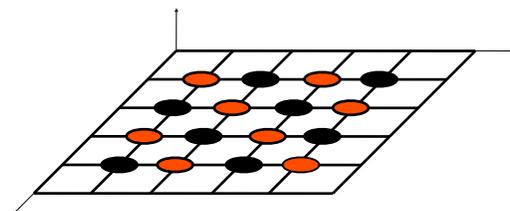
$$v_{2i} \leftarrow \frac{1}{2}(v_{2i-1} + v_{2i+1} + h^2 f_{2i})$$

- Update the ODD points:

$$v_{2i+1} \leftarrow \frac{1}{2}(v_{2i} + v_{2i+2} + h^2 f_{2i+1})$$



- 2-D:



## Test?

$$Au = f$$

Need to know how we're doing!!!

- What  $f$  ?
- What  $v$  ?

$$Au = 0$$

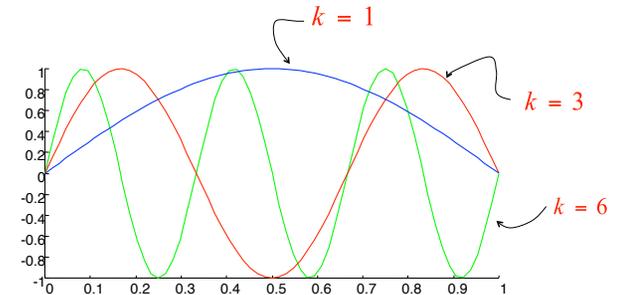
$$v = \text{rand}$$

## Numerical experiments

- Solve  $Au = 0$ ,  $-u_{i-1} + 2u_i - u_{i+1} = 0$
- Use Fourier modes as initial iterates, with  $n = 64$ :

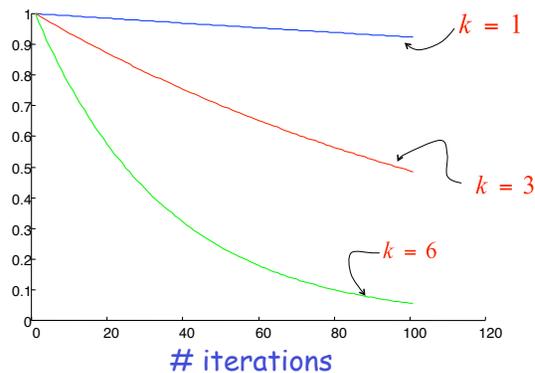
$$v_k = \sin(k\pi x_i), \quad x_i = i/n, \quad 1 \leq i \leq n-1, \quad 1 \leq k \leq n-1$$

component      mode



## Convergence factors differ for different error components

Error,  $\|e\|_\infty$ , in weighted ( $\omega = 2/3$ ) Jacobi on  $Au = 0$  using initial guesses  $v_1, v_3, & v_6$  &  $n = 64$ :



## Stalling convergence

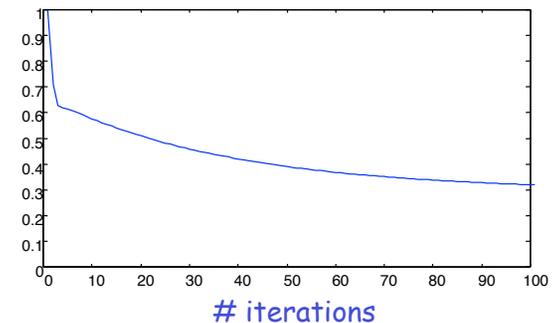
relaxation shoots itself in the foot

- Weighted ( $\omega = 2/3$ ) Jacobi on 1-D problem &  $n = 64$ .

- Initial guess:

$$v = \frac{1}{3} \left[ \left( \sin \left( \frac{1i\pi}{n} \right) \right) + \left( \sin \left( \frac{6i\pi}{n} \right) \right) + \left( \sin \left( \frac{32i\pi}{n} \right) \right) \right]$$

- Error,  $\|e\|_\infty$ , plotted against iteration number:



## Analysis of stationary linear iteration

- The iteration is  $\mathbf{v}^{(new)} = \mathbf{R}\mathbf{v}^{(old)} + \mathbf{g}$ .
- Exact solution doesn't change:  $\mathbf{u} = \mathbf{R}\mathbf{u} + \mathbf{g}$ .
- Subtracting:  $\mathbf{e}^{(new)} = \mathbf{R}\mathbf{e}^{(old)}$ .
- Let  $\mathbf{e}^{(0)}$  be the initial error &  $\mathbf{e}^{(i)}$  be the error after the  $i^{\text{th}}$  iteration. After  $n$  iterations, we have  $\mathbf{e}^{(m)} = \mathbf{R}\mathbf{e}^{(m-1)} = \mathbf{R}^2\mathbf{e}^{(m-2)} = \dots = \mathbf{R}^m\mathbf{e}^{(0)}$ .

We can deal with  $2^4$ , but  $\begin{pmatrix} a & b & c \\ d & e & f \\ g & h & i \end{pmatrix}^4$  ???

What if  $\begin{pmatrix} a & b & c \\ d & e & f \\ g & h & i \end{pmatrix}^4 \mathbf{e} = 2^4 \mathbf{e}$  ???

## Review of eigenvectors & eigenvalues

- The real number  $\lambda$  is an eigenvalue of matrix  $\mathbf{B}$  &  $\mathbf{w} \neq \mathbf{0}$  is its associated eigenvector if  $\mathbf{B}\mathbf{w} = \lambda\mathbf{w}$ . Bold for vectors here temporarily
- The eigenvalues & eigenvectors are characteristics of a given matrix.
- Eigenvectors are linearly independent, & if there is a complete set of  $N$  distinct eigenvectors for an  $N \times N$  matrix, then they form a basis: for any  $\mathbf{v}$ , there exist unique scalars  $v_k$  such that

$$\mathbf{v} = \sum_{k=1}^N v_k \mathbf{w}_k.$$

- Propagation:

$$\mathbf{B}^m \mathbf{v} = \sum_{k=1}^N \lambda^m v_k \mathbf{w}_k.$$

Why is an eigenvector useful???

## "Fundamental Theorem of Iteration"

$\mathbf{R}$  is convergent ( $\mathbf{R}^m \rightarrow \mathbf{0}$  as  $m \rightarrow \infty$ ) iff  $\rho(\mathbf{R}) = \max |\lambda| < 1$ .

Thus,  $\mathbf{v}^{(m)} = \mathbf{R}^m \mathbf{v}^{(0)} \rightarrow \mathbf{0}$  for any initial vector  $\mathbf{v}^{(0)}$  iff  $\rho(\mathbf{R}) < 1$ .

$\rho(\mathbf{R}) < 1$  assures convergence of  $\mathbf{R}$  iteration.

$\rho(\mathbf{R})$  is the spectral convergence factor.

But  $\rho$  doesn't tell you much by itself--it's generally valid only asymptotically. It's useful for the symmetric case in particular because it's equal to  $\|\mathbf{R}\|_2$ , so we'll use it here.

## Rayleigh quotient vs. spectral radius

assume  $\mathbf{A}$  is symmetric ( $\mathbf{w}_k$  orthonormal) & nonnegative definite ( $\lambda \geq 0$ )

- $RQ(\mathbf{v}) \leq \rho(\mathbf{A})$ :  $\mathbf{v} = \sum v_k \mathbf{w}_k$

$$\begin{aligned} RQ(\mathbf{v}) &= \frac{\langle \mathbf{A}\mathbf{v}, \mathbf{v} \rangle}{\langle \mathbf{v}, \mathbf{v} \rangle} = \frac{\langle \mathbf{A} \sum v_k \mathbf{w}_k, \sum v_k \mathbf{w}_k \rangle}{\langle \sum v_k \mathbf{w}_k, \sum v_k \mathbf{w}_k \rangle} \\ &= \frac{\langle \sum \lambda_k v_k \mathbf{w}_k, \sum v_k \mathbf{w}_k \rangle}{\langle \sum v_k \mathbf{w}_k, \sum v_k \mathbf{w}_k \rangle} = \frac{\sum \lambda_k v_k^2}{\sum v_k^2} \leq \lambda_N = \rho(\mathbf{A}) \end{aligned}$$

- $\sup_{\mathbf{v} \neq \mathbf{0}} RQ(\mathbf{v}) = \rho(\mathbf{A})$ :

$$RQ(\mathbf{w}_N) = \frac{\langle \mathbf{A}\mathbf{w}_N, \mathbf{w}_N \rangle}{\langle \mathbf{w}_N, \mathbf{w}_N \rangle} = \frac{\langle \lambda_N \mathbf{w}_N, \mathbf{w}_N \rangle}{\langle \mathbf{w}_N, \mathbf{w}_N \rangle} = \lambda_N = \rho(\mathbf{A})$$

# Euclidean norm vs. spectral radius

use RQ

•  $\|R\|_2 = \rho^{1/2}(R^T R)$ :

$$\begin{aligned} \|R\|_2^2 &= \sup_{e \neq 0} \|Re\|_2^2 / \|e\|_2^2 \\ &= \sup_{e \neq 0} \langle Re, Re \rangle / \langle e, e \rangle \\ &= \sup_{e \neq 0} \langle R^T R e, e \rangle / \langle e, e \rangle = \rho(R^T R) \end{aligned}$$

note:  $\|Re\|_2 \leq \|R\|_2 \cdot \|e\|_2$

•  $\|A\|_2 = \rho^{1/2}(A^2) = \rho(A)$  for symmetric **A!!!**

# $\rho(R)$ vs. $\|R\|_2$

$\rho(R) = \sup |\lambda(R)|$ .

Definition?

$\|R\|_2 = \sup_{e \neq 0} \|Re\|_2 / \|e\|_2 = \rho^{1/2}(R^T R)$ .

Norm independent.

Norm?

Depends on  $\|\cdot\|_2$ .

Asymptotic:

$\rho(R) < 1 \ll \|R\|_2$

means that it will converge someday.

Error bound?

$\|e^{(m+1)}\|_2 \leq \|R\|_2 \|e^{(m)}\|_2$ :  
worst case! Probably pessimistic initially, but sharp sooner or later.

$\rho(R) = \inf \|R\|$   
over all norms.

$\rho(R) = \|R\|_2$   
for symmetric R.

**Example:**  $R = \begin{pmatrix} 0 & K \\ 0 & 0 \end{pmatrix}$ ,  $K \gg 0$  large

$\rho(R) = 0$  but  $\|R\|_2 = \rho^{1/2}(R^T R) = \rho^{1/2} \begin{pmatrix} K^2 & 0 \\ 0 & 0 \end{pmatrix} = K!$

$$e^{(0)} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \Rightarrow \frac{\|e^{(1)}\|_2}{\|e^{(0)}\|_2} = \frac{\|Re^{(0)}\|_2}{\|e^{(0)}\|_2} = K$$

Thus, 1 iteration with  $e^{(0)}$  shows dramatic  $L^2$  divergence!

But  $R^2 = 0$ , so  $e^{(2)} = Re^{(1)} = R^2 e^{(0)} = 0!$

Thus, 2 iterations with  $e^{(0)}$  show complete convergence!

On one hand, this is special ( $\lambda=0$ , large  $K$ ,  $2 \times 2$ ), so this behavior would be more subtle & persistent in general.

On the other, this behavior would vanish for symmetric **R**.

# Convergence factor & rate

• How many iterations are enough to guarantee reduction of the initial error by  $10^{-d}$ ?

$$\frac{\|e^{(m)}\|}{\|e^{(0)}\|} \leq \|R^m\| \leq \|R\|^m \sim 10^{-d}$$

• So, including the asymptotic estimate, we have

$$m \sim \frac{d}{-\log_{10} \|R\|} \text{ or } \frac{d}{-\log_{10} \rho(R)}$$

• Convergence factor =  $\|R\|$  or  $\rho(R)$  error reduction/iterate.

• Convergence rate =  $-\log_{10}(\|R\|)$  or  $-\log_{10}(\rho(R))$  digits/iterate.

# Convergence analysis: Weighted Jacobi

$$R_\omega = (1-\omega)I + \omega D^{-1}(L + U)$$

$$= I - \omega D^{-1}A$$

$$R_\omega = I - \frac{\omega}{2} \begin{pmatrix} 2 & -1 & & & \\ -1 & 2 & -1 & & \\ & -1 & 2 & -1 & \\ & & \ddots & \ddots & \ddots \\ & & & -1 & 2 \end{pmatrix} \quad \text{1-D}$$

For our 1-D model, the eigenvectors of weighted Jacobi  $R_\omega$  & the eigenvectors of  $A$  are the same!

Why???

Special!!!

$$\lambda(R_\omega) = 1 - \frac{\omega}{2} \lambda(A)$$

Remember that  $A$  is without  $h^2$  here!

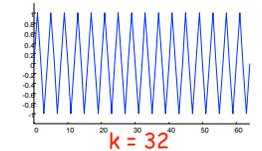
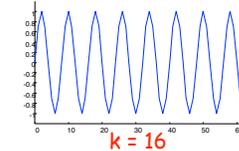
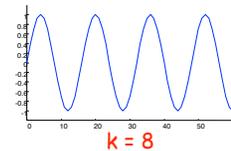
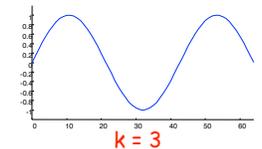
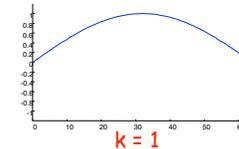
# Eigenpairs of (scaled) $A$

The eigenvectors of  $A$  are (discrete) Fourier modes!

$$\lambda_k(A) = 4\sin^2\left(\frac{k\pi}{2n}\right), \quad w_{k,i} = \sin\left(\frac{ik\pi}{n}\right)$$

$$\begin{bmatrix} -1 & 2 & -1 \\ \lambda_{n-1} \equiv 4 \\ \lambda_1 \equiv \pi^2 h^2 \end{bmatrix}$$

$n = 64$



# Eigenvectors of $R_\omega =$ eigenvectors of $A$

$$\lambda_k(R_\omega) = 1 - 2\omega \sin^2\left(\frac{k\pi}{2n}\right)$$

- Expand the initial error in terms of the eigenvectors:

$$e^{(0)} = \sum_{k=1}^n c_k w_k \leftarrow \begin{matrix} \text{drop bold} \\ \text{for vectors} \end{matrix}$$

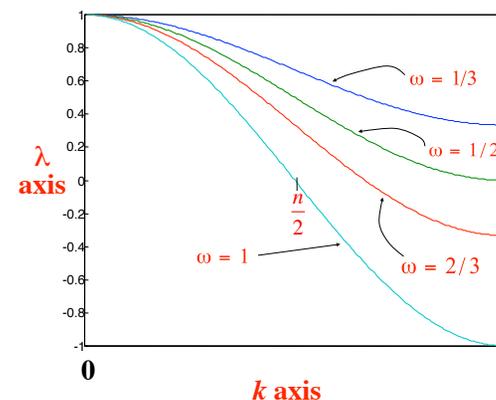
- After  $m$  iterations:

$$R^m e^{(0)} = \sum_{k=1}^n c_k \lambda_k^m w_k$$

- The  $k^{\text{th}}$  error mode is reduced by  $\lambda_k(R_\omega)$  each iteration.

# Relaxation suppresses eigenmodes unevenly

Look carefully at  $\lambda_k(R_\omega) = 1 - 2\omega \sin^2\left(\frac{k\pi}{2n}\right)$ .



Note that if  $0 < \omega \leq 1$ , then  $|\lambda_k(R_\omega)| < 1$  for  $k = 1, 2, \dots, n-1$ .

For  $0 < \omega \leq 1$ ,

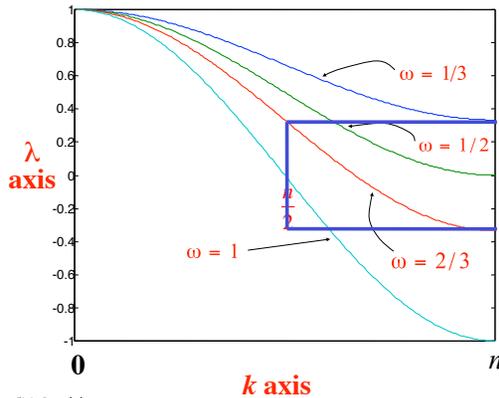
$$\lambda_1 = 1 - 2\omega \sin^2\left(\frac{\pi}{2n}\right)$$

$$= 1 - 2\omega \sin^2\left(\frac{\pi h}{2}\right)$$

$$= 1 - O(h^2) \approx 1$$

# Low frequencies are "undamped"

Notice that no value of  $\omega$  will efficiently damp out long waves or low frequencies.



What value of  $\omega$  gives the best damping of short waves or high frequencies  $n/2 \leq k \leq n-1$ ?  
Choose  $\omega$  such that  $\lambda_{N/2}(R_\omega) = -\lambda_N(R_\omega)$

$$\Rightarrow \omega \approx \frac{2}{3}$$

$$\text{For 2D: } \omega \approx \frac{4}{5}$$

# Smoothing factor

- The smoothing factor is the largest magnitude of the iteration matrix eigenvalues corresponding to the oscillatory Fourier modes:  
smoothing factor =  $\max |\lambda_k(R)|$  for  $n/2 \leq k \leq n-1$ .

"MG" spectral radius?

- Why only the upper spectrum?

- For  $R_\omega$  with  $\omega = 2/3$ , the smoothing factor is  $1/3$ :

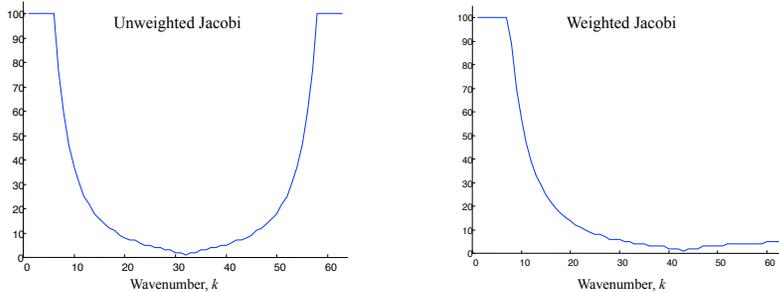
$$|\lambda_{n/2}| = |\lambda_{n-1}| = 1/3 \text{ \& } |\lambda_k| < 1/3 \text{ for } n/2 < k < n-1.$$



- But  $|\lambda_k| \approx 1 - \omega k^2 \pi^2 h^2$  for long waves ( $k \ll n/2$ ).



# Convergence of Jacobi on $Au = 0$



- Jacobi on  $Au = 0$  with  $n = 64$ . Number of iterations needed to reduce initial error  $\|e\|_\infty$  by 0.01.

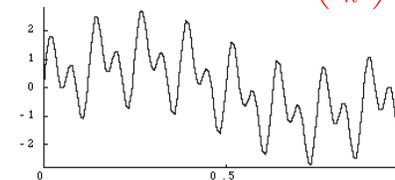
- Initial guess :

$$v_{k,i} = \sin\left(\frac{ik\pi}{n}\right)$$

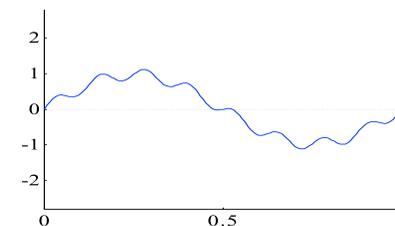
# Weighted Jacobi = smoother (error)

- Initial error:  $v_i = \sin\left(\frac{2i\pi}{n}\right) + \frac{1}{2}\sin\left(\frac{16i\pi}{n}\right) + \frac{1}{2}\sin\left(\frac{32i\pi}{n}\right)$

not solution or approximation!!!



- Error after 35 iteration sweeps:



Many relaxation schemes are smoothers: oscillatory error modes are quickly eliminated, but smooth modes are slowly damped.

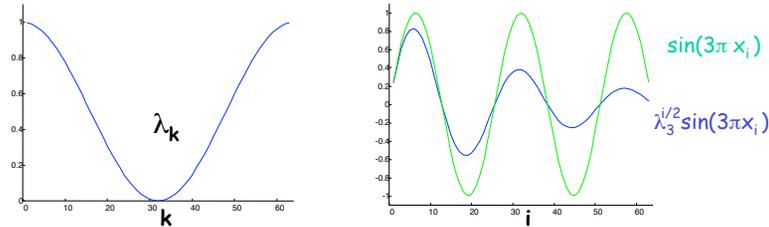
# Similar analysis for other smoothers

- Gauss-Seidel relaxation applied to the 3-point difference matrix  $A$  (1-D model problem):

$$R_G = (D - L)^{-1} U.$$

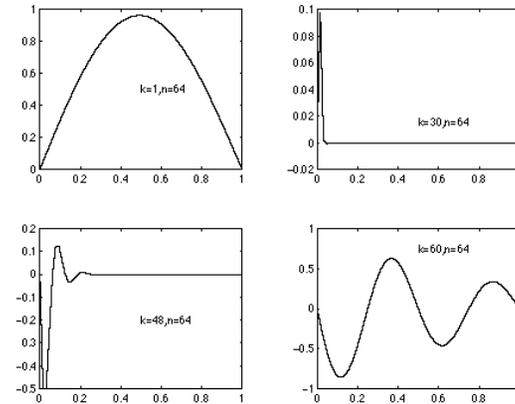
- A little algebra & trigonometry shows that

$$\lambda_k(R_G) = \cos^2\left(\frac{k\pi}{n}\right) \quad (w_k)_i = \cos^i\left(\frac{k\pi}{n}\right) \sin\left(\frac{ik\pi}{n}\right) = \lambda_k^{i/2} \sin(k\pi x_i)$$



What's  $w_k$  look like for large  $k$ ?

# Gauss-Seidel eigenvectors



These are VERY different from Jacobi's eigenvectors. It's not clear how smoothness depends on  $k$ . You cannot expect G-S to quickly reduce Fourier modes. You can only hope that G-S produces smooth results!

# Gauss-Seidel convergence

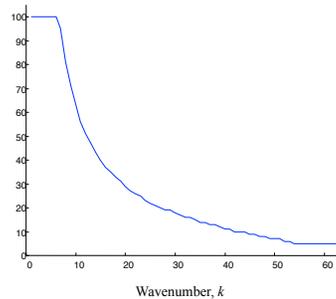
Eigenvectors of  $R_G$  are not the same as those of  $A$ !!!

Gauss-Seidel mixes the modes of  $A$ .

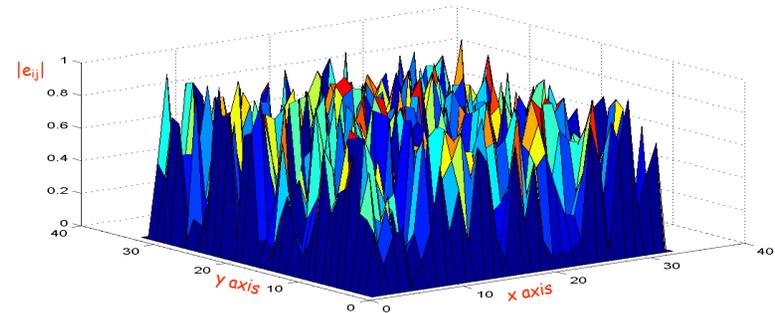
Gauss-Seidel on  $Au = 0$ , with  $n = 64$ . Number of iterations needed to reduce initial error  $\|e\|_\infty$  by 0.01.

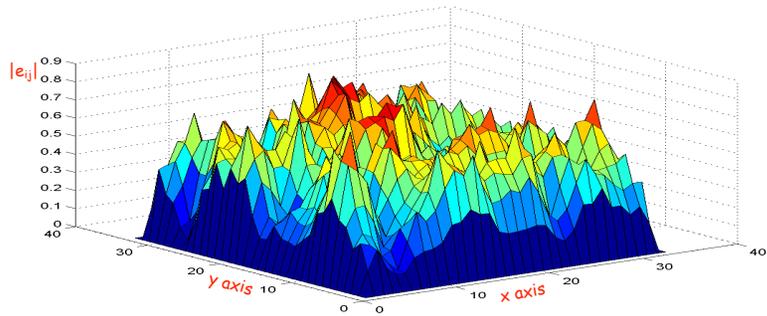
Initial guess (modes of  $A$ ):

$$v_{ki} = \sin\left(\frac{ik\pi}{n}\right)$$



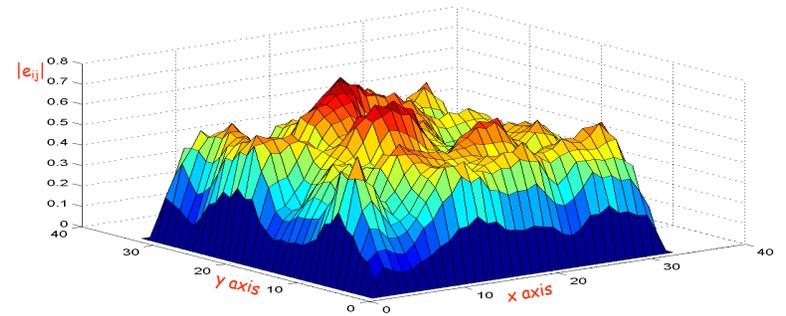
So G-S **does** reduce oscillatory Fourier modes.





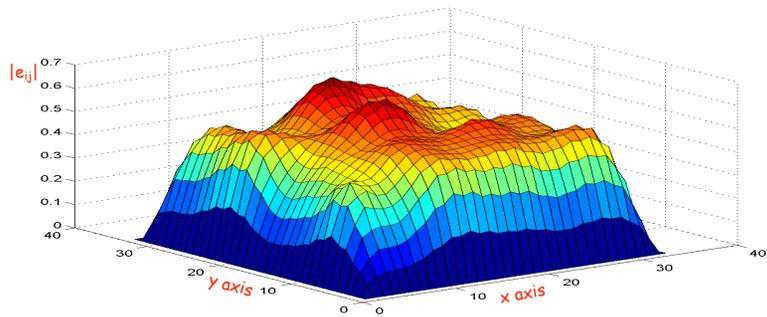
CU-Boulder

101 of 396



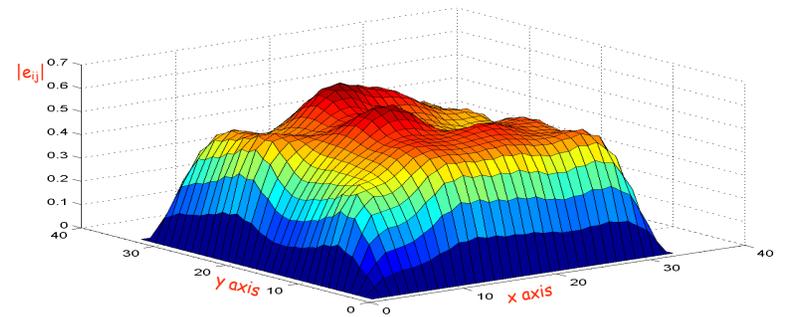
CU-Boulder

102 of 396



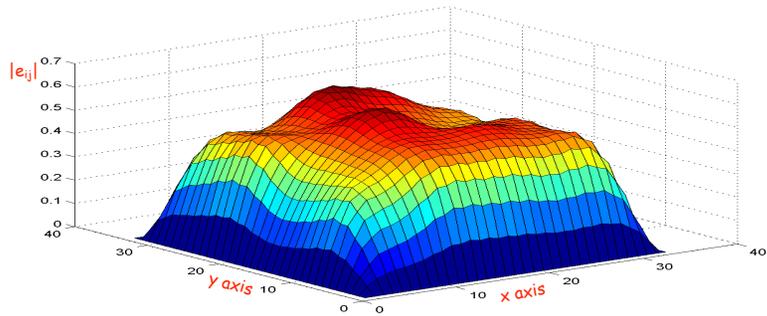
CU-Boulder

103 of 396



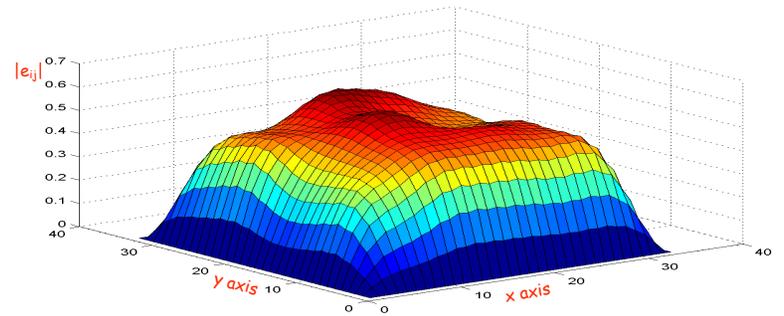
CU-Boulder

104 of 396



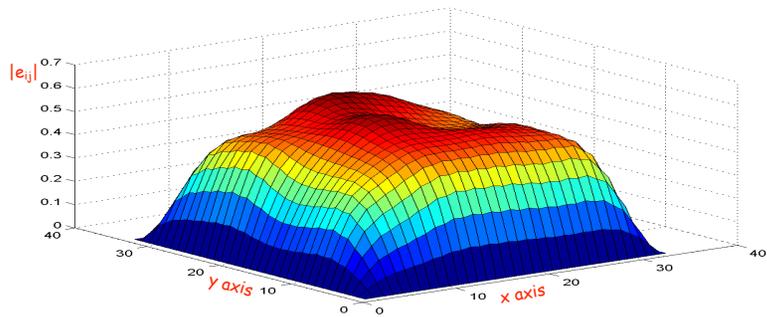
CU-Boulder

105 of 396



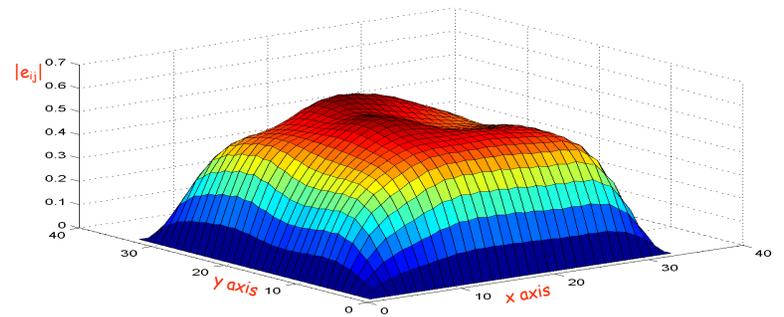
CU-Boulder

106 of 396



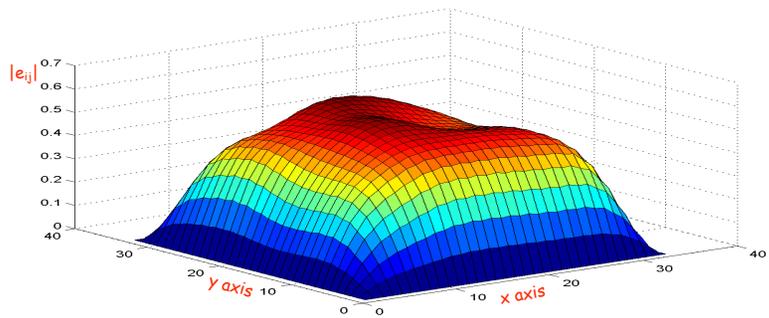
CU-Boulder

107 of 396



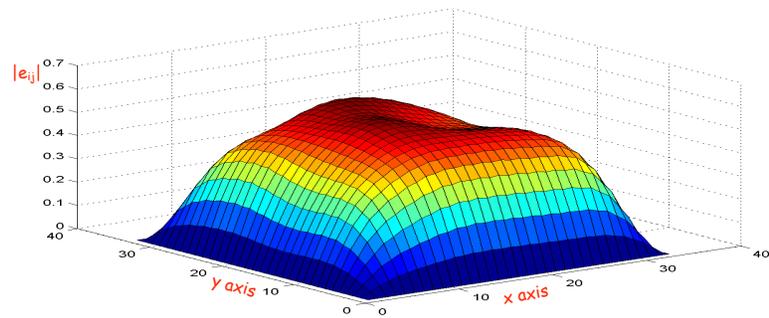
CU-Boulder

108 of 396



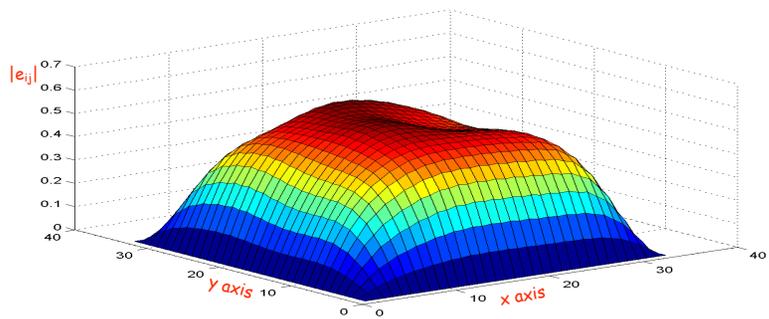
CU-Boulder

109 of 396



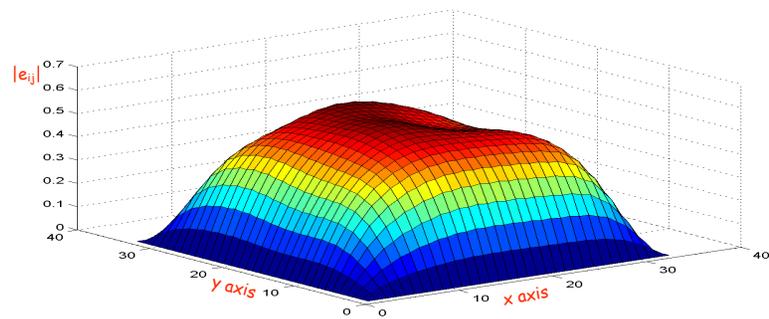
CU-Boulder

110 of 396



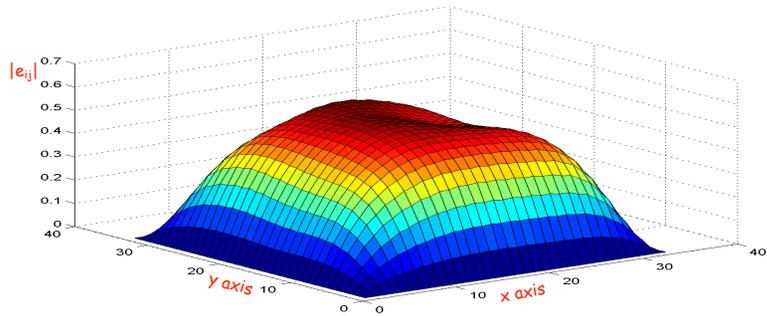
CU-Boulder

111 of 396



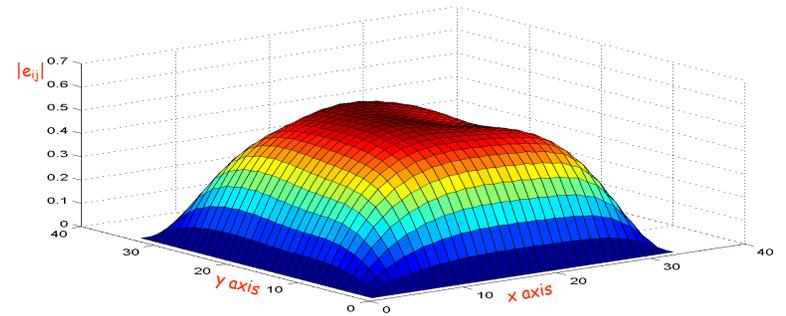
CU-Boulder

112 of 396



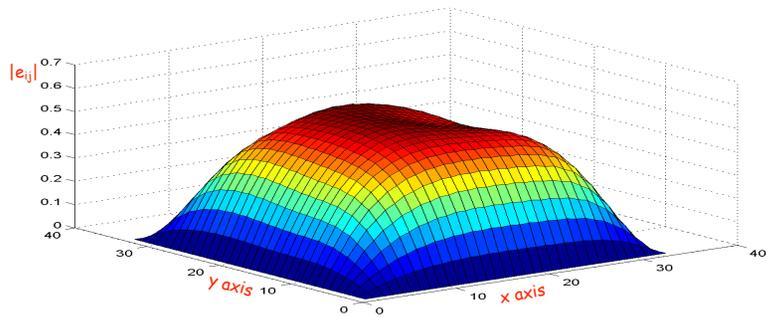
CU-Boulder

113 of 396



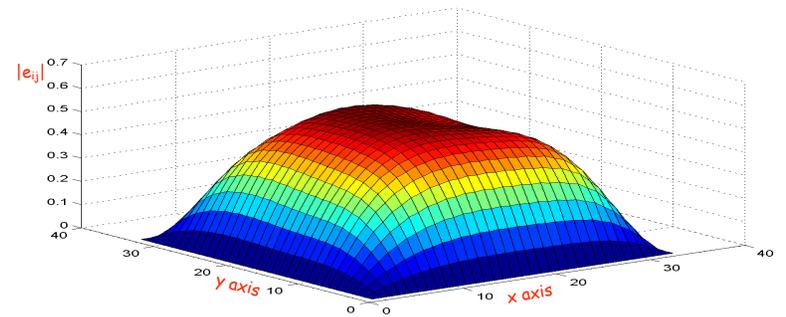
CU-Boulder

114 of 396



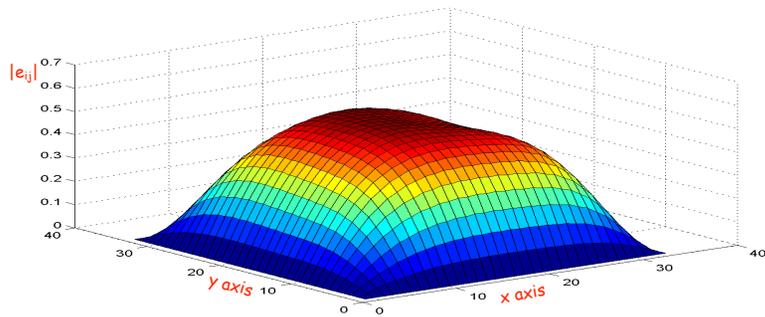
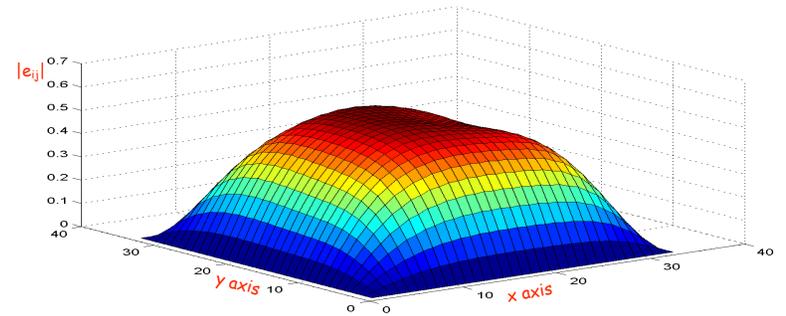
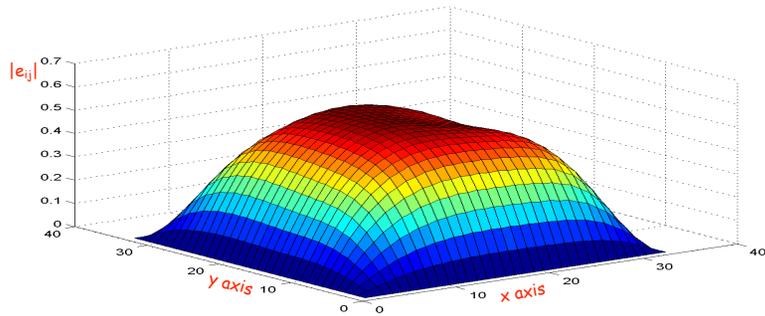
CU-Boulder

115 of 396



CU-Boulder

116 of 396



## Outline

### Chapters 1-5:

- $\sqrt{\quad}$  Model Problems
- $\sqrt{\quad}$  Basic Iterative Methods
  - Convergence tests
  - Analysis
- Elements of Multigrid
  - Relaxation
  - Coarsening
- Implementation
  - Complexity
  - Diagnostics
- Some Theory
  - Spectral vs. algebraic

### Chapters 6-10:

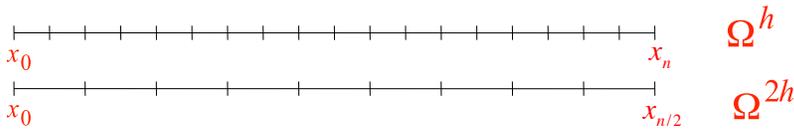
- Nonlinear Problems
  - Full approximation scheme
- Selected Applications
  - Neumann boundaries
  - Anisotropic problems
  - Variable meshes
  - Variable coefficients
- Algebraic Multigrid (AMG)
  - Matrix coarsening
- Multilevel Adaptive Methods
  - FAC
- Finite Elements
  - Variational methodology

Homework Due!

### 3. Elements of multigrid

#### 1<sup>st</sup> observation toward multigrid

- Many relaxation schemes have the smoothing property: oscillatory error modes are quickly eliminated, while smooth modes are often very slow to disappear.
- We'll turn this adversity around: the idea is to use coarse grids to take advantage of smoothing.



How?

### Reason #1 for coarse grids: Nested iteration

- Coarse grids can be used to compute an improved initial guess for the fine-grid relaxation. This is advantageous because:
  - Relaxation on the coarse-grid is much cheaper: half as many points in 1-D, one-fourth in 2-D, one-eighth in 3-D,...
  - Relaxation on the coarse grid has a marginally faster convergence factor ( $|\lambda_1(R)| \approx 1 - \omega\pi^2h^2$ ):

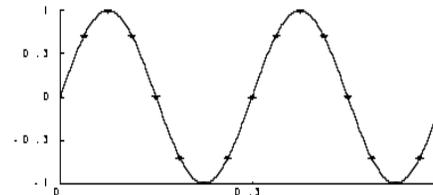
$$1 - O(4h^2) \text{ instead of } 1 - O(h^2).$$

### Idea! Nested iteration

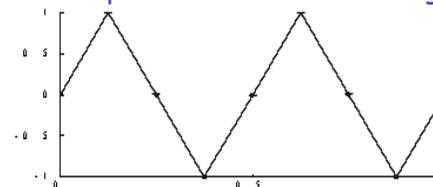
- Relax on  $Au = f$  on  $\Omega^{4h}$  to obtain initial guess  $v^{2h}$ .
- Relax on  $Au = f$  on  $\Omega^{2h}$  to obtain initial guess  $v^h$ .
- Relax on  $Au = f$  on  $\Omega^h$  to obtain ... final solution???
- What is  $A^{2h}u^{2h} = f^{2h}$ ? Analogous to  $A^hu^h = f^h$  for now.
- How do we migrate between grids? Hang on...
- What if the error still has large smooth components when we get to the fine grid  $\Omega^h$ ? Hang on...

### Reason #2 for coarse grids:

- A smooth function:



can be represented by linear interpolation from a coarser grid:



On the coarse grid, smooth error appears to be relatively higher in frequency: in this example, it's the 4-mode out of a possible 15 on the fine grid,  $\sim 1/4$  the way up the spectrum. On the coarse grid, it's the 4-mode out of a possible 7,  $\sim 1/2$  the way up the spectrum.

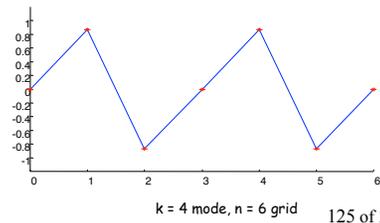
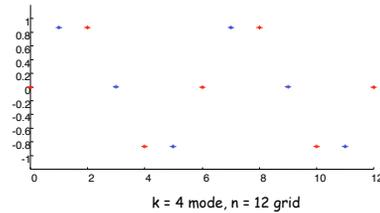
Relaxation on  $2h$  is cheaper & faster on this mode!!!

We need to be careful when we see  $w_k^h$  on the coarse grid  $2h$  and  $w_k^{2h}$  on the fine grid  $h$ .  
 When  $k > n/2$ ,  $w_k^h$  is disguised on the coarse grid: **aliasing!!!**

$$w_{k,2i}^h = \sin\left(\frac{2ik\pi}{n}\right) = \sin\left(\frac{ik\pi}{n/2}\right) = w_{k,i}^{2h}$$

Also, note that  $w_{n/2}^h \rightarrow 0$  on the coarse grid.

What happens to the modes  $k > n/2$ ?

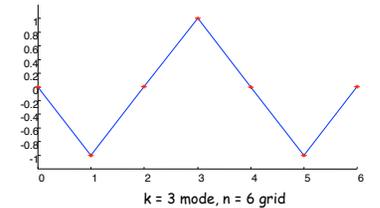
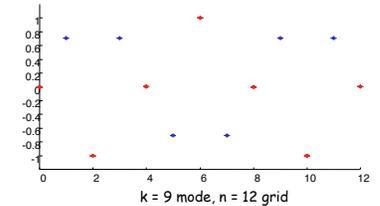


125 of 396

For  $k > n/2$ ,  $w_k^h$  is disguised on the coarse grid: **aliasing!!!**

For  $k > n/2$ , the  $k^{\text{th}}$  mode on the fine grid is aliased & appears as the  $(n - k)^{\text{th}}$  mode on the coarse grid:

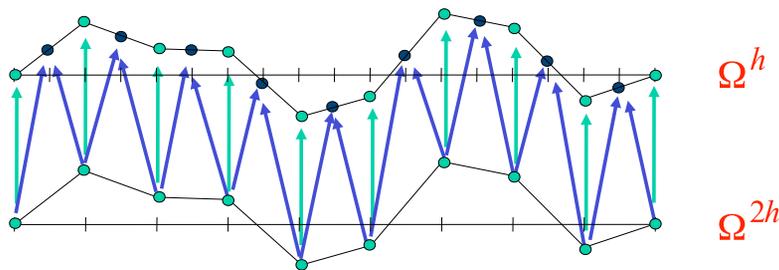
$$\begin{aligned} (w_k^h)_{2i} &= \sin\left(\frac{(2i)\pi k}{n}\right) \\ &= -\sin\left(\frac{2i\pi(n-k)}{n}\right) \\ &= -\sin\left(\frac{i\pi(n-k)}{n/2}\right) \\ &= -(w_{n-k}^{2h})_i \end{aligned}$$



126 of 396

More precisely, let's relate  $1$ -D interpolation (prolongation) grid  $2h$  functions to grid  $h$  functions...

- Values at points on the coarse grid map unchanged to the fine grid.
- Values at fine-grid points NOT on the coarse grid are the averages of their coarse-grid neighbors.



We will often identify  $\Omega^{2h}$  with a subset of  $\Omega^h$ .

127 of 396

## 1-D interpolation (prolongation)

to migrate from coarse to fine grids

- Mapping from the coarse grid to the fine grid:

$$I_{2h}^h : \Omega^{2h} \rightarrow \Omega^h \quad (\Omega^h = \mathbb{R}^{n-1})$$

- Let  $v^h, v^{2h}$  be defined on  $\Omega^h, \Omega^{2h}$ . Then

$$I_{2h}^h v^{2h} = v^h$$

where

$$v_{2i}^h = v_i^{2h} \quad \text{for } 0 \leq i \leq \frac{n}{2} \text{ (including boundaries),}$$

$$v_{2i+1}^h = \frac{1}{2}(v_i^{2h} + v_{i+1}^{2h}) \quad \text{for } 0 \leq i \leq \frac{n}{2} - 1.$$

128 of 396



## Where do we stand?

	smooth components	oscillatory components
relaxation		
nested iteration		

*Note: A red arrow points from the heart in the 'relaxation' row to the heart in the 'nested iteration' row.*

CU-Boulder

133 of 396

## The Key Step to Multigrid

- If what we want to compute is **smooth**, a coarse-grid interpolant **could** do very well.
- If what we want to compute is **oscillatory**, a coarse-grid interpolant **cannot** do very well.
- What if **u** is not smooth? Can we make it so?
- Can we make something smooth? →error←
- Can we smooth **e**? Can we get **e** & use it to get **u**?  
 $Ae = r$  &  $u \leftarrow v + e$ !
- So, use **nested iteration** on the **residual equation** to approximate the error after smoothing!!!
- Just because the coarse grid **can** approximate **e** well doesn't mean we know **how** to do it! But we will soon!

CU-Boulder

134 of 396

## 2<sup>nd</sup> observation toward multigrid

- The residual equation: Let **v** be an approximation to the solution of  $Au = f$ , where the residual  $r = f - Av$ . Then the error  $e = u - v$  satisfies  $Ae = r$ .
- After relaxing on  $Au = f$  on the fine grid, **e** will be smooth, so the coarse grid can approximate **e** well. This will be cheaper & **e** should be more oscillatory there, so relaxation will be more effective.
- Therefore, we go to a coarse grid & relax on the residual equation  $Ae = r$ .

What's a good initial guess on grid  $2h$ ?  $e = 0$ !

How do we get to grid  $2h$ ? Stay tuned...

CU-Boulder

135 of 396

## Idea! Coarse-grid correction 2-grid

- Relax on  $Au = f$  on  $\Omega^h$  to get an approximation  $v^h$ .
- Compute  $r = f - Av^h$ .
- Transfer  $Ae = r$  to  $\Omega^{2h}$  somehow & relax on it to obtain an approximation to the error,  $e^{2h}$ .
- Correct the approximation  $v^h \leftarrow v^h + I_{2h}^h e^{2h}$ .

This is the essence of multigrid.

We need a way to transfer  $Ae = r$  to  $\Omega^{2h}$ .

CU-Boulder

136 of 396

## A way to coarsen $Ae = r$ Galerkin

- Assume we've relaxed so much that  $e$  is smooth.  $P$  when grids understood, else  $I_{2h}^h$ .
- Ansatz:  $e = Pv^{2h}$  for some coarse-grid  $v^{2h}$ .
- How do we characterize  $e$  so we can hope to compute it?

$$Ae = r \quad \Rightarrow \quad A P v^{2h} = r$$

$7 \times 7 \quad 7 \times 3 \quad 3 \times 1 \quad = \quad 7 \times 1$

- Too many equations now & too few unknowns!
- How about just eliminating every other equation?
- How about multiplying both sides by some  $3 \times 7$  matrix?

$$\begin{pmatrix} P^T & A^{2h} & P^T \\ P^T & A & P \end{pmatrix} v^{2h} = P^T r$$

$3 \times 7 \quad 7 \times 7 \quad 7 \times 3 \quad 3 \times 1 \quad = \quad 3 \times 1$

We might write  $R$  instead of  $P^T$  or maybe  $I_h^{2h}$ .

CU-Boulder

137 of 396

## 1-D restriction by injection

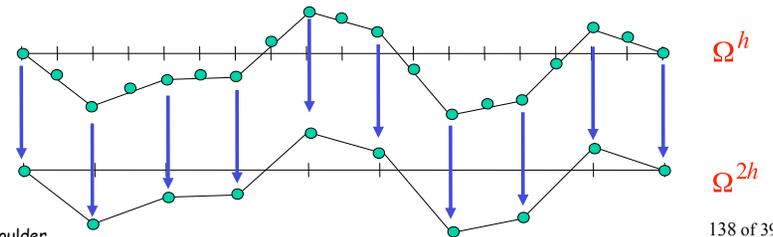
- Mapping from the fine grid to the coarse grid:

$$R = I_h^{2h} : \Omega^h \rightarrow \Omega^{2h}. \quad R \text{ is not } P^T \text{ here!!!}$$

- Let  $v^h, v^{2h}$  be defined on  $\Omega^h, \Omega^{2h}$ . Then

$$Rv^h = I_h^{2h}v^h = v^{2h},$$

where  $v_1^{2h} = v_{2i}^h$ .



CU-Boulder

138 of 396

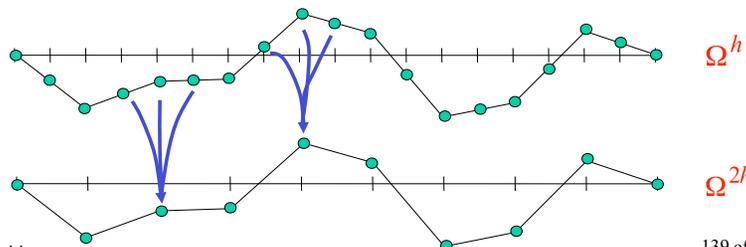
## 1-D restriction by full weighting

- Let  $v^h, v^{2h}$  be defined on  $\Omega^h, \Omega^{2h}$ . Then

$$Rv^h = I_h^{2h}v^h = v^{2h}, \quad R \text{ is } cP^T \text{ here!!!}$$

where

$$v_i^{2h} = \frac{1}{4}(v_{2i-1}^h + 2v_{2i}^h + v_{2i+1}^h).$$



CU-Boulder

139 of 396

## 1-D restriction (full-weighting)

- $R = I_h^{2h}$  is a linear operator:  $\mathfrak{R}^{n-1} \rightarrow \mathfrak{R}^{n/2-1}$ .

Don't confuse  $R$  here with error propagator notation.

- $n = 8$ :

$$\begin{pmatrix} \frac{1}{4} & \frac{1}{2} & \frac{1}{4} & & & & & \\ & \frac{1}{4} & \frac{1}{2} & \frac{1}{4} & & & & \\ & & \frac{1}{4} & \frac{1}{2} & \frac{1}{4} & & & \\ & & & \frac{1}{4} & \frac{1}{2} & \frac{1}{4} & & \\ & & & & \frac{1}{4} & \frac{1}{2} & \frac{1}{4} & \\ & & & & & \frac{1}{4} & \frac{1}{2} & \frac{1}{4} \\ & & & & & & \frac{1}{4} & \frac{1}{2} \\ & & & & & & & \frac{1}{4} \end{pmatrix}_{3 \times 7} \begin{pmatrix} v_1^h \\ v_2^h \\ v_3^h \\ v_4^h \\ v_5^h \\ v_6^h \\ v_7^h \end{pmatrix}_{7 \times 1} = \begin{pmatrix} v_1^{2h} \\ v_2^{2h} \\ v_3^{2h} \end{pmatrix}_{3 \times 1}.$$

- $I_h^{2h}$  has rank  $\frac{n}{2}-1$  because  $\dim(\text{Range}(R)) = \frac{n}{2}-1$ .

Look at the columns of  $R$  associated with grid  $2h$ .

CU-Boulder

140 of 396

# Prolongation & restriction are often nicely related

- For the 1-D examples, linear interpolation & full weighting are

$$I_{2h}^h = \frac{1}{2} \begin{pmatrix} 1 & & & & \\ 2 & & & & \\ 1 & 1 & & & \\ & 2 & & & \\ & 1 & 1 & & \\ & & 2 & & \\ & & 1 & 2 & \\ & & & 1 & 2 & 1 \\ & & & & 1 & 2 & 1 \end{pmatrix}, \quad I_h^{2h} = \frac{1}{4} \begin{pmatrix} 1 & 2 & 1 & & & & \\ & 1 & 2 & 1 & & & \\ & & 1 & 2 & 1 & & \\ & & & 1 & 2 & 1 & \\ & & & & 1 & 2 & 1 \\ & & & & & 1 & 2 & 1 \\ & & & & & & 1 & 2 & 1 \end{pmatrix}$$

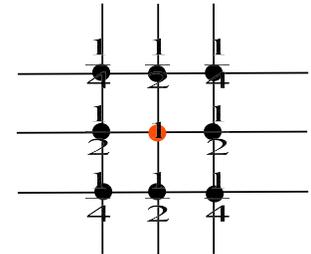
- So they're related by the variational condition

$$I_{2h}^h = c (I_h^{2h})^T, \quad c \text{ in } \mathfrak{R}, \quad P = cR^T$$

# 2-D prolongation

$$\begin{bmatrix} v_{2i,2j}^h = v_{ij}^{2h} \\ v_{2i+1,2j}^h = \frac{1}{2}(v_{ij}^{2h} + v_{i+1,j}^h) \\ v_{2i,2j+1}^h = \frac{1}{2}(v_{ij}^{2h} + v_{i,j+1}^h) \\ v_{2i+1,2j+1}^h = \frac{1}{4}(v_{ij}^{2h} + v_{i+1,j}^h + v_{i,j+1}^h + v_{i+1,j+1}^h) \end{bmatrix} \begin{bmatrix} \frac{1}{4} & \frac{1}{2} & \frac{1}{4} \\ \frac{1}{2} & 1 & \frac{1}{2} \\ \frac{1}{4} & \frac{1}{2} & \frac{1}{4} \end{bmatrix}$$

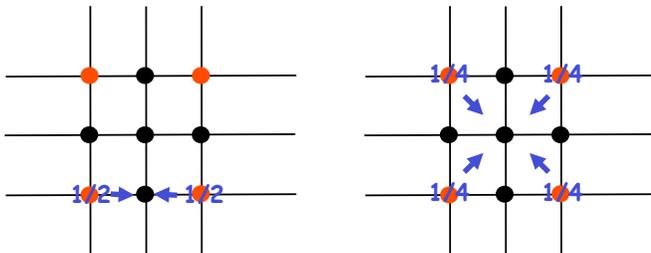
We denote the operator by using a "scatter" stencil  $\begin{bmatrix} \cdot \\ \cdot \\ \cdot \end{bmatrix}$ . Centered over a C-point  $\bullet$ , it shows what fraction of the C-point's value contributes to a neighboring F-point  $\bullet$ .



# "Gather" interpolation stencil

$$\begin{bmatrix} 1/2 & & & & \\ & 1/2 & & & \\ & & 1/4 & & 1/4 \\ & & & 1/4 & & 1/4 \end{bmatrix}$$

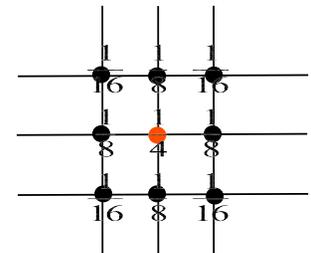
Centered over a fine-grid point  $\bullet$ .



# 2-D restriction (full weighting)

$$\begin{bmatrix} \frac{1}{16} & \frac{1}{8} & \frac{1}{16} \\ \frac{1}{8} & \frac{1}{4} & \frac{1}{8} \\ \frac{1}{16} & \frac{1}{8} & \frac{1}{16} \end{bmatrix}$$

We denote the operator by using a "gather" stencil  $\begin{bmatrix} \cdot \\ \cdot \\ \cdot \end{bmatrix}$ . Centered over a C-point  $\bullet$ , it shows what fraction of the value of the neighboring F-point  $\bullet$  contributes to the value at the C-point.



# Now we put all these ideas together

- Nested Iteration (Relaxation on Coarse Grids)
  - effective on smooth solution (components).
- Relaxation on Fine Grid
  - effective on oscillatory error (components).
- Residual Equation on Fine Grid
  - characterizes the error.
  - enables nested iteration for smooth error (components)!!!
- Prolongation (variables) & Restriction (equations)
  - provides pathways between coarse & fine grids.

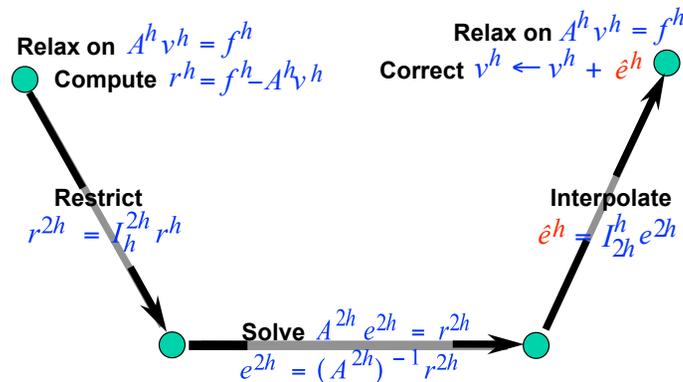
# 2-grid coarse-grid correction

$$v^h \leftarrow CG(v^h, f^h)$$

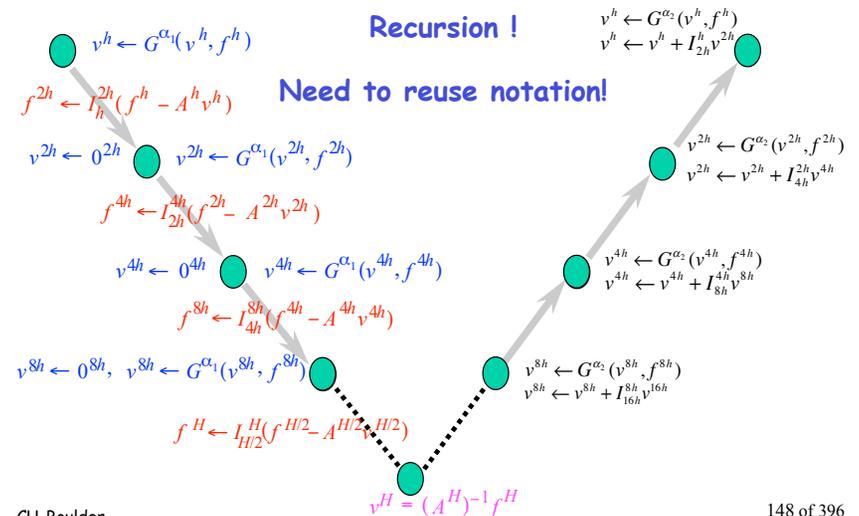
- 1) Relax  $\alpha_1$  times on  $A^h v^h = f^h$  on  $\Omega^h$  with arbitrary initial guess  $v^h$ . If  $h = h_{\text{coarsest}}$ , then go to 6.
- 2) Compute  $r^h = f^h - A^h v^h$ .  $A^{2h} = I_h^{2h} A^h I_h^h = \text{RAP}$  (Galerkin) or direct discretization
- 3) Compute  $r^{2h} = I_h^{2h} r^h$ .
- 4) "Solve"  $A^{2h} e^{2h} = r^{2h}$  on  $\Omega^{2h}$ .
- 5) Correct fine-grid solution  $v^h \leftarrow v^h + I_{2h}^h e^{2h}$ .
- 6) Relax  $\alpha_2$  times on  $A^h v^h = f^h$  on  $\Omega^h$ .

What does  $e^{2h}$  represent here?

# 2-grid coarse-grid correction



# How do we "solve" $A^{2h} e^{2h} = r^{2h}$ ?



# V-cycle (recursive form)

$$v^h \leftarrow MV^h(v^h, f^h)$$

1) Relax  $\alpha_1$  times on  $A^h u^h = f^h$  with initial  $v^h$  given.

2) If  $\Omega^h$  is the coarsest grid, go to 4;

else:  $f^{2h} \leftarrow I_h^{2h}(f^h - A^h v^h)$

$$v^{2h} \leftarrow 0$$

$$v^{2h} \leftarrow MV^{2h}(v^{2h}, f^{2h})$$

3) Correct:  $v^h \leftarrow v^h + I_{2h}^h v^{2h}$ .

4) Relax  $\alpha_2$  times on  $A^h u^h = f^h$  with initial guess  $v^h$ .

# Outline

## Chapters 1-5:

- √ Model Problems
- √ Basic Iterative Methods
  - Convergence tests
  - Analysis
- √ Elements of Multigrid
  - Relaxation
  - Coarsening
- Implementation
  - Complexity
  - Diagnostics
- Some Theory
  - Spectral vs. algebraic

## Chapters 6-10:

- Nonlinear Problems
  - Full approximation scheme
- Selected Applications
  - Neumann boundaries
  - Anisotropic problems
  - Variable meshes
  - Variable coefficients
- Algebraic Multigrid (AMG)
  - Matrix coarsening
- Multilevel Adaptive Methods
  - FAC
- Finite Elements
  - Variational methodology

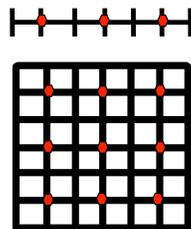
Homework Due!

## 4. Implementation

Storage cost:  $v^h$  &  $f^h$  on each level.

Estimates are approximate ( $n, d, \dots$ ).

- In 1-D, a coarse grid has about half as many points as the fine grid.
- In 2-D, a coarse grid has about one-fourth as many points as the fine grid.
- In  $d$ -dimensions, a coarse grid has about  $2^{-d}$  as many points as the fine grid.



- Total storage cost:  $n^d (1 + 2^{-d} + 2^{-2d} + 2^{-3d} + \dots + 2^{-md}) < \frac{n^d}{1 - 2^{-d}}$  less than 2, 4/3, & 8/7 the cost of storage on the fine grid for 1-D, 2-D, & 3-D problems, respectively.

## Computational cost

- Let one Work Unit (WU) be the cost of one relaxation sweep on the fine grid.
- Ignore the cost of restriction & interpolation (typically about 20% of the total cost).
- Consider a V-cycle with 2 pre-coarse-grid correction sweep ( $\alpha_1 = 2$ ) & 1 post-coarse-grid correction sweep ( $\alpha_2 = 1$ ).
- Cost of a V-cycle (in WUs):

$$3(1 + 2^{-d} + 2^{-2d} + 2^{-3d} + \dots + 2^{-md}) < \frac{3}{1 - 2^{-d}}$$

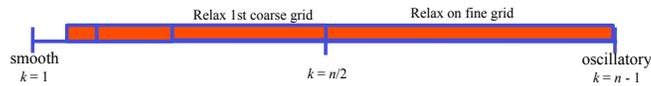
- Cost is about 2, 4/3, & 8/7 X 3 WUs per V-cycle in 1, 2, & 3 dimensions, respectively.

# Convergence analysis

- First, a heuristic argument:
  - The convergence factor for the oscillatory error modes (smoothing factor) is small & bounded uniformly in  $h$ .

smoothing factor =  $\max |\lambda_k(R)|$  for  $n/2 \leq k \leq n-1$ .

- Multigrid focuses the relaxation process on attenuating the oscillatory components on each level.



⇒ The overall multigrid convergence factor is small & bounded uniformly in  $h$ !

Bounded uniformly in  $h \neq$  independent of  $h$ .

# Reminder: Approximate $u''(x)$ via Taylor series

$O(h^2)$  means a quantity bounded in norm by  $Ch^2$  for some constant  $C$ .

- The BVP:
- Approximate 2<sup>nd</sup> derivative using Taylor series:

$$u(0) = u(1) = 0$$

$$u(x_{i+1}) = u(x_i) + h u'(x_i) + \frac{h^2}{2!} u''(x_i) + \frac{h^3}{3!} u'''(x_i) + O(h^4)$$

The finite difference scheme:

$$u(x_{i+1}) - 2u(x_i) + u(x_{i-1})) = \frac{h^2}{2!} u''(x_i) - \frac{h^4}{24} u^{(4)}(x_i) + O(h^6)$$

- Summing & solving:
- Truncation error:

$$\frac{u''(x_i) + 2u''(x_{i+1}) - u''(x_{i-1}))}{h^2} + O(h^2) = f_i \quad i = 1, 2, \dots, n-1$$

# Actual error

$A$  has  $h^{-2}$  in it here

$$Au^{(h)} + O(h^2) = f \Rightarrow Au^{(h)} = f + O(h^2)$$

$$\Rightarrow Au^h = f \quad (u^h = \text{discrete sol'n})$$

$$\Rightarrow A(u^{(h)} - u^h) = O(h^2)$$

$$\Rightarrow AE = O(h^2). \leftarrow \text{consistency}$$

( $E = \text{discretization error}$ )

So  $\|E\| = \|A^{-1} O(h^2)\|$

$$\leq \|A^{-1}\| \cdot \|O(h^2)\|$$

$$= \lambda_{\max}(A^{-1}) \cdot O(h^2) = O(h^2) / \lambda_{\min}(A) \sim O(h^2) / \tau^2$$

or  $\|E\| = O(h^2). \leftarrow \text{convergence}$

↑  
stability

# Overall goal of computation

- Continuous problem:  $Au = f, u_i = u(x_i)$
- Discrete problem:  $A^h u^h = f^h, v^h \approx u^h$
- Global/discretization error:  $E_i = u(x_i) - u_i^h$   
 $\|E\| \leq Kh^p$   
 ( $p = 2$  for model problem & proper norm)
- Algebraic error:  $e_i^h = u_i^h - v_i^h$
- For tolerance  $\epsilon$ , assume the aim is to find  $v^h$  so that the total error,  $\|e\| = \|u^{(h)} - v^h\| \leq \epsilon$ , where  $u^{(h)} = (u(x_i))$ .
- Then this objective is assured as follows:  
 $\|e\| \leq \|u^{(h)} - u^h\| + \|u^h - v^h\| = \|E\| + \|e^h\| \leq \epsilon.$

## We can satisfy the convergence objective by imposing two conditions

- 1)  $\|E\| \leq \epsilon/2$ . Achieve this condition by choosing an appropriately small grid spacing  $h$ :  

$$Kh^p = \epsilon/2.$$
- 2)  $\|e^h\| \leq \epsilon/2$ . Achieve this condition by iterating until  
 $\|e^h\| \leq \epsilon/2 = Kh^p$  on grid  $h$ ; then we've

converged to the level of discretization error.

→ Once discretization error & algebraic error are in balance, then it would be better to go to grid  $h/2$  than to iterate more!

## Convergence to the level of discretization error

- Use an MV scheme with convergence factor  $\gamma < 1$  bounded uniformly in  $h$  (fixed  $\alpha_1$  &  $\alpha_2$ ).
- Assume a  $d$ -dimensional problem on an  $n^d$  grid with  $h = 1/n$ .
- Initial relative error:  $\|e^h\|/\|u^h\| = \|u^h - 0\|/\|u^h\| = 1$ .
- Must reduce this to  $\|e^h\|/\|u^h\| = O(h^p) = O(n^{-p})$ .
- We can determine the number of V-cycles needed for this if we can bound the convergence factor,  $\gamma$ .

## Work to converge to the level of discretization error

- Using  $\theta$  V-cycles with convergence factor  $\gamma$  gives an overall convergence factor of  $\gamma^\theta$ .
- We therefore have  $\gamma^\theta = O(n^{-p})$ , or  $\theta = O(\log n)$ .
- Since 1 V-cycle costs  $O(1)$  WUs & 1 WUs is  $O(n^d)$ , then converging to the level of discretization error using the MV method cost  

$$O(n^d \log n).$$
- Compares to fast direct methods (fast Poisson solvers).

But multigrid can do even better...

## Numerical example

- Consider the 2-D model problem (with  $\sigma = 0$ ):  

$$-u_{xx} - u_{yy} = 2[(1 - 6x^2)y^2(1 - y^2) + (1 - 6y^2)x^2(1 - x^2)]$$
in the unit square, with  $u = 0$  Dirichlet boundary.
- The solution to this problem is  

$$u(x,y) = -(x^4 - x^2)(y^4 - y^2).$$
- We examine effectiveness of MV cycling to solve this problem on  $(n+1) \times (n+1)$  grids  $[(n-1) \times (n-1)$  interior points] for  $n = 16, 32, 64, 128$ .

V-cycle	n = 16				n = 32			
	$\ r^h\ _h$	ratio	$\ e\ _h$	ratio	$\ r^h\ _h$	ratio	$\ e\ _h$	ratio
0	6.75e+02		5.45e-01		2.60e+03		5.61e-01	
1	4.01e+00	0.01	1.05e-02	0.02	1.97e+01	0.01	1.38e-02	0.02
2	1.11e-01	0.03	4.10e-04	0.04	5.32e-01	0.03	6.32e-04	0.05
3	3.96e-03	0.04	1.05e-04	0.26	2.06e-02	0.04	4.41e-05	0.07
4	1.63e-04	0.04	1.03e-04	0.98*	9.79e-04	0.05	2.58e-05	0.59
5	7.45e-06	0.05	1.03e-04	1.00*	5.20e-05	0.05	2.58e-05	1.00*
6	3.75e-07	0.05	1.03e-04	1.00*	2.96e-06	0.06	2.58e-05	1.00*
7	2.08e-08	0.06	1.03e-04	1.00*	1.77e-07	0.06	2.58e-05	1.00*
8	1.24e-09	0.06	1.03e-04	1.00*	1.10e-08	0.06	2.58e-05	1.00*
9	7.74e-11	0.06	1.03e-04	1.00*	7.16e-10	0.06	2.58e-05	1.00*
10	4.99e-12	0.06	1.03e-04	1.00*	4.79e-11	0.07	2.58e-05	1.00*
11	3.27e-13	0.07	1.03e-04	1.00*	3.29e-12	0.07	2.58e-05	1.00*
12	2.18e-14	0.07	1.03e-04	1.00*	2.31e-13	0.07	2.58e-05	1.00*
13	2.33e-15	0.11	1.03e-04	1.00*	1.80e-14	0.06	2.58e-05	1.00*
14	1.04e-15	0.45	1.03e-04	1.00*	6.47e-15	0.36	2.58e-05	1.00*
15	6.61e-16	0.63	1.03e-04	1.00*	5.11e-15	0.79	2.58e-05	1.00*

V-cycle	n = 64				n = 128			
	$\ r^h\ _h$	ratio	$\ e\ _h$	ratio	$\ r^h\ _h$	ratio	$\ e\ _h$	ratio
0	1.06e+04		5.72e-01		4.16e+04		5.74e-01	
1	7.56e+01	0.01	1.39e-02	0.02	2.97e+02	0.01	1.39e-02	0.02
2	2.07e+00	0.03	6.87e-04	0.05	8.25e+00	0.03	6.92e-04	0.05
3	8.30e-02	0.04	4.21e-05	0.06	3.37e-01	0.04	4.22e-05	0.06
4	4.10e-03	0.05	7.05e-06	0.17	1.65e-02	0.05	3.28e-06	0.08
5	2.29e-04	0.06	6.44e-06	0.91*	8.99e-04	0.05	1.63e-06	0.50
6	1.39e-05	0.06	6.44e-06	1.00*	5.29e-05	0.06	1.61e-06	0.99*
7	8.92e-07	0.06	6.44e-06	1.00*	3.29e-06	0.06	1.61e-06	1.00*
8	5.97e-08	0.07	6.44e-06	1.00*	2.14e-07	0.06	1.61e-06	1.00*
9	4.10e-09	0.07	6.44e-06	1.00*	1.43e-08	0.07	1.61e-06	1.00*
10	2.87e-10	0.07	6.44e-06	1.00*	9.82e-10	0.07	1.61e-06	1.00*
11	2.04e-11	0.07	6.44e-06	1.00*	6.84e-11	0.07	1.61e-06	1.00*
12	1.46e-12	0.07	6.44e-06	1.00*	4.83e-12	0.07	1.61e-06	1.00*
13	1.08e-13	0.07	6.44e-06	1.00*	3.64e-13	0.06	1.61e-06	1.00*
14	2.60e-14	0.24	6.44e-06	1.00*	1.03e-13	0.28	1.61e-06	1.00*
15	2.30e-14	0.88	6.44e-06	1.00*	9.19e-14	0.89	1.61e-06	1.00*

CU-Boulder

161 of 396

## Numerical results norms

MV cycling  
 shown across grid levels of 15 V(2,1) cycles. We display, after each cycle, about  $n^2$  grid points, residual norms, total error norms, & ratios of these norms to their values after the previous cycle.

$n = 16, 32, 64, 128$ .  
 $\|r^h\|_h \approx 1$

$$\|r^h\|_h = h \|r^h\|_2$$

scaled residual error

$$\|e\|_h = h \|u^{(h)} - v^h\|_2$$

scaled discrete total error

## A warning about bounds

- Bounds like  $\|e_{n+1}\| \leq \gamma \|e_n\|$  &  $\|u^{(h)} - u^h\| = O(h)$  are only just that--bounds!
- If you see behavior that suggests that these bounds are sharp (e.g., halving  $h$  halves the discretization error), then great. If you don't see this behavior, don't assume things are wrong.
- Think about this:  
 $O(h^2) = O(h)$  but generally  $O(h) \neq O(h^2)$  !!!  
 (Any process that is  $O(h^2)$  is also  $O(h)$ , but the converse isn't necessarily true.)

CU-Boulder

162 of 396

## Reconsideration

You want to approximate  $u^h$ .

A good iteration is the V-cycle.

What's a good way to start it?

Can you do better than  $v^h \leftarrow 0$ ?

→ Start on the coarse grid. ←

Use nested iteration for the V-cycle.

## Look again at nested iteration

- Idea: It's cheaper to solve a problem (fewer iterations) if the initial guess is good.
- How to get a good initial guess:
  - "Solve" the problem on the coarse grid first.
  - Interpolate the coarse solution to the fine grid.
- Now, let's use the V-cycle as the solver on each grid level! This defines the Full Multigrid (FMG) cycle.

CU-Boulder

163 of 396

CU-Boulder

164 of 396

## Full multigrid (FMG)

$$v^h \leftarrow \text{FMG}(f^h)$$

• Initialize  $f^h, f^{2h}, f^{4h}, \dots, f^H$

• Solve on coarsest grid

$$v^H = (A^H)^{-1} f^H$$

• Interpolate initial guess

$$v^{2h} \leftarrow I_{4h}^{2h} v^{4h}$$

• Perform V-cycle

$$v^{2h} \leftarrow MV^{2h}(v^{2h}, f^{2h})$$

• Interpolate initial guess

$$v^h \leftarrow I_{2h}^h v^{2h}$$

• Perform V-cycle

$$v^h \leftarrow MV^h(v^h, f^h)$$

## FMG-cycle (recursive form)

$$v^h \leftarrow \text{FMG}(f^h), \eta$$

1) Initialize  $f^h, f^{2h}, \dots, f^H$ .

2) If  $h = H$ , then go to 4 (where  $MV$  is a direct solve);

else:  $v^{2h} \leftarrow \text{FMG}(f^{2h})$ .

3) Set initial guess:  $v^h \leftarrow I_{2h}^h v^{2h}$ .

We use  
 $\eta = 1$ .

4) Perform  $v^h \leftarrow MV(v^h, f^h)$   $\eta$  times.

## FMG cycle cost

One V(2,1)-cycle is performed per level, at a cost of  $3/(1 - 2^{-d})$  WUs per grid (where the WU is for the size of the finest grid involved).

The size for the WU for coarse-grid  $j$  is  $2^{-jd}$  times the size for the WU for the fine grid (grid 0).

Hence, the cost of the FMG(2,1) cycle in WUs is less than

$$[3/(1 - 2^{-d})](1 + 2^{-d} + 2^{-2d} + \dots) = 3/(1 - 2^{-d})^2.$$

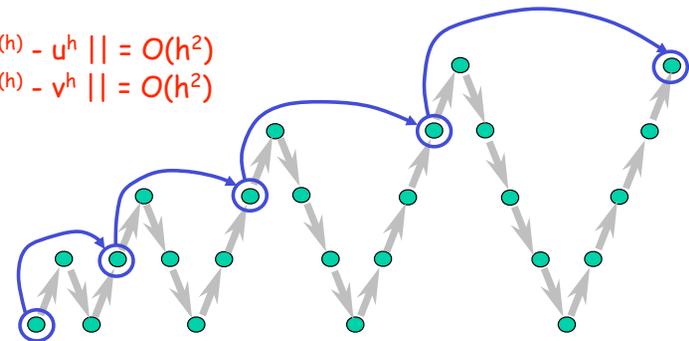
$d = 1$ : 12 WUs;  $d = 2$ : 16/3 WUs;  $d = 3$ : 192/49 WUs.

## Has discretization error been reached by FMG?

If discretization error is achieved, then  $\|e^h\| = O(h^2)$  & the V-cycle approximation converges to the solution of the PDE about as well as the discrete solution does:

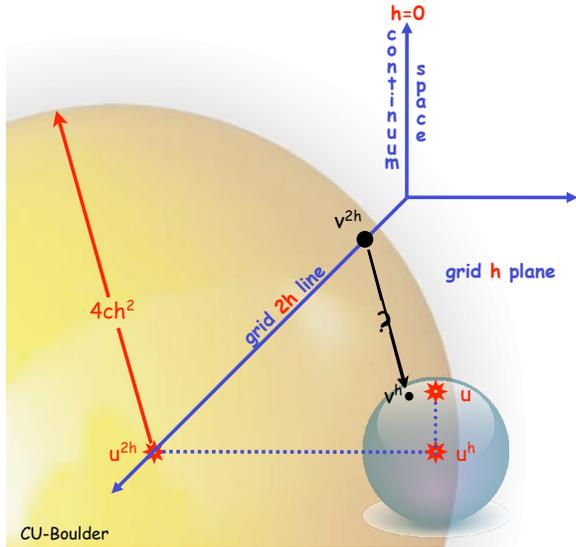
$$\|u^{(h)} - u^h\| = O(h^2)$$

$$\|u^{(h)} - v^h\| = O(h^2)$$



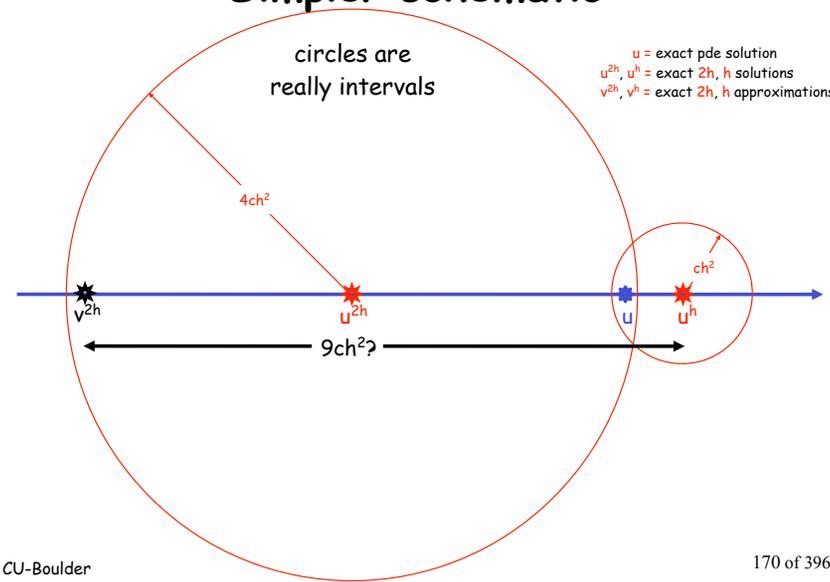
We need to be more careful...

# The basic FMG principle



$u$  = exact pde solution  
 $u^{2h}, u^h$  = exact  $2h, h$  solutions  
 $v^{2h}, v^h$  = exact  $2h, h$  approximations

# Simpler schematic



## Comparing the right things

- **Problem:** We are thinking that  $u^{2h}$  approximates  $u^h$  to order  $O(h^2)$ , when all we really know is that  $u^h$  approximates  $u^{(h)}$  to order  $O(h^2)$  (any  $h$ ).
- We know that  $u^{(2h)}$  &  $u^{(h)}$  are the "same", right? So, if  $u^{2h}$  approximates  $u^{(2h)}$  to order  $O(4h^2)$  &  $u^h$  approximates  $u^{(h)}$  to order  $O(h^2)$ , shouldn't  $u^{2h}$  approximate  $u^h$  to order  $O(4h^2)$ ? How, exactly?
- When we interpolate  $u^{2h}$  to grid  $h$ , what errors does interpolation introduce?
- Sorting out these comparisons is a bit technical.
- In other words, here comes the algebra...

## Interpolation stability

how interpolation affects error

• Property:  $\|Pe^{2h}\| \leq \beta \|e^{2h}\|$

• Reasoning:

$$\|Pe^{2h}\| \leq \|P\| \|e^{2h}\| = \|P^T P\|^{1/2} \|e^{2h}\|$$

$$\begin{aligned} \|P\|^2 &= \max \|Pe\|^2 / \|e\|^2 \\ &= \max \langle Pe, Pe \rangle / \langle e, e \rangle \\ &= \max \langle P^T P e, e \rangle / \langle e, e \rangle \\ &= \rho(P^T P) \\ &= \|P^T P\| \end{aligned}$$

$$P^T P = \begin{pmatrix} 1/2 & 1 & 1/2 & & & \\ & 1/2 & 1 & 1/2 & & \\ & & 1/2 & 1 & 1/2 & \\ & & & & & & \\ & & & & & & \\ & & & & & & & & & & \end{pmatrix} \begin{pmatrix} 1/2 & & & & & \\ 1 & 1/2 & & & & \\ & 1 & & & & \\ & & 1/2 & 1/2 & & \\ & & & 1 & & \\ & & & & 1/2 & \end{pmatrix} = \begin{pmatrix} 3/2 & 1/4 & & & & \\ 1/4 & 3/2 & 1/4 & & & \\ & 1/4 & 3/2 & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & & & & & \end{pmatrix}$$

$$\Rightarrow \|P^T P\|^{1/2} \leq \sqrt{2}$$

In practice,  $\beta \approx 1$ .

## Approximation property

how the discrete solutions approximate each other

$$\|u^h - Pu^{2h}\| \leq \alpha Kh^2$$

$$|(u^{(h)} - Pu^{(2h)})_i| = |u(x_i) - (u(x_{i-1}) + u(x_{i+1}))/2| \leq |u''(\xi_i)|h^2/2$$

$$\Rightarrow \|u^{(h)} - Pu^{(2h)}\| \leq cKh^2 \quad (\|u^h - u^{(h)}\| \approx Kh^2)$$

(We need  $\|u''\| \ll \infty$ , so the norm is scaled by  $h$  here.)

$$\begin{aligned} \Rightarrow \|u^h - Pu^{2h}\| &\leq \|u^h - u^{(h)}\| + \|u^{(h)} - Pu^{(2h)}\| + \|Pu^{(2h)} - Pu^{2h}\| \\ &\leq Kh^2 + cKh^2 + \|P\| \cdot \|u^{(2h)} - u^{2h}\| \\ &\leq Kh^2 + cKh^2 + \beta K(2h)^2 \\ &\leq (1 + c + 4\beta)Kh^2 \end{aligned}$$

In practice,  $\alpha = 1 + c + 4\beta \approx 5$ .

## FMG accuracy

$$\|e^h\| \leq Kh^2$$

Assume:

$$\|e^{2h}\| \leq K(2h)^2$$

induction hypothesis

$$\|u^h - Pu^{2h}\| \leq \alpha Kh^2$$

approximation property ( $\alpha \approx 5$ )

$$\|Pw^{2h}\| \leq \beta \|w^{2h}\|$$

interpolation stability ( $\beta \approx 1$ )

Triangle inequality:

$$\begin{aligned} \|e^h\| &= \|u^h - Pv^{2h}\| \\ &\leq \|u^h - Pu^{2h}\| + \|P(u^{2h} - v^{2h})\| \\ &\leq \alpha Kh^2 + \beta K(2h)^2 \\ &= (\alpha + 4\beta)Kh^2 \end{aligned}$$

$\Rightarrow$

$$\|e^h\| \leq "9"Kh^2$$

So we need only reduce  $\|e^h\|$  by "0.1"!!!

## Numerical example

- Consider again the 2-D model problem (with  $\sigma = 0$ ):  
 $-u_{xx} - u_{yy} = 2[(1 - 6x^2)y^2(1 - y^2) + (1 - 6y^2)x^2(1 - x^2)]$   
 inside the unit square, with  $u = 0$  on the boundary.
- We examine the effectiveness of FMG cycling to solve the problem on  $(n+1) \times (n+1)$  grids [ $(n-1) \times (n-1)$  interior points] for  $n = 2, 4, \dots, 2048$ .

## FMG results

FMG cycle results & comparison with MV cycle costs

N	FMG(1,0)		FMG(1,1)		FMG(2,1)		FMG(1,1)	V(2,1)	V(2,1)
	$\ e\ _h$	ratio	$\ e\ _h$	ratio	$\ e\ _h$	ratio	WU	cycles	WU
2	5.86e-03		5.86e-03		5.86e-03				
4	5.37e-03	0.917	2.49e-03	0.424	2.03e-03	0.347	7/2	3	12
8	2.78e-03	0.518	9.12e-04	0.367	6.68e-04	0.328	7/2	4	16
16	1.19e-03	0.427	2.52e-04	0.277	1.72e-04	0.257	1.03e-04	4	16
32	4.70e-04	0.395	6.00e-05	0.238	4.00e-05	0.233	2.58e-05	5	20
64	1.77e-04	0.377	1.36e-05	0.227	9.36e-06	0.234	6.44e-06	5	20
128	6.49e-05	0.366	3.12e-06	0.229	2.26e-06	0.241	1.61e-06	6	24
256	2.33e-05	0.359	7.35e-07	0.235	5.56e-07	0.246	7/2	7	28
512	8.26e-06	0.354	1.77e-07	0.241	1.38e-07	0.248	7/2	7	28
1024	2.90e-06	0.352	4.35e-08	0.245	3.44e-08	0.249	7/2	8	32
2048	1.02e-06	0.351	1.08e-08	0.247	8.59e-09	0.250	7/2	9	36

$$\|e\|_h = h \|u^{(h)} - v^h\|_2$$

scaled discrete total error

## Successful Scientific Inquiry

- **Attitude**
  - Knowledge is good, but understanding rules!
  - Look for the underlying principle!
  - You can do it! Be positive.
  - But is it really right? Be critical.
  - Don't hope or guess. Think!
  - Control your emotions! Expect ups & downs.
- **Method**
  - Start simply. Reduce issue to the simplest possible case.
  - Take tiny steps, but keep the big picture in mind.
  - Study concrete examples.
  - Look for analogies. Can A be done in any way like how B was done?
- **Creativity**
  - What do you really want? What end are you really aiming for?
  - What do you really need? What you're trying may be sufficient to do what you want, but would an easier weak result do instead?
- **Intelligence**
  - It doesn't hurt to try to be "smart" too.

CU-Boulder

177 of 396

## Diagnostic tools

for debugging the code, the method, the problem

- Finding mistakes in codes, algorithms, concepts, & the problem itself challenges our scientific abilities.
- This challenge is especially tough for multigrid:
  - Interactions between multilevel processes can be very subtle.
  - It's often not easy to know how well multigrid should perform.
- **Achi Brandt:**
  - "The amount of computational work should be proportional to the amount of real physical changes in the computed solution."
  - "Stalling numerical processes must be wrong."
- The "computational culture" is best learned by lots of experience & interaction, but some discussion helps.

CU-Boulder

178 of 396

## Tool # 1: Be methodical

- Modularize your code.
- Test the algebraic solver first.
- Test the discretization next.
- Test the FMG solver last.
- Beware of boundaries, scales, & concepts.
- Ask whether the problem itself is well posed.

CU-Boulder

179 of 396

## Tool # 2: Start simply

- Start from something that already works if you can.
- Introduce complexities slowly & methodically, testing thoroughly along the way.
- Start with a very coarse fine grid (no oxymoron intended).
- Start with two levels if you can, using a direct solver or lots of cycles on coarse grids if nothing else.

If you find trouble, your first job is to find the simplest case where that trouble is still evident!!!

CU-Boulder

180 of 396

## Tool # 3: Expose trouble

Start simply, but don't let niceties mask trouble:

- Set reaction/Helmholtz terms to zero.
- Take infinite or very big time steps.
- Don't take 1-D too seriously, not even 2-D.

## Tool # 4: Test fixed point property

Relaxation shouldn't alter the exact solution of the linear system (up to machine precision).

- Create a right side:  $f^h = A^h u^h$  with  $u^h$  given.
- Make sure  $u^h$  satisfies the right boundary conditions.
- Test relaxation starting with  $u^h$ : Is  $r^h = 0$ , is it zero after relaxation, does  $u^h$  change?
- Test coarse-grid correction starting with  $u^h$ : Is the correction zero?

## Tool # 5: Test on $Au^h = 0$

- The exact solution  $u^h = 0$  is known!
- Residual norm  $\|Av^h\|$  & error norm  $\|v^h\|$  are computable.
- Norms  $\|Av^h\|$  &  $\|v^h\|$  should eventually decrease steadily with a rate that might be predicted by mode analysis.
- Multigrid can converge so fast that early stalling suggests trouble when it's just that all machine-representable numbers in a nonzero  $v^h$  have already been computed! Computing  $r^h = f^h - Av^h$  & updating  $v^h$  shouldn't have trouble with machine precision if you have  $u^h = 0$  & thus  $f^h = 0$ .

## Tool # 6: Zero out residual

- Using a normal test, try multiplying the residual by 0 before you go to the coarse grid.
- Check to see that the coarse-grid corrections are 0.
- Compare this test with a relaxation-only test--the results should be identical.

## Tool # 7: Print out residual norms

dropping superscript  $h$  when it's clear by context

- Use the discrete  $L^2$  norm:

$$\|r\|_h = (h^d \sum r_i^2)^{1/2} = h^{d/2} \|r\|_2.$$

- Output  $\|r\|_h$  after each pre- & post-relaxation sweep.
- These norms should decline to zero steadily for each  $h$ .
- The norm after post-relaxation should be consistently smaller than after pre-relaxation--by the predictable convergence factor at least.

## Tool # 8: Graph the error

- Run a test on a problem with known solution ( $Au = 0$ ).
- Plot algebraic error before & after fine-grid relaxation.
- Is the error oscillatory after coarse-grid correction?
- Is the error much smoother after fine-grid relaxation?
- Are there any strange characteristics near boundaries, interfaces, or other special phenomena?

example  $h^{-2} (-1 \ 2 \ -1)$  **Beware of residuals**  $e_k = \sin(k\pi x)$   
 $Ae = r$   $\lambda_1 \approx \pi^2, \lambda_n \approx 4h^{-2}$

- Relative errors:  $\frac{\|e\|_h}{\|u\|_h}$  vs.  $\frac{\|Ae\|_h}{\|Au\|_h} = \frac{\|r\|_h}{\|f\|_h}$
- Absolute range:  $\|e_1\|_h \approx 1$  &  $\|Ae_1\|_h \approx \pi^2$   
 $\|e_n\|_h \approx 1$  &  $\|Ae_n\|_h \approx 4h^{-2}$  !!!
- Relative errors: consider the case  $u = e_n$   
 $\frac{\|e_n\|_h}{\|u\|_h} = 1$  &  $\frac{\|Ae_n\|_h}{\|Au\|_h} = 1$   
 $\frac{\|e_1\|_h}{\|u\|_h} = 1$  &  $\frac{\|Ae_1\|_h}{\|Au\|_h} \approx (\pi^2/4)h^2$  !!!

Moral: residuals can falsely signal convergence  
 when the error is smooth.

## Tool # 9: Test two-level cycling

- Replace the coarse-grid V-cycle recursive call with a direct solver if possible, or iterate many times with some method known to "work" (test  $\|r\|$  to be sure it's very small), or use many recursive V-cycle calls.
- This can be used to test performance between two coarser levels, especially if residual norm behavior identifies trouble on a particular level.

## Tool # 10: Beware of boundaries

- Boundaries usually require special treatment of the stencils, intergrid transfers, & sometimes relaxation.
- Special treatment often means special trouble, typically exposed in later cycles as it begins to infect the interior.
- Replace the boundary by periodic or Dirichlet conditions.
- Relax more at the boundary, perhaps using direct solvers.
- Make sure your coarse-grid approximation at the boundary is guided by good discretization at the fine-grid boundary.

## Tool # 11: Test for symmetry

- If your problem is symmetric or includes a symmetric case, test for it.
- Check symmetry of the fine-grid & coarse-grid matrices: are  $a_{ij}$  &  $a_{ji}$  relatively equal (to machine precision).
- Be especially watchful for asymmetries near boundaries.

## Tool # 12: Check for compatibility

a bit ahead of schedule, but...

- Consider the problem
$$-u'' = f \quad \text{with} \quad u'(0) = u'(1) = 0.$$
- It's singular: If  $u = 1$ , then  $-u'' = 0$  &  $u'(0) = u'(1) = 0$ .
- It's solvable iff  $f \in \text{Range}(\partial_{xx}) = \eta^\perp(\partial_{xx}) = \{1\}^\perp$  or  $f \perp 1$ .
- First fix the grid  $h$  right side:  $f^h \leftarrow f^h - \langle f^h, 1 \rangle / \langle 1, 1 \rangle 1$ .
- Do this on coarse grids too:  $f^{2h} \leftarrow f^{2h} - \langle f^{2h}, 1 \rangle / \langle 1, 1 \rangle 1$ .
- Uniqueness is also a worry:  $u^h \leftarrow u^h - \langle u^h, 1 \rangle / \langle 1, 1 \rangle 1$ .

## Tool # 13: Test for linearity

also a bit ahead of schedule...

- If you're writing a nonlinear FAS code, it should agree with the linear code when you test it on a linear problem. Try it.
- Move gradually to the target nonlinear test problem by putting a parameter in front of the nonlinear term, then running tests as the parameter changes slowly from 0 to 1.

## Tool # 14: Use a known PDE solution

- Set up the source term ( $f = -u''$  in  $\Omega$ ) & data ( $g = u$  on  $\Gamma$ ).
- Do multigrid results compare qualitatively with sampled  $u$ ?
- Monitor  $\|u - u^h\|_h$ .
- Test a case with no discretization error ( $u = ax^2 + bx + c$ ). The algebraic error should tend steadily to 0.
- Test discretization error ( $u^{iv} \neq 0$ ). The algebraic error should decrease rapidly at first, then stall at discretization error level. Check error behavior as you decrease  $h$ . Does it behave like  $O(h^2)$  ( $h \rightarrow h/2 \Rightarrow e \rightarrow e/2$ ) or however it should?

## Tool # 15: Test FMG accuracy

- Make sure first that the algebraic solver converges as predicted, with uniformly bounded convergence factors.
- Test the discretization using Tool # 14.
- Compare FMG total error to discretization error for various  $h$ . You might need to tune the FMG process here (play with the number of cycles & relaxation sweeps).

## Computing assignments

- Document: norms/weights,  $V(v_1, v_2)$ , errors, labels (table, graph)
- Use various scenarios:  $Ax = 0$ ,  $Ax = f$ , varying  $n$  &  $v_i$  &  $\omega$ , Jacobi/Gauss-Seidel
- Thoroughly test: don't stop until you get what you expect. compare with known solution, text, others. study discretization & algebraic errors. report on "asymptotic" factors.
- Be kind to the reader:   

code = zzz...	tables = +	tables & graphs = ++
		by hand ok
		↓            ↓
		↓            ↓
- Discuss, discuss, discuss: what do you see & think? what did you learn?

## Outline

### Chapters 1-5:

- ✓ Model Problems
- ✓ Basic Iterative Methods
  - Convergence tests
  - Analysis
- ✓ Elements of Multigrid
  - Relaxation
  - Coarsening
- ✓ Implementation
  - Complexity
  - Diagnostics
- Some Theory
  - Spectral vs. algebraic

### Chapters 6-10:

- Nonlinear Problems
  - Full approximation scheme
- Selected Applications
  - Neumann boundaries
  - Anisotropic problems
  - Variable meshes
  - Variable coefficients
- Algebraic Multigrid (AMG)
  - Matrix coarsening
- Multilevel Adaptive Methods
  - FAC
- Finite Elements
  - Variational methodology

Homework Due!

## 5. Some theory

### What is $A^{2h}$ ?

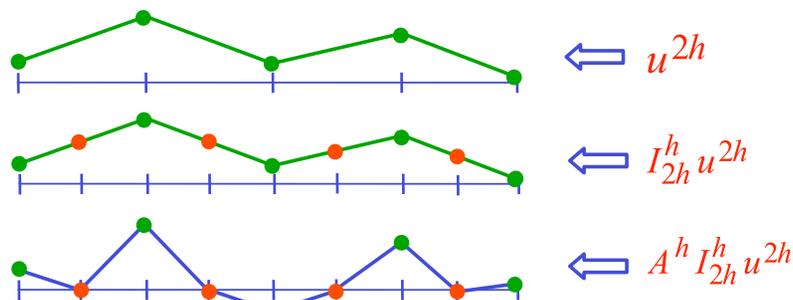
- Recall the 2-grid coarse-grid correction scheme:
  - 1) Relax on  $A^h u^h = f^h$  on  $\Omega^h$  to get  $v^h$ .
  - 2) Compute  $f^{2h} = I_{2h}^{2h}(f^h - A^h v^h)$ .
  - 4) Solve  $A^{2h} u^{2h} = f^{2h}$  on  $\Omega^{2h}$ .
  - 5) Correct fine-grid solution  $v^h \leftarrow v^h + I_{2h}^h u^{2h}$ .
- Assume that  $e^h \in \text{Range}(I_{2h}^h)$ , i. e.,  $e^h = I_{2h}^h u^{2h}$  for some  $u^{2h} \in \Omega^{2h}$ . Then the residual equation can be written

$$A^h e^h = A^h I_{2h}^h u^{2h} = r^h.$$

This characterizes  $u^{2h}$ , but with too many equations.

- How does  $A^h$  act on  $I_{2h}^h$ ?

### How does $A^h$ act on $\text{Range}(I_{2h}^h)$ ?



Thus, the odd rows of  $A^h I_{2h}^h$  are zero (1-D only) &  $r_{2i+1} = 0$ . So we keep the even rows of  $A^h I_{2h}^h$  for the residual equations on  $\Omega^{2h}$ . We do this by applying restriction, either injection or full weighting:

$$I_h^{2h} A^h I_{2h}^h u^{2h} = I_h^{2h} r^h.$$

We use full weighting from now on unless otherwise stated.

### Building $A^{2h}$ : The Galerkin condition

- The residual equation on the coarse grid is

$$I_h^{2h} A^h I_{2h}^h u^{2h} = I_h^{2h} r^h$$

- We thus identify the coarse-grid operator as

$$A^{2h} = I_h^{2h} A^h I_{2h}^h \quad A^{2h} = R A^h P$$

- RAP** is symmetric:

If  $P^T = \alpha R$  (so that  $R^T = (1/\alpha)P$ ), then  
 $(RAP)^T = P^T A^T R^T = \alpha(1/\alpha)RAP = RAP$ .

- RAP** is positive definite:  $P$  full rank  $\Rightarrow Px \neq 0$ .

If  $x \neq 0$ , then  $\langle RAPx, x \rangle = (1/\alpha) \langle APx, Px \rangle > 0!$

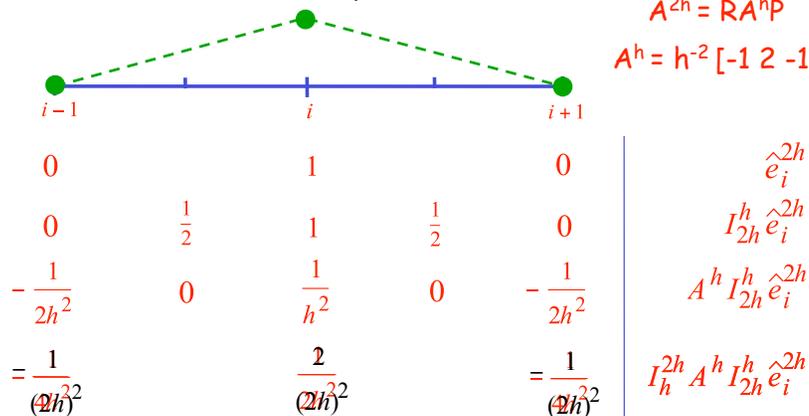
### Computing the $i^{\text{th}}$ row of $A^{2h}$

- Compute  $A^{2h} \hat{e}_i^{2h}$ , where  $\hat{e}_i^{2h} = (0, 0, \dots, 0, 1, 0, \dots, 0)^T$ .

↑ Why is this the  $i^{\text{th}}$  row of  $A^{2h}$ ?

$$A^{2h} = R A^h P$$

$$A^h = h^{-2} [-1 \ 2 \ -1]$$



The  $i^{\text{th}}$  row of  $A^{2h}$  looks a lot like the  $i^{\text{th}}$  row of  $A^h$ !

- The  $i^{\text{th}}$  row of  $A^{2h}$  is  $\frac{1}{(2h)^2} [-1 \ 2 \ -1]$ ,

which is the  $\Omega^{2h}$  version of  $A^h$ .

- Note that **IF** relaxation on  $\Omega^h$  leaves only error in the range of interpolation, then solving

$$A^{2h} u^{2h} = f^{2h}$$

determines the error exactly!

- This is generally not feasible, but this logic motivates wanting  $e^h \in \text{Range}(I_{2h}^h)$  & it leads to a very plausible representation for  $A^{2h}$ .

## 2-D RAP

scale  $\downarrow$

$$\begin{bmatrix} \frac{11}{164} & \frac{11}{82} & \frac{1}{16} \\ \frac{11}{82} & \frac{1}{4} & \frac{1}{8} \\ \frac{1}{164} & \frac{1}{82} & \frac{1}{16} \end{bmatrix} \frac{1}{h^2} \begin{bmatrix} -1 & -1 & -1 \\ -1 & 4 & -1 \\ -1 & -1 & -1 \end{bmatrix} \begin{bmatrix} \frac{1}{4} & \frac{1}{2} & \frac{1}{4} \\ \frac{1}{2} & 1 & \frac{1}{2} \\ \frac{1}{4} & \frac{1}{2} & \frac{1}{4} \end{bmatrix} e_i^{2h}$$

Try  $\frac{1}{h^2} \begin{bmatrix} -1/3 & -1/3 & -1/3 \\ -1/3 & 8/3 & -1/3 \\ -1/3 & -1/3 & -1/3 \end{bmatrix}$  homework!

$\frac{11}{(2h)^2} \otimes$

-1/4	-1/4	-1/2	-1/4	-1/4
-1/4	0	1/2	0	-1/4
-1/2	1/2	3	1/2	-1/2
-1/4	0	1/2	0	-1/4
-1/4	-1/4	-1/2	-1/4	-1/4

## Variational properties of coarsening

The definition for  $A^{2h}$  that resulted from the foregoing line of reasoning is useful for both theoretical & practical reasons. Together with the commonly used relationship between restriction & prolongation, we have the variational properties:

$$A^{2h} = I_h^{2h} A^h I_{2h}^h$$

$$I_{2h}^h = c (I_h^{2h})^T$$

Galerkin Condition

$c \in \mathfrak{R}$

## Properties of restriction

in a little more detail...

- Full Weighting:  $I_h^{2h} : \Omega^h \rightarrow \Omega^{2h}$  or  $I_h^{2h} : \mathfrak{R}^{n-1} \rightarrow \mathfrak{R}^{n/2-1}$

- $n = 8$ : 
$$I_h^{2h} = \frac{1}{4} \begin{bmatrix} 1 & 2 & 1 & & & & & \\ & & & 1 & 2 & 1 & & \\ & & & & & & 1 & 2 & 1 \\ & & & & & & & & & 1 & 2 & 1 \end{bmatrix}_{3 \times 7}$$

- $I_h^{2h}$  has rank  $\frac{n}{2} - 1$  & null space  $\eta(I_h^{2h})$  with  $\dim \frac{n}{2}$ .

## Spectral properties of restriction

- How does  $I_h^{2h}$  act on the eigenvectors of  $A^h$ ?
- Consider  $w_{k,j}^h = \sin\left(\frac{jk\pi}{n}\right)$ ,  $1 \leq k \leq n-1$ ,  $0 \leq j \leq n-1$ .
- A little algebra & trigonometry shows that

$$(I_h^{2h} w_k^h)_j = \cos^2\left(\frac{k\pi}{2n}\right) w_{k,j}^{2h}$$

$$\equiv c_k w_{k,j}^{2h}$$

for  $1 \leq k \leq n/2-1$ .

$$c_k = O(1)$$

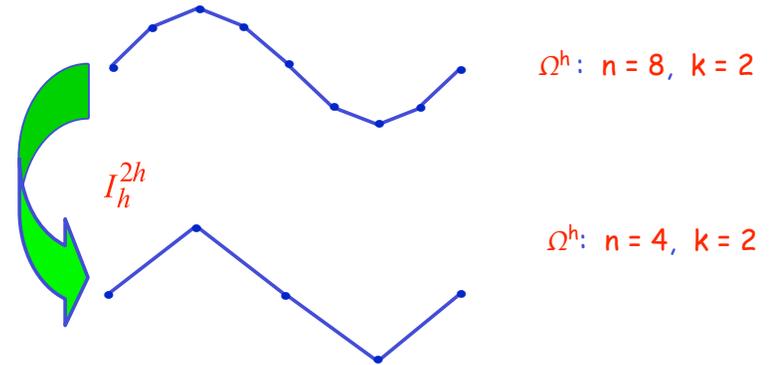
$$c_1 \approx 1 \dots c_n = O(h^2)$$

CU-Boulder

205 of 396

## Spectral properties (cont'd)

- i.e.,  $I_h^{2h}$  [ $k^{\text{th}}$  mode on  $\Omega^h$ ] =  $c_k$  [ $k^{\text{th}}$  mode on  $\Omega^{2h}$ ]



CU-Boulder

206 of 396

## Spectral properties (cont'd)

- Let  $k' = n - k$  for  $1 \leq k \leq n/2-1$ , so that  $n/2+1 \leq k' \leq n-1$ .
- A little algebra & trigonometry shows that

$$(I_h^{2h} w_{k'}^h)_j = -\sin^2\left(\frac{k\pi}{2n}\right) w_{k',j}^{2h}$$

$$\equiv -s_k w_{k',j}^{2h}$$

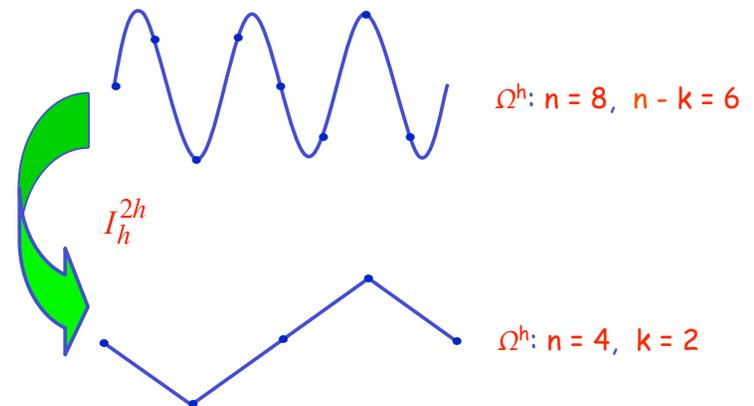
$$s_1 = O(h^2) \dots s_{n/2-1} = O(1)$$

CU-Boulder

207 of 396

## Spectral properties (cont'd)

- i.e.,  $I_h^{2h}$  [ $(n-k)^{\text{th}}$  mode on  $\Omega^h$ ] =  $-s_k$  [ $k^{\text{th}}$  mode on  $\Omega^{2h}$ ]



CU-Boulder

208 of 396

## Spectral properties (cont'd)

- Summarizing:
 
$$\left. \begin{aligned} I_h^{2h} w_k^h &= c_k w_k^{2h} \\ I_h^{2h} w_{k'}^h &= -s_k w_{k'}^{2h} \end{aligned} \right\} \begin{cases} 1 \leq k \leq \frac{n}{2}-1 \\ k' = n-k \end{cases}$$

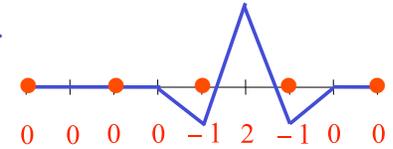
$c_k = O(1)$  (pointing to  $c_k$ )  
 $s_1 = O(h^2) \dots s_{n/2-1} \approx O(1)$  (pointing to  $s_k$ )
- Complementary modes:

$$W_k = \text{span} \{ w_k^h, w_{k'}^h \}$$

$$I_h^{2h} W_k \implies \{ w_k^{2h} \}$$

## Null space of restriction

- Observe that  $\eta(I_h^{2h}) = \text{span}(A^h \hat{e}_i^h)$ , where  $i$  is odd &  $\hat{e}_i^h$  is the  $i^{\text{th}}$  unit vector.



- Let  $\eta_i = A^h \hat{e}_i^h$ .
- While the  $\eta_i$  looks oscillatory, it generally contains all Fourier modes of  $A^h$ :

$$\eta_i = \sum_{k=1}^N a_k w_k \quad a_k \neq 0$$

- All the Fourier modes of  $A^h$  are needed to represent the null space of restriction!

## Properties of interpolation

- Interpolation:  $I_{2h}^h : \Omega^{2h} \implies \Omega^h$  or  $I_{2h}^h : \mathfrak{R}^{n/2-1} \implies \mathfrak{R}^{n-1}$

$n = 8$ :

$$I_{2h}^h = \frac{1}{2} \begin{bmatrix} 1 & & & & & & & \\ 2 & & & & & & & \\ 1 & 1 & & & & & & \\ & 2 & & & & & & \\ & 1 & 1 & & & & & \\ & & & 2 & & & & \\ & & & 1 & 1 & & & \\ & & & & & 2 & & \\ & & & & & & 1 & \end{bmatrix}$$

- $I_{2h}^h$  has full rank & null space  $\{0\}$ .

## Spectral properties of interpolation

- How does  $I_{2h}^h$  act on the eigenvectors of  $A^{2h}$ ?
- Consider  $(w_k^{2h})_j = \sin\left(\frac{jk\pi}{n/2}\right)$ ,  $1 \leq k \leq n/2-1$ ,  $0 \leq j \leq n/2$ .
- A bit of work shows that the modes of  $A^{2h}$  are NOT "preserved" by  $I_{2h}^h$ , but that the space  $W_k$  is "preserved":

$$\begin{aligned} I_{2h}^h w_k^{2h} &= \cos^2\left(\frac{k\pi}{2n}\right) w_k^h - \sin^2\left(\frac{k\pi}{2n}\right) w_{k'}^h \\ &= c_k w_k^h - s_k w_{k'}^h \end{aligned}$$

# Spectral properties of interpolation

$$I_{2h}^h w_k^{2h} = c_k w_k^h - s_k w_{k'}^h$$

- Interpolation of smooth  $\Omega^{2h}$  modes excites oscillatory modes on  $\Omega^h$ .

- Note that if  $k \ll \frac{n}{2}$ , then

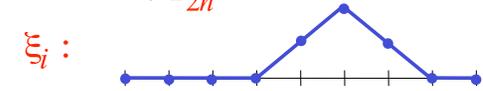
$$I_{2h}^h w_k^{2h} = \left(1 - O\left(\frac{k^2}{(n-1)^2}\right)\right) w_k^h + O\left(\frac{k^2}{(n-1)^2}\right) w_{k'}^h \approx w_k^h$$

- $I_{2h}^h$  is 2<sup>nd</sup>-order interpolation.

# Range of interpolation

- The range of  $I_{2h}^h$  is the span of the columns of  $I_{2h}^h$ .

- Let  $\xi_i$  be the  $i^{\text{th}}$  column of  $I_{2h}^h$ .



$$\xi_i^h = \sum_{k=1}^{n-1} b_k w_k^h, \quad b_k \neq 0$$

- All the Fourier modes of  $A^h$  are needed to represent  $\text{Range}(I_{2h}^h)$ .

# Use all the facts to analyze the coarse-grid correction scheme

relax only before correction

- 1) Relax once on  $\Omega^h$ :  $v^h \leftarrow R_\omega v^h + B f^h$
- 2) Compute & restrict residual:  $f^{2h} \leftarrow I_h^{2h} (f^h - A^h v^h)$
- 3) Solve residual equation:  $v^{2h} = (A^{2h})^{-1} f^{2h}$
- 4) Correct fine-grid solution:  $v^h \leftarrow v^h + I_{2h}^h v^{2h}$

error propagation matrix not restriction!

The entire process appears as

$$v^h \leftarrow R_\omega v^h + B f^h + I_{2h}^h (A^{2h})^{-1} I_h^{2h} (f^h - A^h (R_\omega v^h + B f^h))$$

The exact solution satisfies

$$u^h = R_\omega u^h + B f^h + I_{2h}^h (A^{2h})^{-1} I_h^{2h} (f^h - A^h (R_\omega u^h + B f^h))$$

# CG error propagation

- Subtracting the previous two expressions, we get

$$e^h \leftarrow \left[ I - I_{2h}^h (A^{2h})^{-1} I_h^{2h} A^h \right] R_\omega e^h$$

$$e^h \leftarrow CG e^h$$

- How does CG act on the modes of  $A^h$ ? Assume  $e^h$  consists of the modes  $w_k^h$  &  $w_{k'}^h$  for  $1 \leq k \leq \frac{n}{2} - 1$  &  $k' = n - k$ .

- We know how  $R_\omega^\alpha$ ,  $A^h$ ,  $I_h^{2h}$ ,  $(A^{2h})^{-1}$ ,  $I_{2h}^h$  act on  $w_k^h$  &  $w_{k'}^h$ .

## CG error propagation

For now, assume no relaxation. Then

$$W_k = \text{span} \{w_k^h, w_{k'}^h\}$$

is invariant under CG:

$$CG w_k^h = s_k w_k^h + s_k w_{k'}^h$$

$$CG w_{k'}^h = c_k w_k^h + c_k w_{k'}^h$$

where

$$s_k = \sin^2\left(\frac{k\pi}{2n}\right) \quad c_k = \cos^2\left(\frac{k\pi}{2n}\right)$$

$$s_1 = O(h^2) \dots s_{n/2-1} \approx O(1) \quad c_k = O(1)$$

## CG error propagation for $k \ll n$

- Consider the case  $k \ll n$  (extremely smooth & oscillatory modes):

$$w_k \rightarrow O\left(\frac{k^2}{(n-1)^2}\right) w_k + O\left(\frac{k^2}{(n-1)^2}\right) w_{k'}$$

$$w_{k'} \rightarrow \left(1 - O\left(\frac{k^2}{(n-1)^2}\right)\right) w_k + \left(1 - O\left(\frac{k^2}{(n-1)^2}\right)\right) w_{k'}$$

- Hence, CG eliminates the smooth modes but does not damp the oscillatory modes of the error!

## CG with relaxation

Next, include one relaxation sweep. Note that error propagator  $R_\omega$  preserves the modes of  $A^h$

(although this is often unnecessary). Let  $\lambda_k$  denote the eigenvalue of  $R_\omega$  associated with  $w_k$ .

For  $k \leq n/2-1$ :

$$w_k \rightarrow \lambda_k \underbrace{(s_k)} w_k + \lambda_k \underbrace{(s_k)} w_{k'} \quad \text{Small!}$$

$$w_{k'} \rightarrow \underbrace{(\lambda_{k'})} c_k w_k + \underbrace{(\lambda_{k'})} c_k w_{k'} \quad \text{Small!}$$

$$s_k = \sin^2\left(\frac{k\pi}{2n}\right) \quad c_k = \cos^2\left(\frac{k\pi}{2n}\right)$$

## Crucial observation

- Between relaxation & coarse-grid correction, both smooth & oscillatory components of the error are effectively damped.
- This is the "spectral" picture of how multigrid works. We examine now another viewpoint, the "algebraic" picture of multigrid.

# Recall the variational properties

All the analysis that follows assumes that the variational properties hold:

$$A^{2h} = I_h^{2h} A^h I_h^h$$

$$I_{2h}^h = c (I_h^{2h})^T$$

Galerkin Condition

$$c \in \mathfrak{R}$$

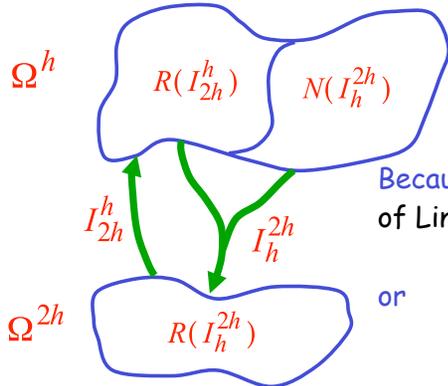
# Fundamental Theorem of Linear Algebra

$$N(B) = R(B^T)^\perp$$

- $x \in N(B)$  any  $y \in R(B^T)$  some  $z$   
 $\Rightarrow \langle x, y \rangle = \langle x, B^T z \rangle = \langle Bx, z \rangle = 0.$   
 $\Rightarrow x \in R(B^T)^\perp$   
 $\Rightarrow N(B) \subset R(B^T)^\perp.$
- $x \in R(B^T)^\perp$  any  $z$   $z = Bx$   
 $\Rightarrow 0 = \langle x, B^T z \rangle = \langle Bx, z \rangle = \langle Bx, Bx \rangle.$   
 $\Rightarrow x \in N(B)$   
 $\Rightarrow N(B) \supset R(B^T)^\perp.$

# Algebraic interpretation of CG

consider the subspaces that make up  $\Omega^h$  &  $\Omega^{2h}$



From now on, 'R( )' means Range & 'N( )' Null Space.

Because the Fundamental Theorem of Linear Algebra shows that

$$N(I_h^{2h}) = R((I_h^{2h})^T)^\perp$$

or

$$N(I_h^{2h}) = R(I_{2h}^h)^\perp$$

$I_{2h}^h$  transfers errors?

Does  $I_h^{2h}$  transfer errors?

So we really care about  $R(I_h^{2h} A^h)$  !?!

# "Energy" inner product & norm

- Inner product symmetry:  
 $\langle Ax, y \rangle = \langle x, Ay \rangle = \langle Ay, x \rangle.$
- Inner product linearity:  
 $\langle A(ax+by), z \rangle = \langle aAx+bAy, z \rangle = a\langle Ax, z \rangle + b\langle Ay, z \rangle.$
- Inner product positive definiteness:  
 $\langle Ax, x \rangle \geq 0$  &  $\langle Ax, x \rangle = 0 \Rightarrow x = 0.$
- Norm:  
 $\langle Ax, x \rangle$  is an inner product  $\Rightarrow \langle Ax, x \rangle^{1/2}$  is a norm.

## Subspace decomposition of $\Omega^h$

- If  $u^h \in N(I_h^{2h} A^h)$ , then, for any  $u^{2h}$ , we have

$$0 = \langle I_h^{2h} A^h u^h, u^{2h} \rangle = \langle A^h u^h, I_{2h}^h u^{2h} \rangle,$$

so

$$R(I_{2h}^h) \perp_{A^h} N(I_h^{2h} A^h),$$

where  $x \perp_{A^h} y$  means  $\langle A^h x, y \rangle = 0$ . "energy" inner product

- Moreover, any  $e^h$  can be written as  $e^h = s^h + t^h$ , where  $s^h \in R(I_{2h}^h)$  &  $t^h \in N(I_h^{2h} A^h)$ .
- Hence, we get the "energy-orthogonal" decomposition

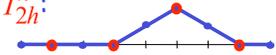
$$\Omega^h = R(I_{2h}^h) \oplus N(I_h^{2h} A^h).$$

errors touched by coarse grid

errors invisible to coarse grid

## Characteristics of the subspaces

- Since  $s^h = I_{2h}^h q^{2h}$  for some  $q^{2h} \in \Omega^{2h}$ , we associate  $s^h$  with the smooth components of  $e^h$ . But,  $s^h$  generally has all Fourier modes in it. Recall the basis vectors for  $I_{2h}^h$ :



- Similarly, we associate  $t^h$  with oscillatory components of  $e^h$ , although  $t^h$  generally has all Fourier modes in it as well. Recall that  $N(I_h^{2h} A^h)$  is spanned by  $\eta_i = A^h \hat{e}_i$ , so  $N(I_h^{2h} A^h)$  is spanned by the unit vectors  $\hat{e}_i^h = (0, 0, \dots, 0, 1, 0, \dots, 0)^T$  for odd  $i$ , which "look" oscillatory.



## Algebraic analysis CG

- Recall that (without relaxation)

$$CG = I - I_{2h}^h (A^{2h})^{-1} I_h^{2h} A^h.$$

- First note that if  $s^h \in R(I_{2h}^h)$ , then  $CG s^h = 0$ . This follows since  $s^h = I_{2h}^h q^{2h}$  for some  $q^{2h} \in \Omega^{2h}$  & therefore

$$CG s^h = \left[ I - I_{2h}^h (A^{2h})^{-1} I_h^{2h} A^h \right] I_{2h}^h q^{2h} = 0.$$

$\underbrace{\hspace{10em}}_{A^{2h} \text{ by Galerkin property}}$

- It follows that  $N(CG) = R(I_{2h}^h)$ , that is, the null space of  $CG$  is the range of interpolation.

What does this imply?

## More algebraic analysis of CG

- Next, note that if  $t^h \in N(I_h^{2h} A^h)$ , then

$$CG t^h = \left[ I - I_{2h}^h (A^{2h})^{-1} I_h^{2h} A^h \right] t^h$$

$\underbrace{\hspace{10em}}_0$

$$\Rightarrow CG t^h = t^h$$

- Thus,  $CG$  is the identity on  $N(I_h^{2h} A^h)$ .

What does this imply?

Together:  $CG(s^h + t^h) = t^h$

# How does the algebraic picture fit with the spectral view?

We may view  $\Omega^h$  in two ways:

$$\Omega^h = \left\{ \begin{array}{l} \text{Low frequency modes} \\ 1 \leq k \leq n/2 \end{array} \right\} \oplus \left\{ \begin{array}{l} \text{High frequency modes} \\ n/2 < k < n \end{array} \right\}$$

that is,

$$\Omega^h = L \oplus H$$

Illuminates relaxation (Jacobi)

or

$$\Omega^h = R(I_{2h}^h) \oplus N(I_h^{2h} A^h).$$

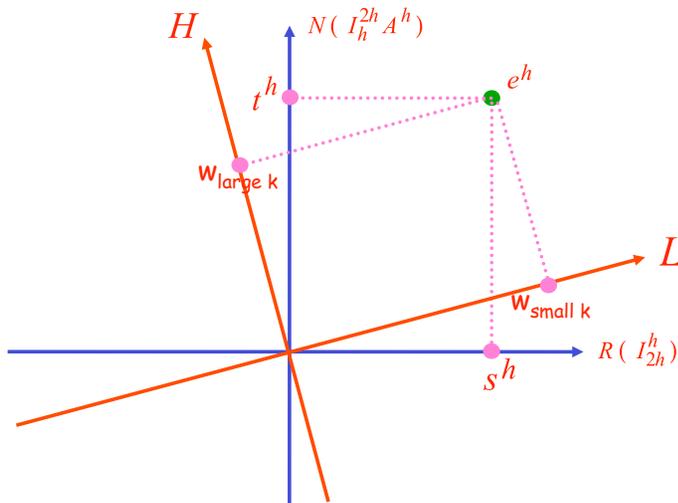
Illuminates CG (exact)

Are these "orthogonal" decompositions?

# Actually, each view is just part of the picture

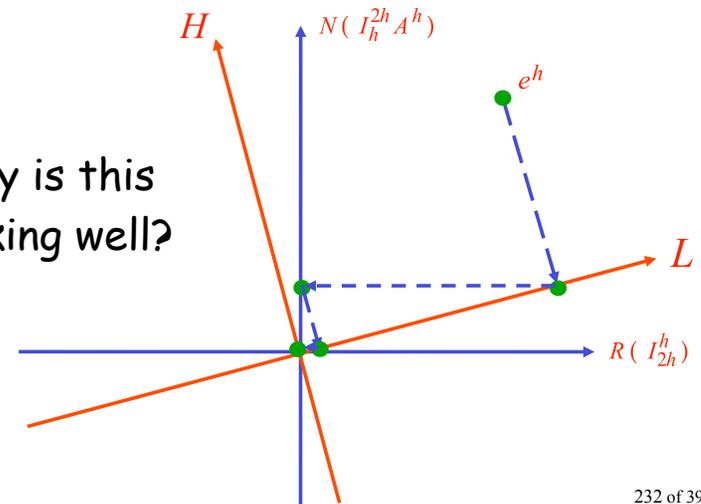
- The operations we've examined work on different spaces!
- While  $N(I_h^{2h} A^h)$  is mostly oscillatory, it isn't  $H$ , & while  $R(I_{2h}^h)$  is mostly smooth, it isn't  $L$ .
- Relaxation "eliminates" error from  $H$ .
- Coarse-grid correction eliminates error from  $R(I_{2h}^h)$ .

# How it actually works (cartoon)

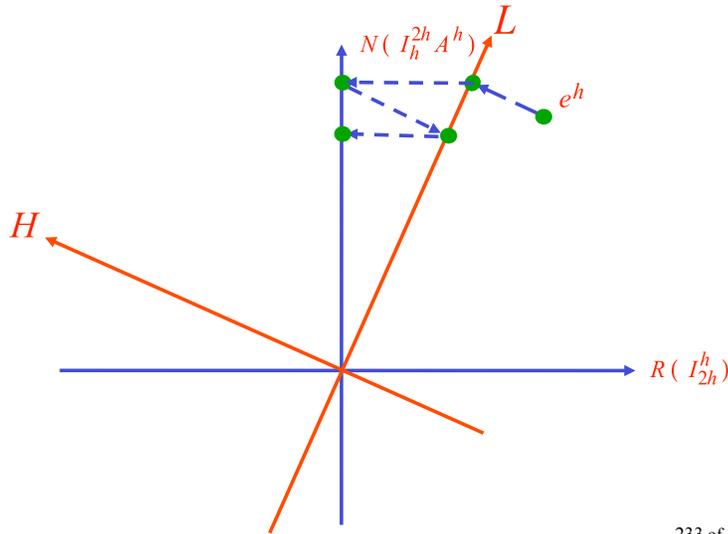


# Relaxation "eliminates" error in $R(I_{2h}^h)$ , but can increase the error in $N(I_h^{2h} A^h)$

Why is this working well?



# What if L points away from $R(I_{2h}^h)$ ?



# Outline

## Chapters 1-5:

- √ Model Problems
- √ Basic Iterative Methods
  - Convergence tests
  - Analysis
- √ Elements of Multigrid
  - Relaxation
  - Coarsening
- √ Implementation
  - Complexity
  - Diagnostics
- √ Some Theory
  - Spectral vs. algebraic

## Chapters 6-10:

- Nonlinear Problems
  - Full approximation scheme
- Selected Applications
  - Neumann boundaries
  - Anisotropic problems
  - Variable meshes
  - Variable coefficients
- Algebraic Multigrid (AMG)
  - Matrix coarsening
- Multilevel Adaptive Methods
  - FAC
- Finite Elements
  - Variational methodology

Homework Due!

## 6. Nonlinear problems

**HANG ON !!!**

- How should we approach the nonlinear system

$$A(u) = f$$

& can we use MG to solve it?

- A fundamental relation we've relied on is the linear residual equation:

$$Au - Av = f - Av \Rightarrow Ae = r.$$

- We can't rely on this now since a nonlinear  $A(u)$  generally means

$$A(u) - A(v) \neq A(e).$$

## The nonlinear residual equation

We still base our development around the residual equation, now the nonlinear residual equation:

$$\begin{aligned}
 & A(u) = f \\
 \Rightarrow & A(u) - A(v) = f - A(v) \\
 \Rightarrow & \boxed{A(u) - A(v) = r}
 \end{aligned}$$

How can we use this equation as the basis for a solution method?

# Newton's method for scalar $F: \mathfrak{R} \rightarrow \mathfrak{R}$

Best known & most important nonlinear solver!

- We wish to solve  $F(x) = 0$ .

Ex:  $F(x) = x e^x - 1, \quad F'(x) = (1+x)e^x$ .

↙ exponent, not error ↘

- Expand  $F$  in a Taylor series about  $x$ :

$$F(x+s) = F(x) + sF'(x) + s^2 F''(\xi)$$

Ex:  $(x+s)e^{(x+s)} - 1 = x e^x - 1 + s(1+x)e^x + \text{h.o.t.}$

- Dropping higher-order terms (h.o.t.), if  $x+s$  is a solution,

$$0 = F(x) + sF'(x) \Rightarrow s = -F(x)/F'(x)$$

- We thus arrive at Newton's method:

$$x \leftarrow x - F(x)/F'(x)$$

# Newton's method for systems

- Ex:  $-u''(x) + u(x) e^{u(x)} = f$  may be discretized as

$$A(u) \equiv \begin{pmatrix} a_1(u) \\ a_2(u) \\ \vdots \\ a_{n-1}(u) \end{pmatrix} \equiv \begin{pmatrix} \frac{2u_1 - u_2}{h^2} + u_1 e^{u_1} \\ \vdots \\ \frac{-u_{i-1} + 2u_i - u_{i+1}}{h^2} + u_i e^{u_i} \\ \vdots \\ \frac{-u_{n-2} + 2u_{n-1}}{h^2} + u_{n-1} e^{u_{n-1}} \end{pmatrix} = \begin{pmatrix} f(x_1) \\ \vdots \\ f(x_i) \\ \vdots \\ f(x_{n-1}) \end{pmatrix} \equiv f.$$

Jacobian

- Taylor series about  $v$ :  $A(v+e) = A(v) + J(v)e + \text{h.o.t.}$

$$J(v) \equiv \left( \frac{\partial a_i(u)}{\partial u_j} \right)_{u=v} = \begin{pmatrix} \frac{2}{h^2} + (1+v_1)e^{v_1} & \frac{-1}{h^2} & & \\ & \vdots & & \\ & \frac{-1}{h^2} & \frac{2}{h^2} + (1+v_i)e^{v_i} & \frac{-1}{h^2} \\ & & \vdots & \\ & & \frac{-1}{h^2} & \frac{2}{h^2} + (1+v_{n-1})e^{v_{n-1}} \end{pmatrix}$$

↑ error, not exponent ↑

## Newton for systems

- The system  $A(u) = f$  in vector form is

$$\begin{pmatrix} a_1(u_1, u_2, \dots, u_{n-1}) \\ a_2(u_1, u_2, \dots, u_{n-1}) \\ \vdots \\ a_N(u_1, u_2, \dots, u_{n-1}) \end{pmatrix} = \begin{pmatrix} f_1 \\ f_2 \\ \vdots \\ f_{n-1} \end{pmatrix}$$

- Expanding  $A(v+e)$  in a Taylor series about  $v$ :

$$A(v+e) = A(v) + J(v)e + \text{h.o.t.}$$

$$J(v) = \left( \frac{\partial a_i}{\partial u_j}(v) \right)$$

## Newton for systems (cont'd)

- $J(v)$  is the Jacobian

$$J(v) = \begin{pmatrix} \frac{\partial a_1}{\partial u_1} & \frac{\partial a_1}{\partial u_2} & \dots & \frac{\partial a_1}{\partial u_{n-1}} \\ \frac{\partial a_2}{\partial u_1} & \frac{\partial a_2}{\partial u_2} & \dots & \frac{\partial a_2}{\partial u_{n-1}} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\partial a_{n-1}}{\partial u_1} & \frac{\partial a_{n-1}}{\partial u_2} & \dots & \frac{\partial a_{n-1}}{\partial u_{n-1}} \end{pmatrix}_{u=v}$$

- If  $u = v + e$  is a solution,  $f = A(v) + J(v)e + \text{h.o.t.}$ , so

$$e \approx [J(v)]^{-1}(f - A(v)).$$

- This leads to the iteration

$$v \leftarrow v + [J(v)]^{-1}(f - A(v))$$

## Newton's via the residual equation

- The nonlinear residual equation is
 
$$A(v + e) - A(v) = r.$$
- Expanding  $A(v + e)$  in a two-term Taylor series about  $v$  & ignoring h.o.t.:
 
$$A(v) + J(v) \hat{e} - A(v) = r \quad \hat{e} \approx e$$
 or
 
$$J(v) \hat{e} = r.$$
- Newton's method is thus:

$$v \leftarrow v + [J(v)]^{-1} r, \quad r = f - A(v)$$

## How does multigrid fit in?

- One obvious method is to use multigrid to solve  $J(v) \hat{e} = r$  at each iteration step. This method is called Newton-MG & can be very effective.
- However, we might want to use multigrid ideas to treat the nonlinearity directly.
- To do that, we need to specialize multigrid components (relaxation & coarsening) for the nonlinear case.

## What is nonlinear relaxation?

$$A(u) = f$$

- Nonlinear Gauss-Seidel:
 

For each  $i = 1, 2, \dots, n-1$  :

Change the value of  $v_i$  so that the  $i^{\text{th}}$  equation is satisfied:  $(A(v))_i = f_i$ .
- Equivalently:
 

For each  $i = 1, 2, \dots, n-1$  :

Find  $s \in \mathbb{R}$  such that  $(A(v + s \varepsilon_i))_i = f_i$ ,

where  $\varepsilon_i$  is the  $i^{\text{th}}$  canonical unit basis vector.

## How is nonlinear Gauss-Seidel done?

- Each  $(A(v))_i = f_i$  is a nonlinear scalar equation for  $v_i$ . We can use the scalar Newton's method to solve!
- Example:  $-u''(x) + u(x) e^{u(x)} = f$  may be discretized so that  $(A(v))_i = f_i$  is given by
 
$$\frac{-v_{i-1} + 2v_i - v_{i+1}}{h^2} + v_i e^{v_i} = f_i \quad 1 \leq i \leq n-1$$
- Newton iteration for  $v_i$  is given by

$$v_i \leftarrow v_i - \frac{\frac{-v_{i-1} + 2v_i - v_{i+1}}{h^2} + v_i e^{v_i} - f_i}{\frac{2}{h^2} + (1 + v_i) e^{v_i}}$$

## How do we coarsen for nonlinear multigrid?

- Recall the nonlinear residual equation  $A(v + e) - A(v) = r$ .
- In multigrid, we obtain an approximate solution  $v^h$  on the fine grid, then solve the residual equation on the coarse grid.
- The residual equation on  $\Omega^{2h}$  appears as

$$A^{2h}(v^{2h} + e^{2h}) - A^{2h}(v^{2h}) = r^{2h}$$

We should have a routine for that.

## Look at the coarse residual equation

- We must evaluate the quantities on  $\Omega^{2h}$  in  $A^{2h}(v^{2h} + e^{2h}) - A^{2h}(v^{2h}) = r^{2h}$ .
- Given  $v^h$ , a fine-grid approximation, we restrict the residual to the coarse grid:  $r^{2h} = I_h^{2h}(f^h - A^h(v^h))$ .
- For  $v^{2h}$ , we restrict  $v^h$  by  $v^{2h} = I_h^{2h} v^h$ .

Thus,

$$A^{2h}(I_h^{2h} v^h + e^{2h}) = A^{2h}(I_h^{2h} v^h) + I_h^{2h}(f^h - A^h(v^h))$$

## We've obtained a coarse-grid equation of the form $A^{2h}(u^{2h}) = f^{2h}$

- Consider the coarse-grid equation:

$$A^{2h}(I_h^{2h} v^h + e^{2h}) = A^{2h}(I_h^{2h} v^h) + I_h^{2h}(f^h - A^h(v^h))$$

$u^{2h}$   
coarse-grid unknown

$f^{2h}$   
all quantities are known

- We solve  $A^{2h}(u^{2h}) = f^{2h}$  for  $u^{2h} = I_h^{2h} v^h + e^{2h}$  & obtain

$$e^{2h} = u^{2h} - I_h^{2h} v^h$$

- We then apply the correction:

$$v^h \leftarrow v^h + I_{2h}^h e^{2h}$$

## Full approximation scheme (FAS) 2-grid form

- Perform nonlinear relaxation on  $A^h(u^h) = f^h$  to obtain an approximation  $v^h$ .
- Restrict the approximation & its residual:  $v^{2h} = I_h^{2h} v^h$   $r^{2h} = I_h^{2h}(f^h - A^h(v^h))$ .
- Solve the coarse-grid equation:  $A^{2h}(u^{2h}) = A^{2h}(v^{2h}) + r^{2h} = I_h^{2h} f^h + A^{2h}(I_h^{2h} v^h) - I_h^{2h} A^h(v^h)$ .
- Extract 2h approximation to h error:  $e^{2h} = u^{2h} - v^{2h}$ .
- Interpolate & correct:  $v^h \leftarrow v^h + I_{2h}^h e^{2h}$ .

## A few observations about FAS

$$A^{2h}(I_h^{2h}v^h + e^{2h}) = A^{2h}(I_h^{2h}v^h) + I_h^{2h}(f^h - A^h(v^h))$$

- If  $A$  is a linear operator, then FAS reduces directly to the linear two-grid correction scheme:

$$A^{2h}(I_h^{2h}v^h + e^{2h}) = A^{2h}(I_h^{2h}v^h) + I_h^{2h}(f^h - A^h(v^h))$$

- An exact solution to the fine-grid problem is a fixed point of the FAS iteration:

$$A^{2h}(I_h^{2h}v^h + e^{2h}) = A^{2h}(I_h^{2h}v^h) + I_h^{2h}(f^h - A^h(v^h))$$

## Still more observations about FAS

- A true multilevel FAS process is recursive, using FAS to solve the nonlinear  $\Omega^{2h}$  problem using  $\Omega^{4h}$ .
- Hence, FAS is generally employed in a V- or W-cycling scheme.

## A few more observations about FAS

$$A^{2h}(I_h^{2h}v^h + e^{2h}) = A^{2h}(I_h^{2h}v^h) + I_h^{2h}(f^h - A^h(v^h))$$

$\tau_h^{2h} \leftarrow A(Pv^h) - PA(v^h)$

- The FAS coarse-grid equation can be written as

$$A^{2h}(u^{2h}) = f^{2h} + \tau_h^{2h}$$

where  $\tau_h^{2h} = A^{2h}(I_h^{2h}v^h) - I_h^{2h}A^h(v^h)$  is the so-called tau correction term &  $f^{2h}$  is the original 2h source term, provided you choose it that way:  $f^{2h} = I_h^{2h}f^h$ .

- In general, since  $\tau_h^{2h} \neq 0$ , the solution  $u^{2h}$  to the FAS coarse-grid equation is not the same as the solution to the original coarse-grid problem

$$A^{2h}(u^{2h}) = f^{2h}$$

- The tau correction is as a way to alter the coarse-grid equation to enhance its approximation properties.

## Even more observations about FAS

- For linear problems, we use FMG to obtain a good initial guess on the fine grid. Convergence of nonlinear iterations depends critically on having a good initial guess.
- When FMG is used for nonlinear problems, the interpolant  $I_{2h}^h u^{2h}$  is generally accurate enough to be in the basin of attraction of the fine-grid solver.
- Thus, whether FAS, Newton, or Newton-multigrid is used on each level, one FMG cycle should provide a solution accurate to the level of discretization, unless the nonlinearity is extremely strong.

## Intergrid transfers for FAS

- Generally speaking, the standard operators (linear interpolation, full weighting) work effectively in FAS schemes.
- For strongly nonlinear problems or for the coarse-grid approximation that is to become a fine-grid initial guess, higher-order interpolation (e.g., cubic interpolation) may be beneficial.

## Example: Newton-MG vs. FAS

- PDE (er, ODE):

$$-u''(x) + u(x) e^{u(x)} = f(x).$$

- Discretization:

$$\frac{-v_{j-1} + 2v_j - v_{j+1}}{h^2} + v_j e^{v_j} = f_j.$$

## What is $A^{2h}(u^{2h})$ in FAS?

As in the linear case, there are two basic possibilities:

1.  $A^{2h}(u^{2h})$  is determined by discretizing the nonlinear operator,  $A(u)$ , in the same fashion as was employed to obtain  $A^h(u^h)$ , except that the coarser mesh spacing is used.

2.  $A^{2h}(u^{2h})$  is determined from the Galerkin condition

$$A^{2h}(u^{2h}) = I_h^{2h} A^{2h}(I_{2h}^h u^{2h})$$

where the action of the Galerkin product can be captured in an implementable formula.

The first method is usually easier & more common.

## One Newton-MG step

- **Step 0.** Given  $v$ , form the grid  $h$  linear correction equation:

$$\frac{-e_{j-1} + 2e_j - e_{j+1}}{h^2} + (1+v_j)e^{v_j} e_j = r_j \equiv f_j - \left( \frac{-v_{j-1} + 2v_j - v_{j+1}}{h^2} + v_j e^{v_j} \right)$$

Initialize the Newton correction approximation:  $e = 0$ .

- **Step 1:** Relax on the grid  $h$  linear equation.

- **Step 2:** Solve the grid  $2h$  error correction equation:

$$\frac{-e_{j-1}^{2h} + 2e_j^{2h} - e_{j+1}^{2h}}{(2h)^2} + (1+v_{2j})e^{v_{2j}} e_j^{2h} = r_{2j} - \left( \frac{-e_{2j-1} + 2e_{2j} - e_{2j+1}}{h^2} + (1+v_{2j})e^{v_{2j}} e_{2j} \right)$$

- **Step 3:** Correct the grid  $h$  Newton correction:

$$e \leftarrow e + I_{2h}^h e^{2h}$$

- **Step 4:** Stop if you've "solved" the linear equation well enough for Newton correction  $e$  & set  $v \leftarrow v + e$ . Else, leave  $v$  alone & return to **Step 1**.

## One FAS step

- Step 0. Given  $v$ , form the grid  $h$  nonlinear equation:

$$\frac{-v_{j-1} + 2v_j - v_{j+1}}{h^2} + v_j e^{v_j} = f_j$$

- Step 1: Relax on the grid  $h$  nonlinear equation to improve  $v$ .

$$v_j \leftarrow v_j - \frac{\frac{-v_{j-1} + 2v_j - v_{j+1}}{h^2} + v_j e^{v_j} - f_j}{\frac{2}{h^2} + (1 + v_j)e^{v_j}}$$

- Step 2: Solve the grid  $2h$  FAS correction equation:

$$\frac{-e_{j-1}^{2h} + 2e_j^{2h} - e_{j+1}^{2h}}{(2h)^2} + (v_{2j}^h + e_j^{2h})e^{v_{2j}^h} - v_{2j}^h e^{v_{2j}^h} = (I_h^h r^h)_j$$

unknowns

- Step 3: Correct the grid  $h$  approximation  $v$ :

$$v \leftarrow v + I_{2h}^h e^{2h}$$

## Nonlinear problems: 2d example

- Consider

$$-\Delta u(x,y) + \gamma u(x,y) e^{u(x,y)} = f(x,y)$$

on the unit square,  $[0,1] \times [0,1]$ , with homogeneous Dirichlet boundary conditions & a regular  $h = 1/128$  Cartesian grid.

- Suppose the exact solution is

$$u(x,y) = (x^2 - x^3) \sin(3\pi y)$$

## Discretization of the nonlinear example

- The operator can be written (sloppily) as

$$\frac{1}{h^2} \begin{pmatrix} -1 & 4 & -1 \\ -1 & 4 & -1 \end{pmatrix} u_{i,j}^h + \gamma u_{i,j}^h e^{u_{i,j}^h} = f_{i,j}$$

$$(A^h(v^h))_{i,j}$$

- Relaxation (nonlinear Gauss-Seidel) is given by

$$v_{i,j}^h \leftarrow v_{i,j}^h - \frac{(A^h(v^h))_{i,j} - f_{i,j}}{\frac{4}{h^2} + \gamma(1 + v_{i,j}^h)e^{v_{i,j}^h}}$$

## FAS & Newton's method on

$$-\Delta u(x,y) + \gamma u(x,y) e^{u(x,y)} = f(x,y)$$

$n = 128$

convergence factor is for the last cycle

- FAS  $V(2,1)$ -cycles until  $\|r\| < 10^{-10}$ .

	$\gamma$			
	1	10	100	1000
convergence factor	0.135	0.124	0.098	0.072
number of FAS cycles	12	11	11	10

- Newton's Method with exact inner solves until  $\|r\| < 10^{-10}$ .

	$\gamma$			
	1	10	100	1000
convergence factor	4.00E-05	7.00E-05	3.00E-04	2.00E-04
number of Newton iterations	3	3	3	4

## Newton, Newton-MG, & FAS on

$$-\Delta u(x,y) + \gamma u(x,y) e^{u(x,y)} = f(x,y)$$

$$n = 128, \gamma = 10$$

- Newton uses exact solves, Newton-MG is with a fixed number of inner V(2,1)-cycles for the Jacobian problem, overall stopping criterion  $\|r\| < 10^{-10}$ .

Method	Outer iterations	Inner iterations	Megaflops
Newton	3		1660.6
Newton-MG	3	20	56.4
Newton-MG	4	10	38.5
Newton-MG	5	5	25.1
Newton-MG	10	2	22.3
Newton-MG	19	1	24.6
FAS	11		27.1

## Compare FMG-FAS & FMG-Newton-MG

$$-\Delta u(x,y) + \gamma u(x,y) e^{u(x,y)} = f(x,y)$$

- We do one FMG cycle using one FAS V(2,1) -cycle as the "solver" at each new level. We then follow that with as many FAS V(2,1)-cycles as is needed to obtain  $\|r\| < 10^{-10}$ .
- Next, we do one FMG cycle using a Newton-MG step at each new level (with one linear V(2,1)-cycle as the Jacobian "solver.") We then follow that with as many Newton-multigrid steps as is needed to obtain  $\|r\| < 10^{-10}$ .

Don't try this at home !!!

## Compare FMG-FAS & FMG-Newton-MG

$$-\Delta u(x,y) + \gamma u(x,y) e^{u(x,y)} = f(x,y)$$

$$n = 128, \gamma = 10$$

Done !!!

Cycle	$\ r\ $	$\ u^{(n)} - v\ $	Mflops	$\ r\ $	$\ u^{(n)} - v\ $	Mflops	Cycle
FMG-FAS	1.10E-02	2.00E-05	3.1	1.06E-02	2.50E-05	2.4	FMG-Newton
FAS V	6.80E-04	2.40E-05	5.4	6.70E-04	2.49E-05	4.1	Newton-MG
FAS V	5.00E-05	2.49E-05	7.6	5.10E-05	2.49E-05	5.8	Newton-MG
FAS V	3.90E-06	2.49E-05	9.9	6.30E-06	2.49E-05	7.5	Newton-MG
FAS V	3.20E-07	2.49E-05	12.2	1.70E-06	2.49E-05	9.2	Newton-MG
FAS V	3.00E-08	2.49E-05	14.4	5.30E-07	2.49E-05	10.9	Newton-MG
FAS V	2.90E-09	2.49E-05	16.7	1.70E-07	2.49E-05	12.6	Newton-MG
FAS V	3.00E-10	2.49E-05	18.9	5.40E-08	2.49E-05	14.3	Newton-MG
FAS V	3.20E-11	2.49E-05	21.2	1.70E-08	2.49E-05	16.0	Newton-MG
				5.50E-09	2.49E-05	17.7	Newton-MG
				1.80E-09	2.49E-05	19.4	Newton-MG
				5.60E-10	2.49E-05	21.1	Newton-MG
				1.80E-10	2.49E-05	22.8	Newton-MG
				5.70E-11	2.49E-05	24.5	Newton-MG

## Remembering coarse-grid correction

- Relax (damped Jacobi) to smooth  $e = u - v$ :

$$v \leftarrow v - \omega D^{-1}(Av - f).$$

- Form the residual equation  $Ae = r$ :

$$Ae = A(u - v) = f - Av = r.$$

- Use premise that smooth error  $\Rightarrow e = I_{2h}^h e^{2h}$ :

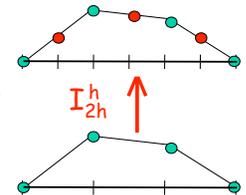
$$AI_{2h}^h e^{2h} = r.$$

fewer unknowns

- Use transpose  $I_h^{2h} = (I_{2h}^h)^T$  to reduce equations:

$$I_h^{2h} AI_{2h}^h e^{2h} = I_h^{2h} r.$$

fewer equations



$A^{2h}$







## The new system is symmetric

- We have the symmetric system  $\hat{A}^h u^h = \hat{f}^h$ :

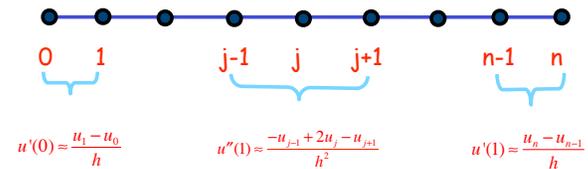
$$\frac{1}{h^2} \begin{pmatrix} 1 & -1 & & & & \\ -1 & 2 & -1 & & & \\ & -1 & 2 & -1 & & \\ & & \ddots & \ddots & \ddots & \\ & & & -1 & 2 & -1 \\ & & & & -1 & 1 \end{pmatrix} \begin{pmatrix} u_0^h \\ u_1^h \\ u_2^h \\ \vdots \\ u_{n-1}^h \\ u_n^h \end{pmatrix} = \begin{pmatrix} f_0^h/2 \\ f_1^h \\ f_2^h \\ \vdots \\ f_{n-1}^h \\ f_n^h/2 \end{pmatrix}$$

- Solvability is guaranteed by ensuring that  $\hat{f}^h$  is orthogonal to the constant vector  $\mathbf{1}^h$ :

$$\langle \hat{f}^h, \mathbf{1}^h \rangle = \sum_{j=0}^n \hat{f}_j^h = 0$$

## One-sided differences at boundary a similar result

- No ghost points:



- This yields the system

$$\begin{aligned} \frac{-u_{j-1} + 2u_j - u_{j+1}}{h^2} &= f_j & 1 \leq j \leq n \\ \frac{u_1 - u_0}{h^2} &= 0 & \frac{u_n - u_{n-1}}{h^2} = 0 \end{aligned}$$

## The well-posed discrete system

back to central differences @ boundary

- The  $(n+2) \times (n+1)$  system is:

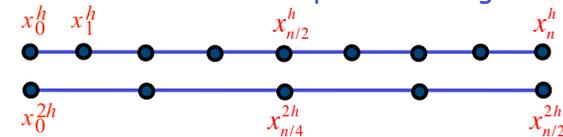
$$\begin{aligned} \frac{-u_{j-1} + 2u_j - u_{j+1}}{h^2} &= f_j & 1 \leq j \leq n-1 \\ \frac{u_0 - u_1}{h^2} &= 0 & \text{for one-sided scheme} \\ \frac{-u_{n-1} + u_n}{h^2} &= 0 \\ \sum_{j=0}^n u_j^h &= 0 & \text{(choose the zero mean solution)} \end{aligned}$$

or, more simply

$$\begin{aligned} \hat{A}^h u^h &= \hat{f}^h \\ \langle u^h, \mathbf{1}^h \rangle &= 0 \end{aligned}$$

## Multigrid for the Neumann problem

- We must have the interval endpoints on all grids



- Relaxation is performed at all points, including endpoints:

$$v_0^h \leftarrow v_1^h + h^2 \hat{f}_0^h \quad v_j^h \leftarrow \frac{v_{j-1}^h + v_{j+1}^h + h^2 \hat{f}_j^h}{2} \quad v_n^h \leftarrow v_{n-1}^h + h^2 \hat{f}_n^h$$

- We add a global Gram-Schmidt single step after relaxation on each level to enforce the zero-mean condition:

$$v^h \leftarrow v^h - \frac{\langle v^h, \mathbf{1}^h \rangle}{\langle \mathbf{1}^h, \mathbf{1}^h \rangle} \mathbf{1}^h$$



# Neumann problem: An example

Consider the problem

$$-u''(x) = 2x - 1, \quad 0 < x < 1, \quad u'(0) = u'(1) = 0,$$

which has  $u(x) = \frac{x^2}{2} - \frac{x^3}{3} + c$  as a solution for any  $c$  ( $c = -1/12$  gives the zero mean solution).

grid size n	$\ r^h\ _h$	average conv. factor	$\ u^{(h)} - v^h\ _h$	number of cycles
32	6.30E-11	0.079	9.70E-05	9
64	1.90E-11	0.089	2.40E-05	10
128	2.60E-11	0.093	5.90E-06	10
256	3.70E-11	0.096	1.50E-06	10
512	5.70E-11	0.100	3.70E-07	10
1024	8.60E-11	0.104	9.20E-08	10
2048	2.10E-11	0.112	2.30E-08	10
4096	5.20E-11	0.122	5.70E-09	10

V(2,1)  
cycles

# 7. Selected applications

## 7b. Anisotropic problems

- All problems considered thus far have had  $-h^{-2}$  as the off-diagonal entries.

- We consider two situations when the matrix has two different constants on the off-diagonals. These situations arise when

- the (2-d) differential equation has constant but different coefficients for the derivatives in the coordinate directions
- the discretization has constant but different mesh spacing in the different coordinate directions

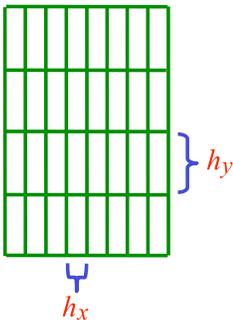
## We consider two types of anisotropy

- Different coefficients on the derivatives

$$-u_{xx} - \epsilon u_{yy} = f$$

discretized on a uniform grid with spacing  $h$ .

- Different mesh spacings:



$$h_x = h = \frac{1}{n}$$

$$h_y = \frac{h_x}{\sqrt{\epsilon}}$$

## Both problems lead to the same stencil

$$\frac{-u_{j-1,k} + 2u_{j,k} - u_{j+1,k}}{h^2} + \epsilon \frac{-u_{j,k-1} + 2u_{j,k} - u_{j,k+1}}{h^2}$$



$$A^h = \frac{1}{h^2} \begin{pmatrix} & -\epsilon & \\ -1 & 2 + 2\epsilon & -1 \\ & -\epsilon & \end{pmatrix}$$



$$\frac{-u_{j-1,k} + 2u_{j,k} - u_{j+1,k}}{h^2} + \frac{-u_{j,k-1} + 2u_{j,k} - u_{j,k+1}}{\left(\frac{h}{\sqrt{\epsilon}}\right)^2}$$

## Why standard multigrid can fail

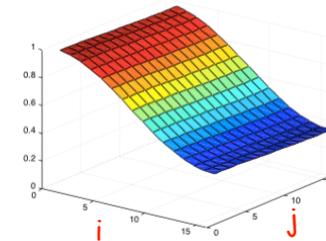
$$A^h = \frac{1}{h^2} \begin{pmatrix} & -\varepsilon & \\ -1 & 2+2\varepsilon & -1 \\ & -\varepsilon & \end{pmatrix}$$

- Note that  $A^h$  has weak connections in the  $y$ -direction. MG convergence factors degrade as  $\varepsilon$  gets small, with poor performance already at  $\varepsilon = 0.1$ .
- Consider the limiting case  $\varepsilon \Rightarrow 0$ :  $A^h = \frac{1}{h^2} \begin{pmatrix} 0 & & \\ -1 & 2 & -1 \\ & 0 & 0 \end{pmatrix}$
- Collection of disconnected 1-D problems!
- Point relaxation smoothes oscillatory errors in the  $x$ -direction (strong connections), but with no connections in the  $y$ -direction, the errors in that direction will generally be random; point relaxation provides no smoothing in the  $y$ -direction.

## We analyze weighted Jacobi

The eigenvalues of the weighted Jacobi iteration matrix for this problem are

$$\lambda_{i,j} = 1 - \frac{2\omega}{1+\varepsilon} \left( \sin^2\left(\frac{i\pi}{2n}\right) + \varepsilon \left( \sin^2\left(\frac{j\pi}{2n}\right) \right) \right)$$

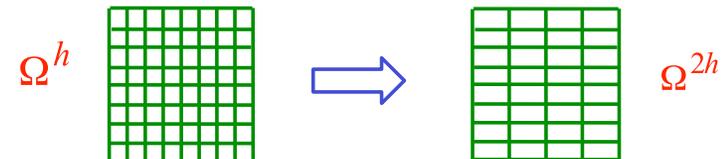


## Two strategies for anisotropy

- **Semicoarsening:** The equations are weakly coupled in the  $y$ -direction, so we can't expect the error after point relaxation to have any connection to the errors above or below it. We therefore can & should coarsen only in the  $x$ -direction.
- **Line relaxation:** The equations are strongly coupled in the  $x$ -direction, so we could solve simultaneously for all the unknowns along lines of constant  $y$ . This should expose whatever weak smoothness there might be in the  $x$ -direction, which should allow standard coarsening.

## Semicoarsening with point relaxation

- Point relaxation on  $A^h = \frac{1}{h^2} \begin{pmatrix} & -\varepsilon & \\ -1 & 2+2\varepsilon & -1 \\ & -\varepsilon & \end{pmatrix}$  smoothes in the  $x$ -direction. Coarsen by removing every other  $y$ -line.



- We do **not** coarsen along the remaining  $y$ -lines.
- Semicoarsening is not as "fast" as full coarsening. The number of points on  $\Omega^{2h}$  is about half the number of points on  $\Omega^h$ , instead of the usual one-fourth.

## Interpolation with semicoarsening

- We interpolate in the 1-D way along each line of constant  $y$ .
- The formulas for interpolating the correction from the coarse to the fine grid for the 2D model problem are

$$v_{2j,k}^h = v_{2j,k}^h + v_{j,k}^{2h}$$

$$v_{2j+1,k}^h = v_{2j+1,k}^h + \frac{v_{j,k}^{2h} + v_{j+1,k}^{2h}}{2}$$

## Line relaxation with full coarsening

$$A^h = \frac{1}{h^2} \begin{pmatrix} & & -\varepsilon \\ -1 & 2+2\varepsilon & -1 \\ & & -\varepsilon \end{pmatrix}$$

- The other approach to this problem is to do the usual full coarsening, but to relax entire  $x$ -lines (constant  $y$ ) of variables simultaneously.
- Consider an  $x$ -line equation specified by a fixed  $j$ :

$$\begin{matrix} & \xrightarrow{-v_{i,j+1}} & & \\ -v_{i-1,j} & + (2+2\varepsilon)v_{i,j} & -v_{i+1,j} & = h^2 f_{i,j} \\ & \xleftarrow{-v_{i,j-1}} & & \end{matrix} \quad 1 \leq i \leq n-1$$

$$\xrightarrow{\hspace{2cm}} -v_{i-1,j} + (2+2\varepsilon)v_{i,j} - v_{i+1,j} = h^2 f_{i,j} + v_{i,j+1} + v_{i,j-1}$$

## Line relaxation

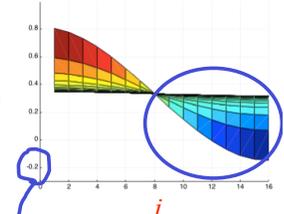
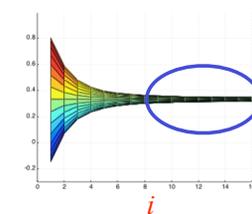
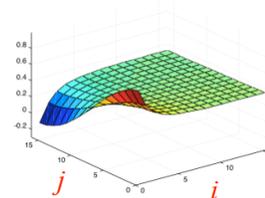
$$-v_{i-1,j} + (2+2\varepsilon)v_{i,j} - v_{i+1,j} = h^2 f_{i,j} + v_{i,j+1} + v_{i,j-1}$$

- Nice 1D system, analogous to the discretization of  $-u'' + \alpha u = g$ ,  $\alpha = 2\varepsilon h^{-2} > 0!$
- One sweep of line relaxation consists of solving a tridiagonal system for each constant  $y$ . Total cost is an optimal  $O(n^2)$ .
- Each solve can be done by Gaussian elimination since the system is tridiagonal, or a 1D multigrid solver (useful for generalization to higher dimensions).
- The individual lines can be solved simultaneously in a Jacobi way or sequentially in a Gauss-Seidel way.

## Why line relaxation works

Eigenvalues of the weighted block Jacobi iteration matrix:

$$\lambda_{i,j} = 1 - \frac{2\omega}{2 \sin^2\left(\frac{i\pi}{2n}\right) + \varepsilon} \left( \sin^2\left(\frac{i\pi}{2n}\right) + \varepsilon \left( \sin^2\left(\frac{j\pi}{2n}\right) \right) \right)$$



0.2

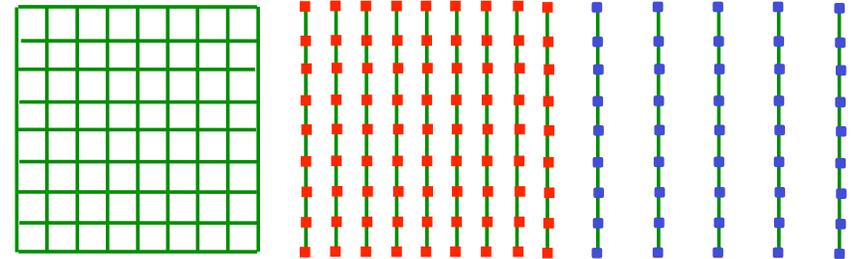
## Semicoarsening & line relaxation

- We might not know the direction of weak coupling or it might vary over the domain.
- Suppose we want a method that can handle either

$$A_1^h = \frac{1}{h^2} \begin{pmatrix} & -\varepsilon & \\ -1 & 2+2\varepsilon & -1 \\ & -\varepsilon & \end{pmatrix} \quad \text{or} \quad A_2^h = \frac{1}{h^2} \begin{pmatrix} & & -1 \\ -\varepsilon & 2+2\varepsilon & -\varepsilon \\ & & -1 \end{pmatrix}$$

- We could use semicoarsening in the  $x$ -direction to handle  $A_1^h$  & line relaxation in the  $y$ -direction to take care of  $A_2^h$ .

## Semicoarsening & line relaxation



- The original grid.
- Original grid viewed as a stack of "pencils." Line relaxation is used to solve problem along each "pencil".
- Coarsening is done by deleting every other pencil.

## An anisotropic example

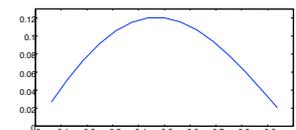
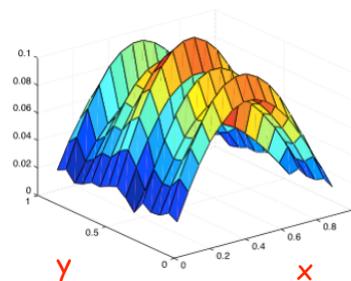
- Consider  $-u_{xx} - \varepsilon u_{yy} = f$  with  $u = 0$  on the boundaries of the unit square, & stencil given by

$$A^h = \frac{1}{h^2} \begin{pmatrix} & -\varepsilon & \\ -1 & 2+2\varepsilon & -1 \\ & -\varepsilon & \end{pmatrix}$$

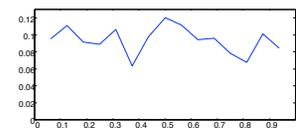
- Suppose that  $f(x,y) = 2(y - y^2) + 2\varepsilon(x - x^2)$  so that the exact solution is  $u(x,y) = (y - y^2)(x - x^2)$ .
- Note: If  $\varepsilon$  is small, then the  $x$ -direction dominates, while if  $\varepsilon$  is large, then the  $y$ -direction dominates.

## What is smooth error?

- Consider  $\varepsilon = 0.001$  & suppose point Gauss-Seidel is applied to a random initial guess. The error after 50 sweeps appears as:



Error along line of constant  $x$



Error along line of constant  $y$

## We experiment with 3 methods

- Standard V(2,1)-cycling, with point Gauss-Seidel relaxation, full coarsening, & linear interpolation.
- Semicoarsening in the  $x$ -direction. Coarse & fine grids have the same number of points in the  $y$ -direction. 1-D full weighting & linear interpolation are used in the  $x$ -direction, with no  $y$ -coupling in the intergrid transfers.
- Semicoarsening in the  $x$ -direction combined with line relaxation in the  $y$ -direction. 1-D full weighting & interpolation.

## With semicoarsening, the operator must change

- To account for unequal mesh spacing, the residual & relaxation operators must use a modified stencil:

$$A = \begin{pmatrix} & -\frac{\epsilon}{h_y^2} & \\ -\frac{1}{h_x^2} & \left( \frac{2}{h_x^2} + \frac{2\epsilon}{h_y^2} \right) & -\frac{1}{h_x^2} \\ & -\frac{\epsilon}{h_y^2} & \end{pmatrix}$$

- Note that, as grids become coarser,  $h_x$  grows while  $h_y$  remains constant.

## How do the 3 methods work for various values of $\epsilon$ ?

$n = 16$

Asymptotic convergence factors of V(2,1)-cycles:

scheme	1000	100	10	$\epsilon$ 1	0.1	0.01	0.001	1E-04
standard method	0.95	0.94	0.58	0.13	0.58	0.90	0.95	0.95
x-semi	0.94	0.99	0.98	0.93	0.71	0.28	0.07	0.07
x-semi & line relax	0.04	0.08	0.08	0.08	0.07	0.07	0.08	0.08

$y$ -direction strong

$x$ -direction strong

why is this bad???

Note: semicoarsening in  $x$  works well for  $\epsilon < 0.001$  but degrades noticeably even at  $\epsilon = 0.1$ .

Poisson

## A semicoarsening subtlety

- Suppose  $\epsilon$  is small, so that semicoarsening in  $x$  is used. As we progress to coarser grids,  $h_x^{-2}$  gets small but  $h_y^{-2}$  remains constant.
- If, on some coarse grid,  $h_x^{-2}$  becomes comparable to  $\epsilon h_y^{-2}$ , then the problem effectively becomes recoupled in the  $y$ -direction. Continued semicoarsening can produce artificial anisotropy, strong in the  $y$ -direction.
- When this occurs, it is best to stop semicoarsening & use full coarsening on any further coarse grids.

## 7. Selected applications

### 7c. Variable meshes

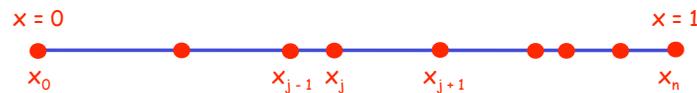
- Non-uniform grids are commonly used for domain or data irregularities or emerging solution features.

- Consider how we might approach the 1-D problem

$$-u''(x) = f(x), \quad 0 < x < 1,$$

$$u(0) = u(1) = 0$$

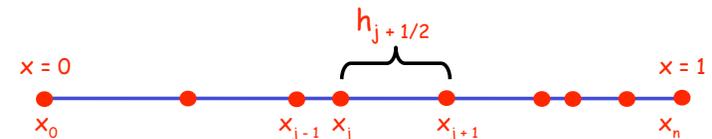
posed on the following nonuniform grid:



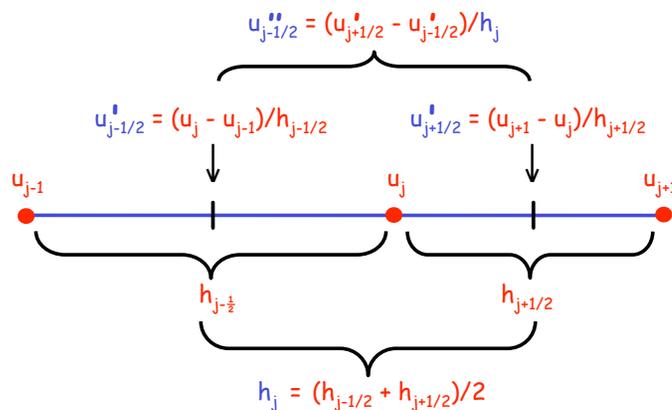
## We need some notation for the mesh spacing

Let  $n$  be a positive integer. We define the spacing interval between  $x_j$  &  $x_{j+1}$ :

$$h_{j+1/2} \equiv x_{j+1} - x_j, \quad j = 0, 1, \dots, n-1.$$



## Building second divided differences



## The discrete differential operator

- Using 2<sup>nd</sup>-order finite differences (& messy algebra!), we obtain the discrete representation

$$-\alpha_j^h u_{j-1}^h + (\alpha_j^h + \beta_j^h) u_j^h - \beta_j^h u_{j+1}^h = f_j^h \quad 1 \leq j \leq n-1$$

$$u_0^h = u_n^h = 0$$

where

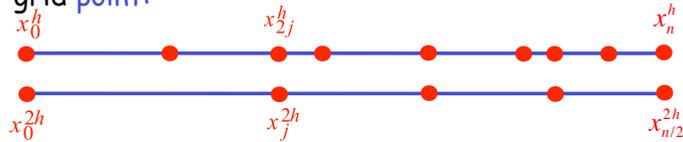
$$\alpha_j^h = \frac{2}{h_{j-\frac{1}{2}}(h_{j-\frac{1}{2}} + h_{j+\frac{1}{2}})} \quad \& \quad \beta_j^h = \frac{2}{h_{j+\frac{1}{2}}(h_{j-\frac{1}{2}} + h_{j+\frac{1}{2}})} .$$

- Multiplying by  $(h_{j-1/2} + h_{j+1/2})/2$  yields an SPD matrix with stencil

$$A^h = \begin{bmatrix} -\frac{1}{h_{j-\frac{1}{2}}} & \frac{1}{h_{j-\frac{1}{2}}} + \frac{1}{h_{j+\frac{1}{2}}} & -\frac{1}{h_{j+\frac{1}{2}}} \end{bmatrix} .$$

## Modify interpolation for variable $h$

- We choose every other fine-grid point as a coarse-grid point:



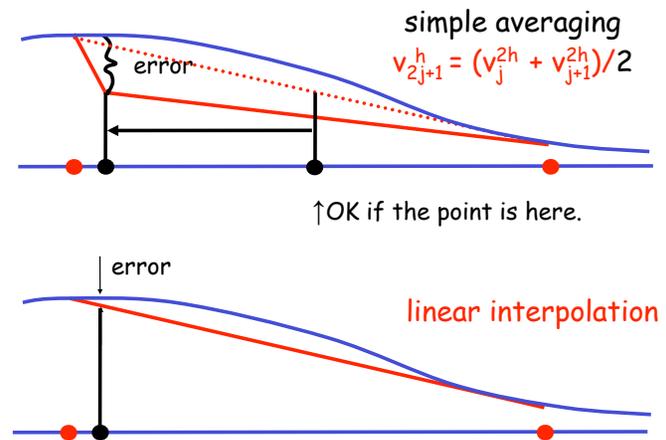
- In  $[x_0^h, x_{2j}^h]$ , linear means

$$v(x) = v_0^{2h} + (v_1^{2h} - v_0^{2h}) (x - x_0^h) / (x_{2j}^h - x_0^h).$$

- Plug in  $x = x_1^h$ : writing  $v^h = I_{2h}^h v^h$  yields

$$v_{2j}^h = v_j^{2h}, \quad v_{2j+1}^h = \frac{h_{2j+3/2} v_j^{2h} + h_{2j+1/2} v_{j+1}^{2h}}{h_{2j+1/2} + h_{2j+3/2}}, \quad 1 \leq j \leq n/2-1.$$

## Proper linear interpolation is needed

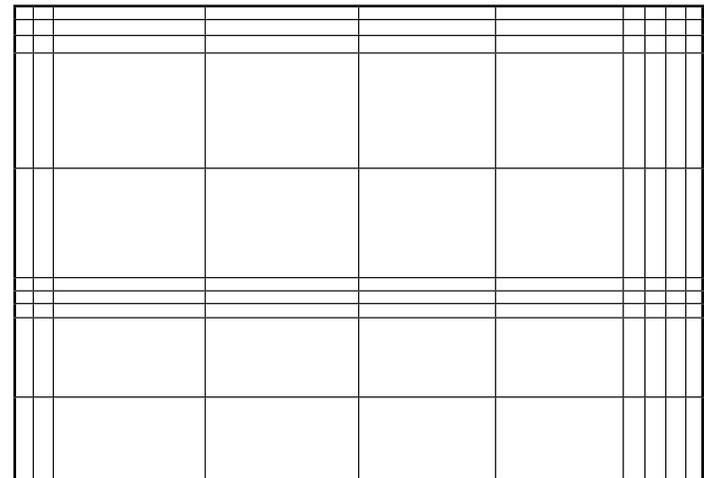


## We use the variational properties to derive restriction & $A^{2h}$

$$A^{2h} = I_h^{2h} A^h I_{2h}^h \quad I_h^{2h} = \frac{1}{2} (I_{2h}^h)^T$$

- This produces a stencil on  $\Omega^{2h}$  that is similar, but not identical, to the fine-grid stencil. If the resulting system is scaled by  $(h_{j-1/2} + h_{j+1/2})$ , then the Galerkin product is the same as the fine-grid stencil.
- For 2-D problems, this approach can be generalized readily to tensor-product grids. However, for general irregular grids, AMG is a better choice.

## Tensor-product grids?



## 7. Selected applications

### 7d. Variable coefficients

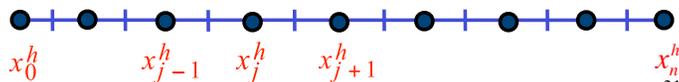
- A common difficulty is variable coefficients, exemplified in 1-D by

$$-(a(x) u'(x))' = f(x), \quad 0 < x < 1, \\ u(0) = u(1) = 0,$$

where  $a(x)$  is a positive function on  $[0,1]$ .

- We seek to develop a conservative, or self-adjoint, method for discretizing this problem.

- Assume we have available to us the values of  $a(x)$  ( $a_{j+1/2} \equiv a(x_{j+1/2})$ ) at midpoints of the uniform grid



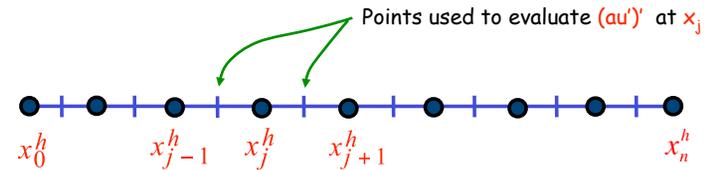
CU-Boulder

313 of 396

## Discretize using central differences

We can use second-order differences to approximate the derivatives. To use a grid spacing of  $h$ , we evaluate  $a(x)u'(x)$  at points midway between the gridpoints:

$$(a(x)u'(x))' \Big|_{x_j} = \frac{(au')|_{x_{j+1/2}} - (au')|_{x_{j-1/2}}}{h} + O(h^2)$$



CU-Boulder

314 of 396

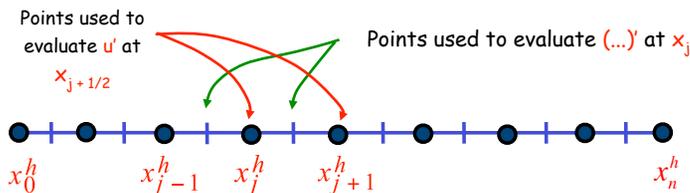
## Discretize using central differences (cont'd)

To evaluate  $(au')|_{x_{j+1/2}}$ , we must sample  $a(x)$  at the point  $x_{j+1/2}$  & use second-order differences:

$$(au')|_{x_{j+1/2}} \approx a_{j+1/2} \frac{u_{j+1} - u_j}{h} \quad (au')|_{x_{j-1/2}} \approx a_{j-1/2} \frac{u_j - u_{j-1}}{h}$$

where

$$a_{j+1/2} \equiv a(x_{j+1/2}).$$



CU-Boulder

315 of 396

## The basic stencil

We combine the differences for  $u'$  & for  $(au)'$  to obtain the operator

$$-(a(x_j)u'(x_j))'(x_j) \approx -\frac{a_{j+1/2} \left( \frac{u_{j+1} - u_j}{h} \right) - a_{j-1/2} \left( \frac{u_j - u_{j-1}}{h} \right)}{h}$$

& the problem becomes

$$\frac{1}{h^2} \left( -a_{j-1/2} u_{j-1} + \left( a_{j-1/2} + a_{j+1/2} \right) u_j - a_{j+1/2} u_{j+1} \right) = f_j \quad 1 \leq j \leq n \\ u_0 = u_n = 0.$$

CU-Boulder

316 of 396

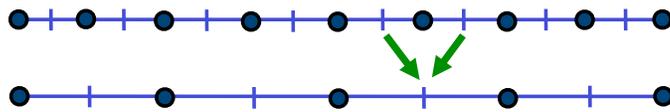
## Coarsening the variable coefficient problem

- A reasonable approach is to use a standard multigrid algorithm with linear interpolation, full weighting, & the stencil

$$A^{2h} = \frac{1}{(2h)^2} \begin{bmatrix} -a_{j-\frac{1}{2}}^{2h} & a_{j-\frac{1}{2}}^{2h} + a_{j+\frac{1}{2}}^{2h} & -a_{j+\frac{1}{2}}^{2h} \end{bmatrix}$$

where

$$a_{j+\frac{1}{2}}^{2h} = \frac{a_{2j+1/2}^h + a_{2j+3/2}^h}{2}$$



- The same stencil is obtained by the Galerkin relation.

## Variable mesh vs. variable coefficients

after scaling by  $\eta_j = (h_{j-1/2} + h_{j+1/2})/2$  &  $h$

- Variable mesh

$$-\frac{1}{h_{j-\frac{1}{2}}} u_{j-1}^h + \left( \frac{1}{h_{j-\frac{1}{2}}} + \frac{1}{h_{j+\frac{1}{2}}} \right) u_j^h - \frac{1}{h_{j+\frac{1}{2}}} u_{j+1}^h = \eta_j f_j^h$$

- Variable coefficients

$$\frac{1}{h} \left( -a_{j-\frac{1}{2}} u_{j-1} + \left( a_{j-\frac{1}{2}} + a_{j+\frac{1}{2}} \right) u_j - a_{j+\frac{1}{2}} u_{j+1} \right) = h f_j$$

- Correspondence

$$\frac{1}{h_{j-1/2}} \Leftrightarrow \frac{a_{j-1/2}}{h} \quad \frac{1}{h_{j+1/2}} \Leftrightarrow \frac{a_{j+1/2}}{h}$$

## A variable coefficient example

- We use V(2,1) cycle, full weighting, linear interpolation.
- We use  $a(x) = 1 + \rho \sin(k\pi x)$  &  $a(x) = 1 + \rho \text{rand}(x)$ .

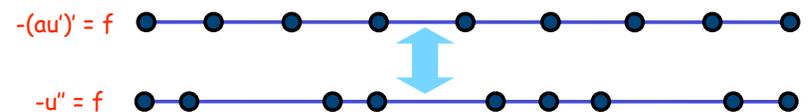
$n = 1024$

$a(x) = 1 + \rho \sin(k\pi x)$        $a(x) = 1 + \rho \text{rand}(x)$

$\rho$	$k=3$	$k=25$	$k=50$	$k=100$	
0	0.085	0.085	0.085	0.085	0.085
0.25	0.084	0.098	0.098	0.094	0.083
0.5	0.093	0.185	0.194	0.196	0.173
0.75	0.119	0.374	0.387	0.391	0.394
0.85	0.142	0.497	0.511	0.514	0.472
0.95	0.191	0.681	0.69	0.694	0.672

## Standard multigrid degrades if $a(x)$ is highly variable

MG for variable coefficients is equivalent to MG (with simple averaging) for Poisson's equation on a variable mesh.



But simple averaging won't accurately represent smooth components if  $x_{2j+1}^h$  is close to  $x_{2j}^h$  but far from  $x_{2j+2}^h$ .



## Pretend variability comes from mesh

$$A^h = \begin{bmatrix} -\frac{a_{j-\frac{1}{2}}}{h} & \frac{a_{j-\frac{1}{2}}}{h} + \frac{a_{j+\frac{1}{2}}}{h} & -\frac{a_{j+\frac{1}{2}}}{h} \\ \frac{1}{h_{j-\frac{1}{2}}} & \frac{1}{h_{j-\frac{1}{2}}} + \frac{1}{h_{j+\frac{1}{2}}} & -\frac{1}{h_{j+\frac{1}{2}}} \end{bmatrix}$$

- We can solve for the mesh sizes:  $h_{j-\frac{1}{2}} = \frac{h}{a_{j-\frac{1}{2}}}$ ,  $h_{j+\frac{1}{2}} = \frac{h}{a_{j+\frac{1}{2}}}$ .

- So linear interpolation yields

$$v_{2j+1}^h = \frac{h_{2j+3/2} v_j^{2h} + h_{2j+1/2} v_{j+1}^{2h}}{h_{2j+1/2} + h_{2j+3/2}} = \frac{a_{2j+1/2} v_j^{2h} + a_{2j+3/2} v_{j+1}^{2h}}{a_{2j+1/2} + a_{2j+3/2}}$$

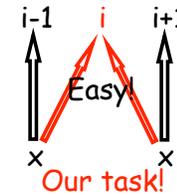
or, assuming that we've already interpolated to C-points  $2j$  &  $2j+2$ , it can be rewritten as

$$\left( \frac{a_{2j+1/2}}{h} + \frac{a_{2j+3/2}}{h} \right) v_{2j+1}^h - \frac{a_{2j+1/2}}{h} v_{2j}^{2h} - \frac{a_{2j+3/2}}{h} v_{2j+2}^{2h} = 0.$$

- We just used operator interpolation!
- Works for any stencil:  $[-\alpha \quad \alpha + \beta \quad -\beta]$ .

## Operator interpolation

- Assume that the error is known on the coarse grid:



So we can assume that  $e$  is known on the fine grid at F-points  $i_{\pm 1}$ . We then just need to relate  $e_i$  to  $e_{i_{\pm 1}}$ .

$$i-1 \rightleftarrows i \rightleftarrows i+1$$

- Assume smooth error with the ansatz that  $r = 0$ .
- Applying this at point  $i$ :  $-\alpha e_{i-1} + (\alpha + \beta) e_i - \beta e_{i+1} = 0$ .

- Solving for  $i$ :

$$e_i = \frac{\alpha}{\alpha + \beta} e_{i-1} + \frac{\beta}{\alpha + \beta} e_{i+1}.$$

- Accidental: F-points connect only to C-points. What do we do otherwise???

## Outline

### Chapters 1-5:

- ✓ Model Problems
- ✓ Basic Iterative Methods
  - Convergence tests
  - Analysis
- ✓ Elements of Multigrid
  - Relaxation
  - Coarsening
- ✓ Implementation
  - Complexity
  - Diagnostics
- ✓ Some Theory
  - Spectral vs. algebraic

### Chapters 6-10:

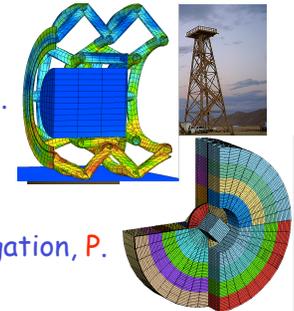
- ✓ Nonlinear Problems
  - Full approximation scheme
- ✓ Selected Applications
  - Neumann boundaries
  - Anisotropic problems
  - Variable meshes
  - Variable coefficients
- Algebraic Multigrid (AMG)
  - Matrix coarsening
- Multilevel Adaptive Methods
  - FAC
- Finite Elements
  - Variational methodology

Homework Due!

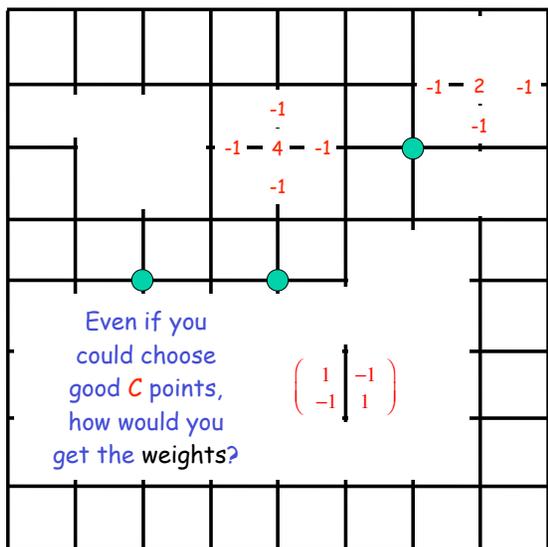
## 8. Algebraic multigrid (AMG)

unstructured grids, variable coefficients, ... assume SPD  $A$

- Automatically determines coarsening.
- AMG has two distinct phases:
  - setup phase: define MG components.
  - solution phase: perform MG cycles.
- AMG differs from geometric MG:
  - fix relaxation: point Gauss-Seidel.
  - choose coarsening: "grids" & prolongation,  $P$ .
- AMG principles:
  - algebraically smooth errors have small residuals:  $Ae \approx 0$ .
  - "strong" connections mean good neighbors: good C-points.
  - smooth error is locally almost constant:  $e \approx c$  for this  $A$ .
  - prolongation must match "smooth" error:  $e \in \text{range}(P)$ .
  - variational conditions apply: given  $P$ , set  $R=P^T$  &  $A_c=RAP$ .
  - only real task is to compute  $C$  &  $P$ : write  $e_f$  in terms of  $e_c$ .

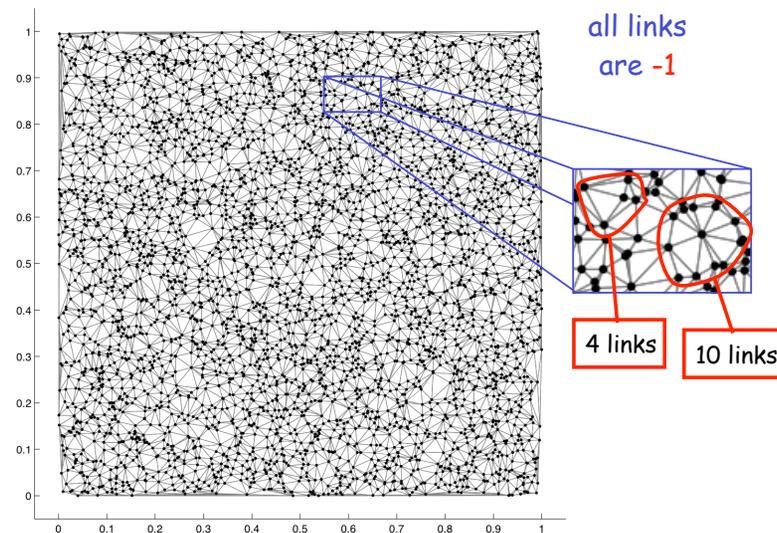


## Why AMG?



Even with this nice stencil, the geometry can give us trouble!

## Graph Laplacian



## AMG has two phases

### • Setup Phase

- Select coarse "grids,"  $\Omega^{m+1}$ ,  $m = 1, 2, \dots$
- Define interpolation,  $I_{m+1}^m$ ,  $m = 1, 2, \dots$
- Define restriction & coarse-grid operators,

$$I_{m+1}^m = (I_{m+1}^m)^T, \quad A^{m+1} = I_{m+1}^m A^m I_{m+1}^m.$$

### • Solve Phase

Standard MG processes: V-cycle, W-cycle, FMG, FAS, ...

- All AMG processes parallelize well, although coarse-grid selection must be done with care.

## AMG fundamental concept:

smooth error = "small" residuals/energy

- Error propagation via weighted Jacobi smoothing:

$$e^{k+1} = (I - \omega D^{-1}A) e^k.$$

- Error that is slow to converge satisfies

$$(I - \omega D^{-1}A) e \approx e \Rightarrow \omega D^{-1}A e \approx 0$$

$$\Rightarrow r \approx 0$$

- A little more precisely, assuming that  $\omega D \approx I$ , then slow-to-converge error has relatively small energy:

$$\langle Ae, e \rangle \ll \langle e, e \rangle \|A\|$$

## AMG uses strong connection to determine MG components

- Smoothing assumption:  
 $r \approx 0$  or  $\langle Ae, e \rangle \approx 0$ .
- We say that  $i$  is strongly connected to  $j$  if

$$-a_{ij} \geq \theta \max_{k \neq i} \{-a_{ik}\}, \quad 0 < \theta \leq 1.$$

- Zero-row-sum "M-matrices" actually satisfy

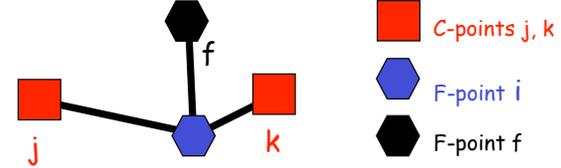
$$\langle Ae, e \rangle \approx \sum_{i \neq j} -\frac{a_{ij}}{2} (e_i - e_j)^2 \approx 0.$$

- So smooth error is more or less constant along strong connections.

We really mean matrices that have stencils like we've seen.

## Operator-induced interpolation

assume a graph of  $A$  & given coarse points



$$r_i = 0 \Rightarrow a_{ii} e_i + a_{ij} e_j + a_{ik} e_k + a_{if} e_f = 0.$$

To define  $e_i$  in terms of  $e_j$  &  $e_k$ , we must eliminate  $e_f$ .

Erase it?!

Use the smoothness principle that  $e$  is locally almost constant.

$$e_f = e_i \Rightarrow (a_{ii} + a_{if})e_i + a_{ij} e_j + a_{ik} e_k = 0$$

$\Rightarrow$  F-points determined by C-points!

We'll instead replace  $e_f$  by  $e_j$  &  $e_k$  in proportion to  $a_{fj}$  &  $a_{fk}$ .

## Some useful set definitions

- The set of strong connections of a variable  $u_i$ , that is, the variables upon whose values the value of  $u_i$  depends, is defined as

$$S_i = \left\{ j : -a_{ij} \geq \theta \max_{k \neq i} \{-a_{ik}\} \right\}.$$

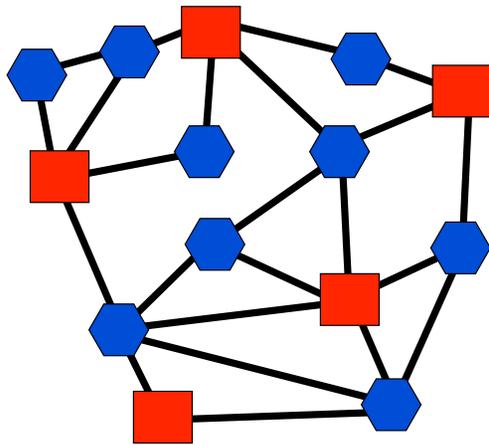
- The set of points strongly connected to variable  $u_i$  is denoted  $S_i^T = \{j : i \in S_j\}$ .
- The set of coarse-grid variables is denoted  $C$ .
- The set of fine-grid variables is denoted  $F$ .
- The set of interpolatory coarse-grid variables used to interpolate the value of the fine-grid variable  $u_i$  is denoted  $C_i$ .

## Choosing the coarse grid

- Two Criteria:
  - (C1) For each  $i \in F$ , every  $j \in S_i$  should either be in  $C$  or strongly connected to at least one point in  $C_i$ .
  - (C2)  $C$  should be a maximal subset with the property that no  $C$ -points are strongly connected to each other.
- Satisfying (C1) & (C2) is sometimes impossible.
- We use (C2) as a guide while enforcing (C1).

# Selecting the coarse-grid points

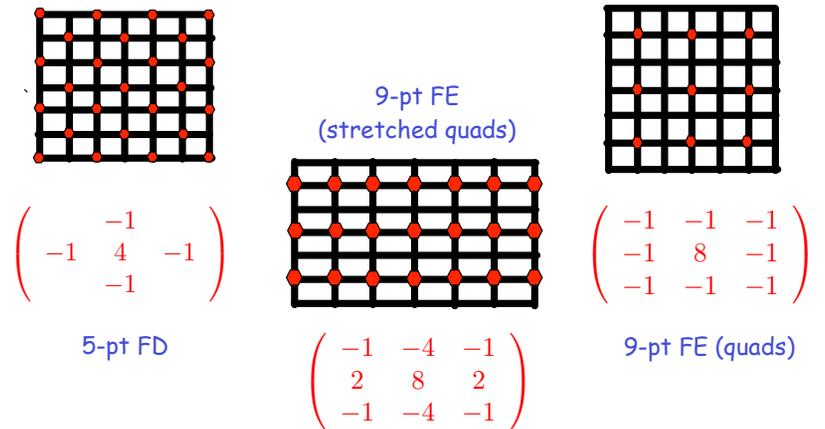
choose C-point to allow most F-points ("value")



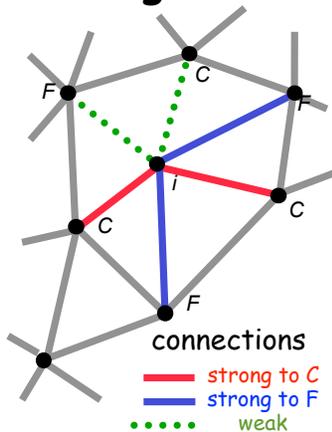
C-point selected (point with largest "value")  
 Neighbors of C-point become F-points  
 Next C-point selected (after updating "values")  
 F-points selected, etc.

# Sample grids for the Laplacian

5-pt FD, 9-pt FE (quads), & 9-pt FE (stretched quads)



# Prolongation is based on smooth error, strong connections (from M-matrices)



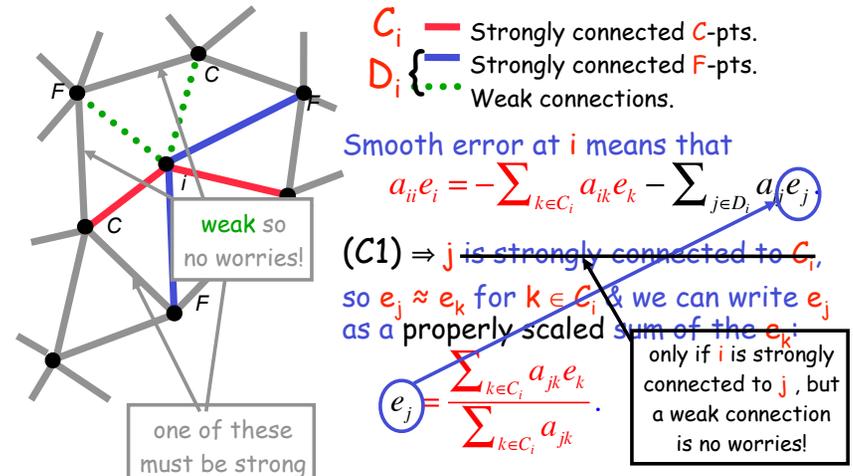
Smooth error is given by:  
 $r_i = a_{ii}e_i + \sum_{j \in C \cup F} a_{ij}e_j \approx 0.$

Prolongation :

$$(Pe)_i = \begin{cases} e_i, & i \in C \\ \sum_{k \in C} w_{ik}e_k, & i \in F \end{cases}$$

Actually, we want to allow for the possibility that we don't interpolate from all of C...

# Prolongation is based on smooth error, strong connections (from M-matrices)



Now we just substitute to get interpolation weights!

# Interpolation weights--the algebra

$$a_{ii}e_i = -\sum_{k \in C_i} a_{ik}e_k - \sum_{j \in D_i} a_{ij}e_j$$

$$+$$

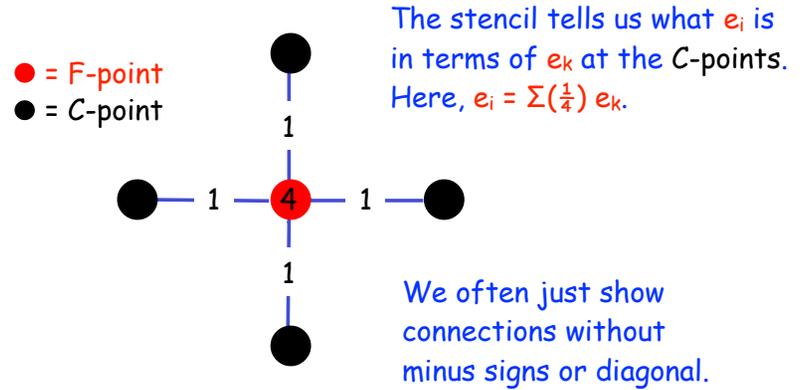
$$e_j = \frac{\sum_{k \in C_i} a_{jk}e_k}{\sum_{k \in C_i} a_{jk}} = \sum_{k \in C_i} \frac{a_{jk}}{\sum_{l \in C_i} a_{jl}} e_k$$

$$\Downarrow$$

$$w_{ik} = -\frac{1}{a_{ii}} \left( a_{ik} + \sum_{j \in D_i} a_{ij} \frac{a_{jk}}{\sum_{l \in C_i} a_{jl}} \right)$$

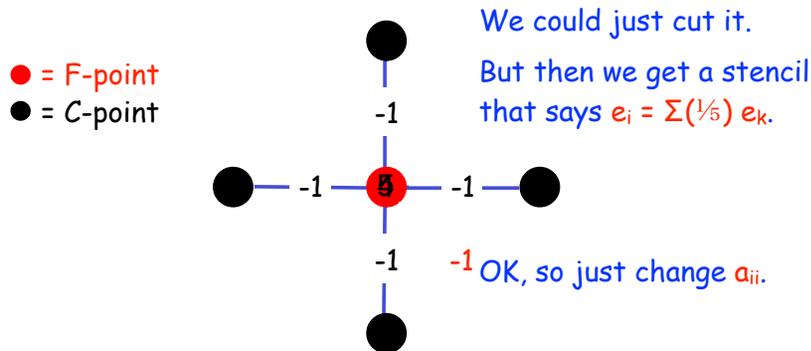
# Ideal setting

Suppose F-points are only connected to C-points.  
Interpolation only care about smooth  $e$ , so assume  $r_i = 0$ .



# Real setting

Suppose F-points are usually connected to other F-points.  
How do we eliminate these F-F connections?

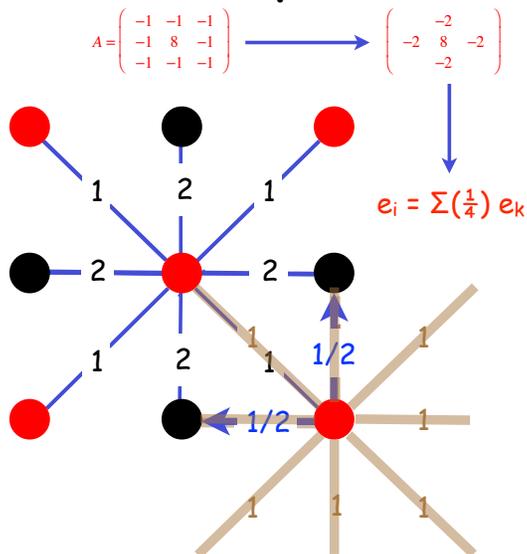


# Can we do better?

- We now have a direct way to determine interpolation.
- We just clip F-F connections & then adjust the diagonal (denominator) to make the weights sum to 1.
- Observe that what we are doing here is trying to write an F-point as a combination of neighboring C-points in a way that reflects the nature of smooth error.
- But if we assume that F-points are in the minority, can we use this our crude direct interpolation idea to eliminate an F-F connection by replacing the offending F-point with a linear combination of C-points in its neighborhood ( $C_i$ )?

## Example of AMG computation of P

● = F point  
● = C point



CU-Boulder

341 of 396

## AMG setup costs

- Many geometric MG methods need to compute prolongation & coarse-grid operators.
- The only additional expense in the AMG setup phase is the coarse-grid selection algorithm.
- So AMG's setup phase is usually only 10-25% more expensive than in geometric MG.
- But AMG is more robust in terms of geometric difficulties.

CU-Boulder

342 of 396

## AMG performance: Sometimes a success story

- AMG performs extremely well on the model problem (Poisson's equation, regular grid): optimal convergence factors (e.g., 0.14) & scalability w.r.t. problem size.
- AMG appears to be both scalable & efficient on diffusion problems on unstructured grids (e.g., 0.1-0.3).
- AMG handles anisotropic diffusion on structured & unstructured grids relatively well (e.g., 0.35).

CU-Boulder

343 of 396

## How does it perform (vol I)?

regular grids, plain, old, vanilla problems,  
unit square,  $n = 64$ , Dirichlet boundaries

- Laplacian:  $-u_{xx} - u_{yy} = 0$

Stencil	Convergence per cycle	Operator Complexity	Time per cycle
$\begin{pmatrix} -1 & -1 \\ -1 & 4 & -1 \\ -1 & -1 & -1 \end{pmatrix}$ 5-pt	0.054	2.21	0.29
$\begin{pmatrix} -1 & -1 & -1 \\ -1 & 8 & -1 \\ -1 & -1 & -1 \end{pmatrix}$ 5-pt skew	0.067	2.12	0.27
$\begin{pmatrix} -1 & -1 & -1 \\ -1 & 8 & -1 \\ -1 & -1 & -1 \end{pmatrix}$ 9-pt (-1,8)	0.078	1.30	0.26
$\begin{pmatrix} -1 & -4 & -1 \\ -4 & 20 & -4 \\ -1 & -4 & -1 \end{pmatrix}$ 9-pt (-1,-4,20)	0.109	1.30	0.26

- Anisotropic 5-Point Laplacian:  $-\epsilon u_{xx} - u_{yy} = 0$

$\epsilon$	0.001	0.01	0.1	0.5	1	2	10	100	1000
Convergence/cycle	0.084	0.093	0.058	0.069	0.054	0.079	0.087	0.093	0.083

CU-Boulder

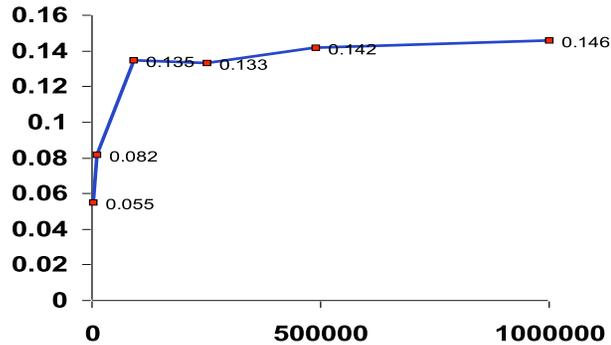
344 of 396

## How does it perform (vol II)?

structured meshes, rectangular domains

### 5-point Laplacian on regular rectangular grids

Convergence factor (y-axis) plotted against number of nodes (x-axis)



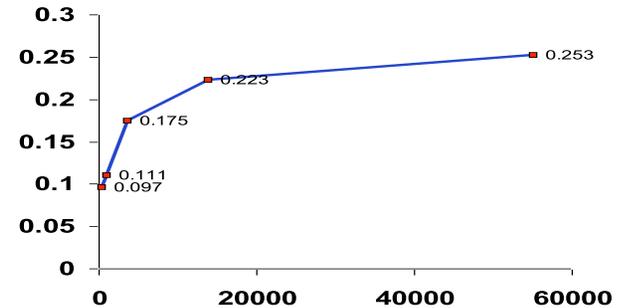
## How does it perform (vol III)?

unstructured meshes, rectangular domains

### Laplacian on random unstructured grids

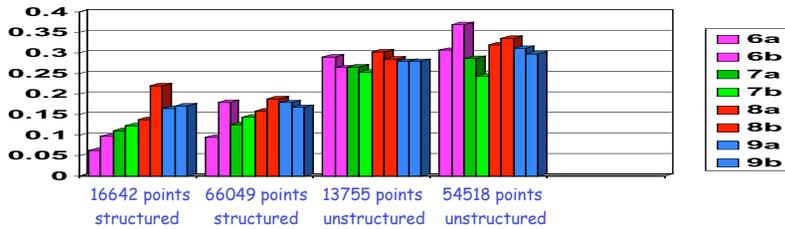
(regular triangulations, 15-20% nodes randomly collapsed into neighboring nodes)

Convergence factor (y-axis) plotted against number of nodes (x-axis)



## How does it perform (vol IV)?

$-\nabla \cdot (d(x,y) \nabla u) = 0$  on structured, unstructured grids



Problems used: "a" means parameter  $c = 10$ , "b" means  $c = 1,000$

$$6: d(x,y) = 1.0 + c|x-y| \quad 8: d(x,y) = \begin{cases} 1.0 & 0.125 \leq \max\{|x-0.5|, |y-0.5|\} \leq 0.25 \\ c & \text{otherwise} \end{cases}$$



$$7: d(x,y) = \begin{cases} 1.0 & x \leq 0.5 \\ c & x > 0.5 \end{cases}$$

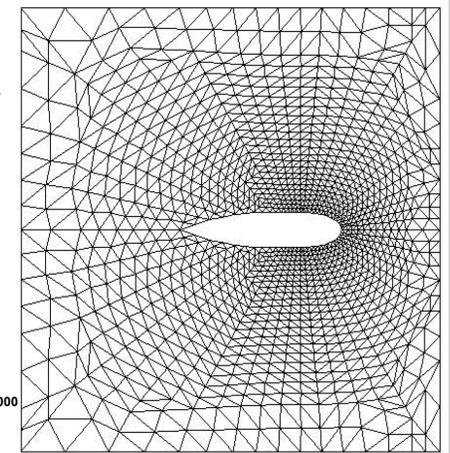
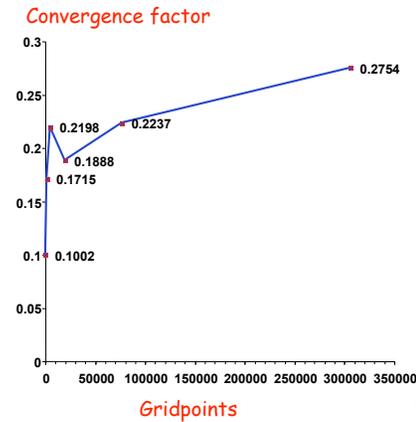


$$9: d(x,y) = \begin{cases} 1.0 & 0.125 \leq \sqrt{(x-0.5)^2 + (y-0.5)^2} \leq 0.25 \\ c & \text{otherwise} \end{cases}$$



## How does it perform (vol V)?

Laplacian operator, unstructured grids



Now for a glimpse at several other AMG topics...

## AMG for systems

- How can we do AMG on systems?

$$\begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \begin{pmatrix} u \\ v \end{pmatrix} = \begin{pmatrix} f \\ g \end{pmatrix}$$

- Naïve approach: "Block" AMG (block Gauss-Seidel, using scalar AMG to "solve" at each cycle)

$$u \leftarrow (A_{11})^{-1}(f - A_{12}v)$$

$$v \leftarrow (A_{22})^{-1}(g - A_{21}u)$$

Great Idea! Except that it often doesn't work!

Block AMG doesn't account for strong inter-variable coupling.

## AMG for systems: A solution

- To solve the system problem, allow interaction between the unknowns at all levels:

$$A^k = \begin{pmatrix} A_{11}^k & A_{12}^k \\ A_{21}^k & A_{22}^k \end{pmatrix} \quad \& \quad I_{k+1}^k = \begin{pmatrix} (I_{k+1}^k)_u & 0 \\ 0 & (I_{k+1}^k)_v \end{pmatrix}$$

- This is called the "unknown-based" approach.
- 2-D biharmonic  $(-\Delta)^2 u = f$ , Dirichlet & Neumann boundaries, unit square, uniform quadrilateral mesh:

Mesh spacing	0.125	0.0625	0.03135	0.015625
Convergence factor	0.22	0.35	0.42	0.44

## Adaptive AMG ( $\alpha$ AMG) to broaden applicability

adaptive interpolation

based on discovering the sense of smoothness

+

adaptive C-point choice

auto-determination of good coarse points

## Standard interpolation

Standard AMG collapses stencils by assuming smooth error is locally constant (Poisson "sense of smoothness"):

$$e_j = \frac{\sum_{k \in C_i} a_{jk} e_k}{\sum_{k \in C_i} a_{jk}}$$

$$a_{ii} e_i = - \sum_{k \in C_i} a_{ik} e_k - \sum_{j \in D_i} a_{ij} e_j$$

Strong C      F & Weak C

## Isn't standard interpolation OK?

- Suppose someone tried to make  $A$  "nice" for relaxation by scaling the diagonal so it's the identity:

$$\underline{A} \leftarrow D^{-1/2} A D^{-1/2}, \quad D = \text{diag}(a_{ii}).$$

- Relaxation still gives small residuals:

$$\underline{A} \underline{e} \approx 0.$$

- But

$$\underline{A} \underline{e} = (D^{-1/2} A D^{-1/2}) \underline{e} \approx 0 \Rightarrow A(D^{-1/2} \underline{e}) \approx 0 \Rightarrow \underline{e} \approx D^{1/2} c.$$

- So "smooth" here means  $\underline{e}_i \approx c \sqrt{a_{ii}}$ . This could vary a lot!

How can we discover what smooth vectors actually look like?

## Discovering smoothness

Relax on  $Ax = 0$  !!!

- What if we found a smooth error  $x$  that's far from  $c$ ?
- If, say,  $x_j = 1.4x_k$  for  $j \in D_i$  & all  $k \in C_i$ , then we could set

$$e_j \rightarrow \left( \sum_{k \in C_i} a_{jk} 1.4 e_k \right) / \left( \sum_{k \in C_i} a_{jk} \right).$$

- If  $x_j/x_k$  varies with  $k$ , then it's just a bit more complicated.

## $\begin{pmatrix} -1 & -4 & -1 \\ 2 & 8 & 2 \\ -1 & -4 & -1 \end{pmatrix}$ Adaptive C-point choice

- What does a "good"  $C$  mean? Systems???
- We want each  $i$  in  $F$  to "depend" on  $C$ : Irregular grids???  
Variable coeffs???
- $e_k$  given  $\forall k \in C \Rightarrow e_i$  well determined  $\forall i \in F$ .

- Let's look at the matrix: 
$$A = \begin{pmatrix} A_{ff} & A_{fc} \\ A_{cf} & A_{cc} \end{pmatrix}.$$

- What do we know about smooth  $e$ ?  $r_f = A_{ff} e_f + A_{fc} e_c = 0$ .
- Given  $e_c$ , is  $e_f$  well determined? Does  $r_f = 0$  determine  $e_f$ ?
- Given  $e_c$ , when does  $r_f = 0$  not determine  $e_f$ ?
- What if  $A_{ff}$  is singular?
- Then  $e_f$  cannot fully depend on  $e_c$ .
- We want  $A_{ff}$  to be well conditioned!!! How can we tell?

## Compatible relaxation (CR)

- To ensure that  $C$  is a really good set of coarse grid points, we want  $A_{ff}$  to be well conditioned. Thus, we can assess whether we have good  $C$ -points by "CR":

relax on  $A_{ff} x_f = 0$ .

- Fast CR means that  $F$  depends on  $C$  in the sense that smooth error ( $r_f = 0$ ) is quickly recovered from  $C$ .
- Fast CR also has the benefit that the  $F$ -point residuals can be made really small after  $F$ -point relaxation.
- We can also show that fast CR means that a  $P$  exists that gives good MG convergence:  $-A_{ff}^{-1} A_{fc}$ .
- This  $P$  isn't local; hopefully, CR can give us a good local  $P$ .

Problem: If  $A_{ff}$  isn't well conditioned, what then???

# Outline

## Chapters 1-5:

- √ Model Problems
- √ Basic Iterative Methods
  - Convergence tests
  - Analysis
- √ Elements of Multigrid
  - Relaxation
  - Coarsening
- √ Implementation
  - Complexity
  - Diagnostics
- √ Some Theory
  - Spectral vs. algebraic

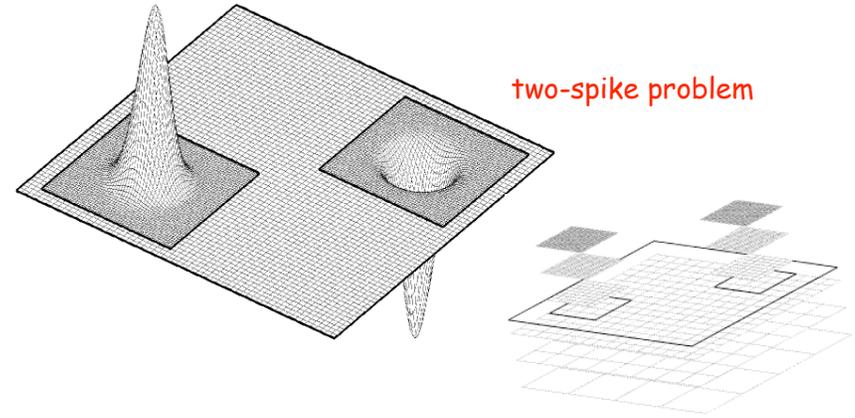
## Chapters 6-10:

- √ Nonlinear Problems
  - Full approximation scheme
- √ Selected Applications
  - Neumann boundaries
  - Anisotropic problems
  - Variable meshes
  - Variable coefficients
- √ Algebraic Multigrid (AMG)
  - Matrix coarsening
- Multilevel Adaptive Methods
  - FAC
- Finite Elements
  - Variational methodology

Homework Due!

# 9. Multilevel adaptive methods

fast adaptive composite grid method (FAC)

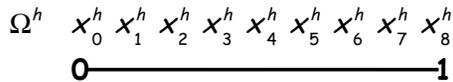


Local enhancement to resolve special regions of activity or interest.

## Model 1-D problem

$$-u''(x) = f(x), \quad 0 < x < 1$$

$$u(0) = u(1) = 0$$

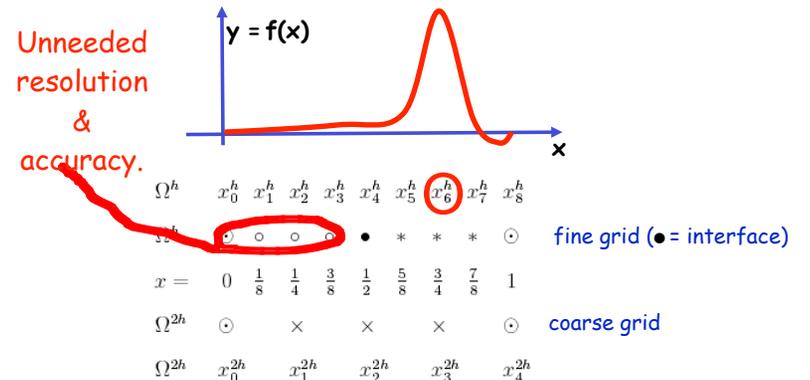


$$h^{-2}(-u_{i-1}^h + 2u_i^h - u_{i+1}^h) = f_i^h, \quad 1 \leq i \leq 7$$

$$u_0^h = u_8^h = 0, \quad h = \frac{1}{8}$$

## Local refinement

suppose  $f(x)$  has a spike at  $x = 3/4$   
but is smooth elsewhere

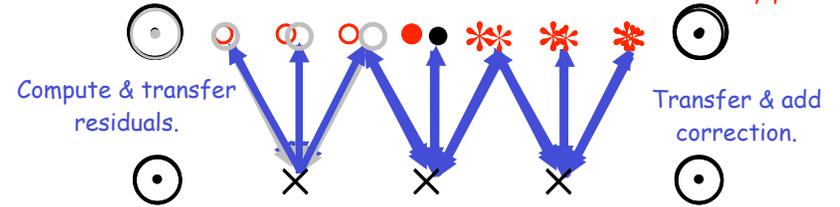


# Strategy

- Recognize that there's little value of having the fine grid in the smooth region,  $[0, 1/2]$ .
- Start with uniform grid & standard MG, then:
  - first eliminate relaxation in  $[0, 1/2]$ .
  - then eliminate intergrid transfers & residual calculations.
- Then interpret this process via the composite grid (=  $2h$ -points in  $[0, 1/2]$  +  $h$ -points in  $[1/2, 1]$ ).
- We'll try absurdly hard to eliminate work @  $x = 1/2$ , but we have in mind multi-dimensions & smaller patches.

# Local-Relaxation/Global-Correction MG

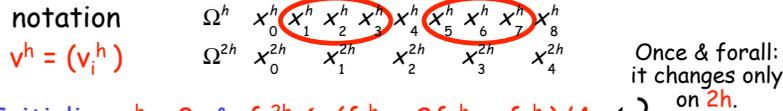
→ → WARNING: You won't get zero residuals on all of grid  $h$ . ← ←  
 Remember that the concept is right here, but we need to make it efficient--without changing the results!!!  
 Relax only in the local region.  
 No longer consider the usual global grids as usual.  $x = 1/2$  becomes a boundary point.



We start by using local relaxation & eliminating unnecessary residual transfers.  
 To save all of the work of computing & transferring residuals & corrections in regions where they don't change, we need some messy ALGEBRA!  
 Trust me on the more messy stuff...

## Eliminate relaxation

so  $r^h$  changes only on  $2h$   $v^h$  changes only here

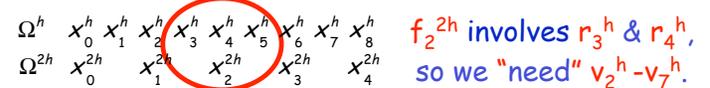


- Initialize  $v^h = 0$  &  $f_1^{2h} \leftarrow (f_1^h + 2f_2^h + f_3^h)/4$ .
- Relax on  $v^h$  on the local fine grid ( $x_5^h, x_6^h, x_7^h$ ).
- Compute  $r^h = f^h - A^h v^h$  (@  $x_3^h - x_7^h$ ) & transfer to  $2h$ :  
 $f_2^{2h} \leftarrow (r_3^h + 2r_4^h + r_5^h)/4$  &  $f_3^{2h} \leftarrow (r_5^h + 2r_6^h + r_7^h)/4$ .
- Compute an approximation,  $v^{2h}$ , to the solution of the  $2h$  residual equation,  $A^{2h} u^{2h} = f^{2h}$ .
- Update the residual at  $x_1^{2h}$  for later cycles:  
 $f_1^{2h} \leftarrow f_1^{2h} - (-v_0^{2h} + 2v_1^{2h} - v_2^{2h})/(2h)^2$ .
- Correct:  $v^h \leftarrow v^h + I_{2h}^h v^{2h}$ .

Why save  $v^h$  outside of local region?

cycle

## Eliminate more



Store  $v^h$  only @  $x_3^h - x_7^h$  & save  $v_2^h$  on  $2h$  (call it  $w_1^{2h}$ ).

- Initialize  $v^h = 0$ ,  $w_1^{2h} = 0$ , &  $f_1^{2h} \leftarrow (f_1^h + 2f_2^h + f_3^h)/4$ .
- Relax on  $v^h$  on the local fine grid ( $x_5^h, x_6^h, x_7^h$ ).
- Compute  $r^h = f^h - A^h v^h$  (@  $x_2^h - x_7^h$ ) & transfer to  $2h$ :  
 $f_2^{2h} \leftarrow (r_3^h + 2r_4^h + r_5^h)/4$  &  $f_3^{2h} \leftarrow (r_5^h + 2r_6^h + r_7^h)/4$ .
- Compute an approximation,  $v^{2h}$ , to the solution of the  $2h$  residual equation,  $A^{2h} u^{2h} = f^{2h}$ .
- Update the residual at  $x_1^{2h}$  for later cycles:  
 $f_1^{2h} \leftarrow f_1^{2h} - (-v_0^{2h} + 2v_1^{2h} - v_2^{2h})/(2h)^2$ .
- Accumulate the  $2h$  approximation:  $w_1^{2h} \leftarrow w_1^{2h} + v_1^{2h}$ .
- Correct:  $v^h \leftarrow v^h + I_{2h}^h v^{2h}$  (@  $x_3^h - x_7^h$ ).

cycle

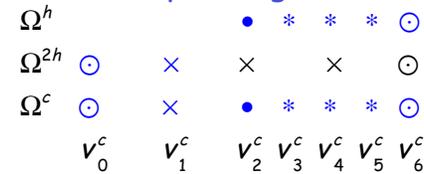
# Eliminate the rest

$\Omega^h$   $x_0^h x_1^h x_2^h x_3^h x_4^h x_5^h x_6^h x_7^h x_8^h$   $r_3^h$  doesn't change on  $h$ .  
 $\Omega^{2h}$   $x_0^{2h} x_1^{2h} x_2^{2h} x_3^{2h} x_4^{2h}$  Compute change in  $r_3^h$  on  $2h$ .

$2h$  residual @  $x_2^{2h} = g_2^{2h} - (-w_1^{2h} + 2v_4^h - v_6^h)/(2h)^2$ ,  
 where  $g_2^{2h} = (f_3^h + 2f_4^h + f_5^h)/4$ . (Messy algebra!)

- Initialize  $v^h = 0$ ,  $w_1^{2h} = 0$ , &  $f_1^{2h} \leftarrow (f_1^h + 2f_2^h + f_3^h)/4$ .
- Relax on  $v^h$  on the local fine grid ( $x_5^h, x_6^h, x_7^h$ ).
- Compute the right sides for  $2h$ :  
 $f_2^{2h} \leftarrow g_2^{2h} - (-w_1^{2h} + 2v_4^h - v_6^h)/(2h)^2$  &  
 $f_3^{2h} \leftarrow (r_5^h + 2r_6^h + r_7^h)/4$ .
- Compute an approximation,  $v^{2h}$ , to the solution of the  $2h$  residual equation,  $A^{2h} u^{2h} = f^{2h}$ .
- Update the residual at  $x_1^{2h}$  for later cycles:  
 $f_1^{2h} \leftarrow f_1^{2h} - (-v_0^{2h} + 2v_1^{2h} - v_2^{2h})/(2h)^2$ .
- Accumulate the  $2h$  approximation:  $w_1^{2h} \leftarrow w_1^{2h} + v_1^{2h}$ .
- Correct:  $v^h \leftarrow v^h + I_{2h}^h v^{2h}$  (@  $x_4^h - x_7^h$ ).

# Interpretation composite grid



What equation in  $v^h$  are we solving?

$(-v_0^c + 2v_1^c - v_2^c)/(2h)^2 = f_1^c \equiv (f_1^h + 2f_2^h + f_3^h)/4$   
 $(-v_{i-1}^c + 2v_i^c - v_{i+1}^c)/h = f_i^c \equiv f_{i+2}^h \quad i = 3, 4, 5$  } usual

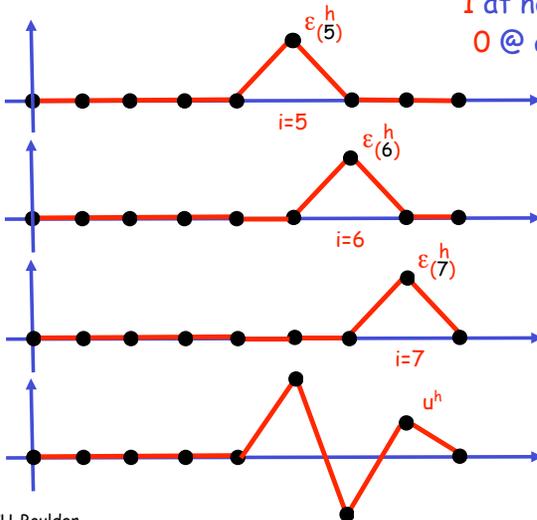
$(-v_1^c + 3v_2^c - 2v_3^c)/(2h)^2 = f_2^c \equiv (f_3^h + 2f_4^h)/4$  @interface

→ → The key is to treat the interface correctly. ← ←

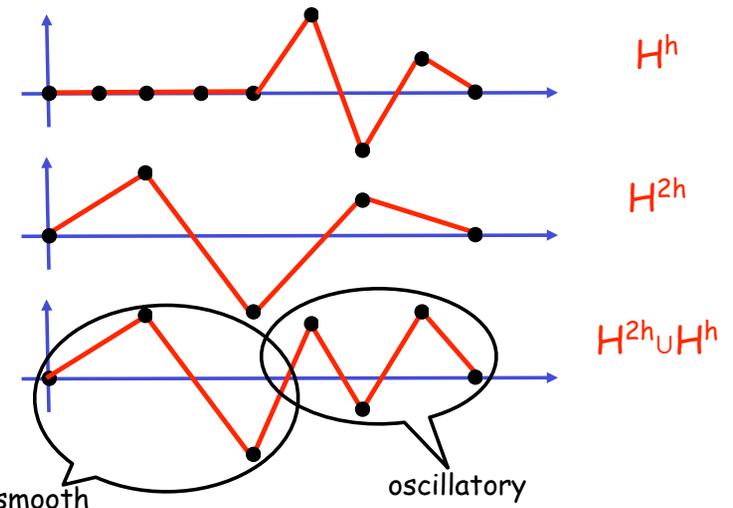
# Finite Element Local Refinement

continuous piecewise linear:

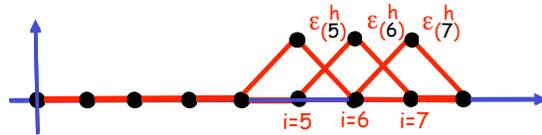
1 at node  $i=5,6,7$  only,  
0 @ all other nodes.



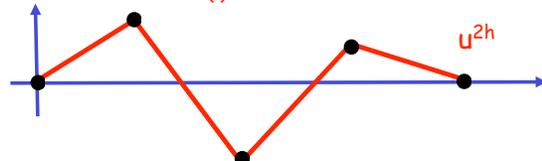
# Local Fine-Grid Refinement



# Abstract FE relaxation & $2h$ correction



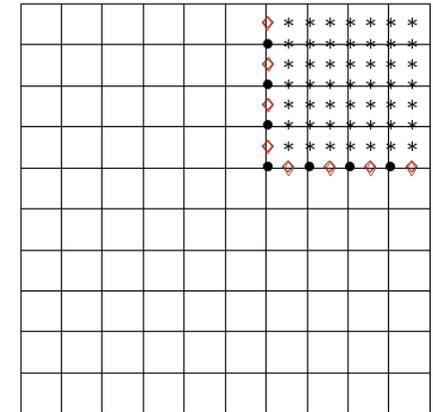
- Minimize  $F(u^h - s \epsilon_i^h)$  over  $s$  for  $i = 5, 6, 7$  in turn.



- Minimize  $F(u^h + u^{2h})$  over  $u^{2h}$  in  $H^{2h}$ .

# Issues

- Adaptivity
- Error estimates
- Norms (proper scaling)
- Multiple dimensions
  - Slave points  $\diamond$
  - More complicated stencils
- Data structures
- Parallel algorithms (AFAC)
- Time-space



# Two-spike example (Laplacian)

- Global grid  $h$  & one double-patch refined level  $h/2$ .
- $V(1,0)$ -cycles, Gauss-Seidel.
- Asymptotic convergence of the solver.
- Scaled  $L^2$  discretization error estimate.

Global $h$	Convergence factor	Discrete $L^2$ norm of discretization error
1/32	0.362	2.34e-2
1/64	0.367	5.742-3
1/128	0.365	1.43e-3

# Outline

## Chapters 1-5:

- Model Problems
- Basic Iterative Methods
  - Convergence tests
  - Analysis
- Elements of Multigrid
  - Relaxation
  - Coarsening
- Implementation
  - Complexity
  - Diagnostics
- Some Theory
  - Spectral vs. algebraic

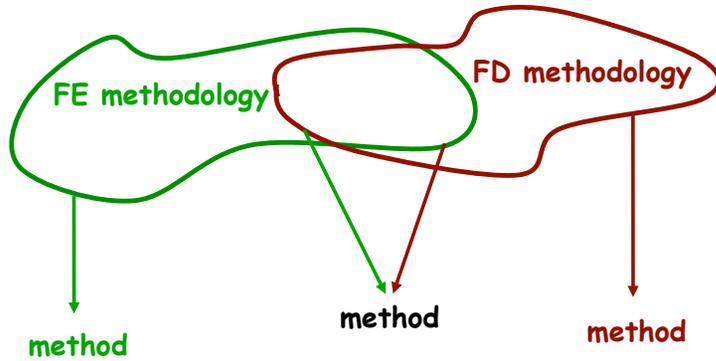
## Chapters 6-10:

- Nonlinear Problems
  - Full approximation scheme
- Selected Applications
  - Neumann boundaries
  - Anisotropic problems
  - Variable meshes
  - Variable coefficients
- Algebraic Multigrid (AMG)
  - Matrix coarsening
- Multilevel Adaptive Methods
  - FAC
- Finite Elements
  - Variational methodology

Homework Due!

# 10. Finite elements

FE, a variational discretization methodology



FD: differences at nodes & truncation error.

FE: weak form on discrete functions & approximation property.

Other methods: finite volume, collocation, spectral, ...

# Discretization

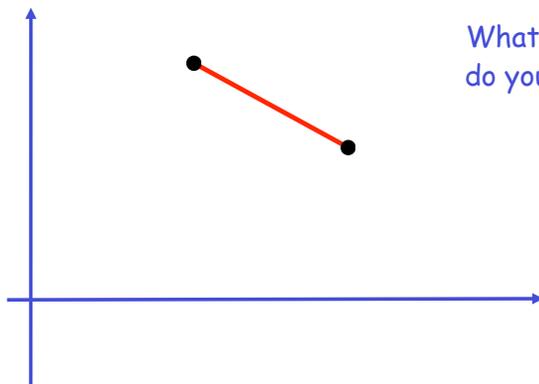
This is a subspace view.  
These discrete functions are in the space of solutions.

Discrete because they are determined by a finite number of nodes



→ → FIRST KEY POINT ← ←  
FE sees grid points (nodes or dofs) • only as characterizations of continuum functions.

# Localize



What choice do you have?

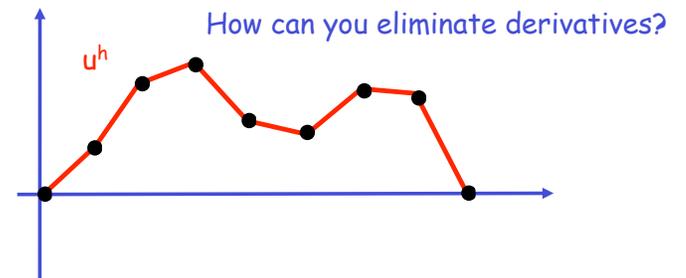
→ → SECOND KEY POINT ← ←

FE functions are localized to ensure a sparse matrix.

# Continuous piecewise linear functions

??? -  $u'' = f$  ???

How do you take 2<sup>nd</sup> derivatives of  $u^h$ ?



→ → THIRD KEY POINT ← ←

FE integrates away 2<sup>nd</sup> derivatives.

# Weak form!

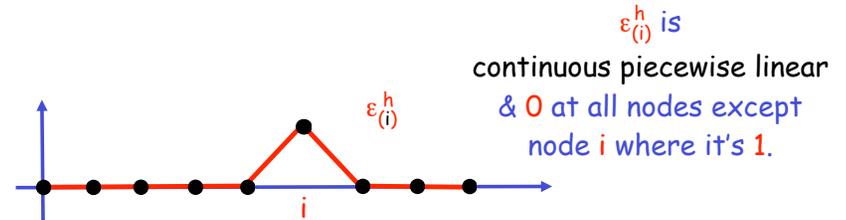
$-u'' = f$   $u =$  trial function  
 $\Rightarrow -u''v = fv$   $\forall v$   $v =$  test function  
 $\Rightarrow \int (-u'')v dx = \int f v dx$   $\forall v$   
 $\Rightarrow \int u'v' dx - u'v|_0^1 = \int f v dx$   $\forall v$  How is this weaker?  
Hint: 2 ways.  
 $\Rightarrow \int u'v' dx = \int f v dx$   $\forall v \ni v(0) = v(1) = 0$   
 $\Rightarrow (u', v') = (f, v)$   $\forall v \ni v(0) = v(1) = 0$

Why we have  $H^1$ !!! ? Need basis for space of admissible  $u^h$  &  $v$ .  
How would we discretize this?  $\Rightarrow$  can actually be reversed to  
 $\leftarrow$  if  $u$  is smooth enough.

CU-Boulder So the weak & strong forms are "equivalent". 377 of 396

# Representation

a basis of "hat" functions



Any continuous piecewise linear function can be represented by

$$u^h = \sum_i u_i^h \epsilon_{(i)}^h$$

Now we're back to using node values.

CU-Boulder

378 of 396

# Weak form & Galerkin discretization

$(u', v') = (f, v)$   $\forall v \ni v(0) = v(1) = 0$ .  
 $(\sum_j u_j^h \epsilon_{(j)}^h, v') = (f, v)$   $\forall v \ni v(0) = v(1) = 0$ .  
 $(\sum_j u_j^h \epsilon_{(j)}^h, \epsilon_{(i)}^h) = (f, \epsilon_{(i)}^h)$   $\forall i$ .  
 $\sum_j (\epsilon_{(i)}^h, \epsilon_{(j)}^h) u_j^h = (f, \epsilon_{(i)}^h)$   $\forall i$ .  $A^h u^h = f^h$

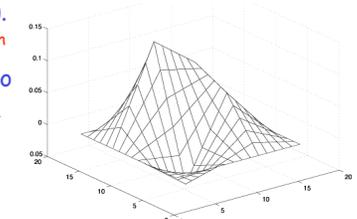
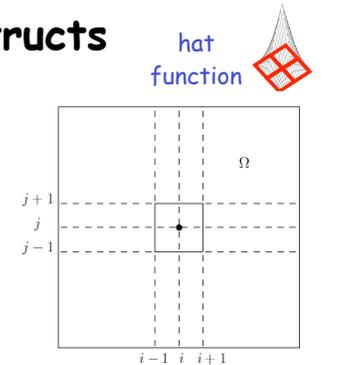
So the unknown is  $u^h = (u_j^h)$ , the matrix is  $A^h = (a_{ij}^h) = ((\epsilon_{(i)}^h, \epsilon_{(j)}^h))$ , & the right side is  $f^h = (f_i^h) = ((f, \epsilon_{(i)}^h))$ .

CU-Boulder

379 of 396

# 2D FE constructs

- Assume that  $\Omega$  is the unit square.
- Consider an  $n \times n$  grid of square "cells".
- Continuous piecewise bilinear elements:  $H^h \subset H_0^1(\Omega)$ .
- Each  $u^h$  in  $H^h$  is determined by its node values. This is how we'll represent them!
- Within each square:  $u^h = axy + bx + cy + d$  ( $u^h$  is linear in each coordinate direction).
- If  $u^h$  on one side of an element matches  $u^h$  on the other at the nodes, then we want to know that it matches on the common edge so that  $u^h$  is continuous: we want to know that specifying a piecewise bilinear function at the nodes gives us continuity.

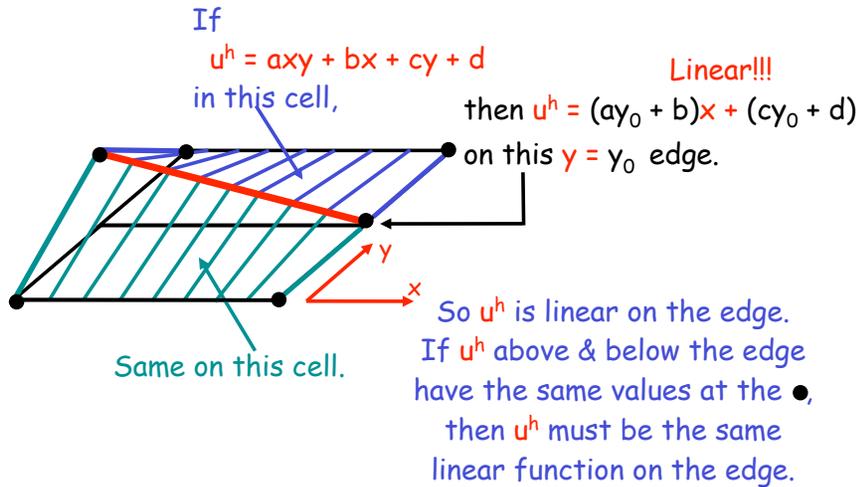


CU-Boulder

Why do we care???

380 of 396

## Continuity



## Weak form

- The Gauss Divergence Theorem & homogeneous boundary conditions yield

$$(Lu, v) = (-u_{xx} - u_{yy}, v) = (-\nabla \cdot \nabla u, v) = (\nabla u, \nabla v).$$

- Note:

$$(\nabla u, \nabla v) = \int_{\Omega} (u_x v_x + u_y v_y) d\Omega.$$

$$\nabla u = \begin{pmatrix} u_x \\ u_y \end{pmatrix}$$

- So the problem becomes

$$(\nabla u^h, \nabla v^h) = (f, v^h) \quad \forall v^h \in H^h$$

or

$$\int_{\Omega} (u_x^h v_x^h + u_y^h v_y^h) d\Omega = \int_{\Omega} f v^h d\Omega \quad \forall v^h \in H^h.$$

## Towards the matrix equation

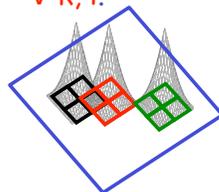
for nodal values

$$(\nabla u^h, \nabla v^h) = (f, v^h) \quad \forall v^h \in H^h$$

- Using  $u^h(x, y) = \sum_{ij} u_{ij}^h \varepsilon_{(ij)}^h(x, y)$  & choosing  $v^h = \varepsilon_{(kl)}^h$ :

$$\sum_{ij} u_{ij}^h (\nabla \varepsilon_{(ij)}^h, \nabla \varepsilon_{(kl)}^h) = (f, \varepsilon_{(kl)}^h) \quad \forall k, l.$$

- Matrix terms  $(\nabla \varepsilon_{(ij)}^h, \nabla \varepsilon_{(kl)}^h)$  are 0 when  $|i - k|$  or  $|j - l| > 1$ .



- We compute

$$(\nabla \varepsilon_{(ij)}^h, \nabla \varepsilon_{(ij)}^h) = 8/3 \quad \& \quad (\nabla \varepsilon_{(ij)}^h, \nabla \varepsilon_{(i \pm 1 j \pm 1)}^h) = -1/3.$$

- Assume  $f$  is fairly smooth locally:

$$(f, \varepsilon_{(kl)}^h) = \int_{\Omega} f \varepsilon_{(kl)}^h d\Omega \approx h^2 f(x_{kl}).$$

$$u^h = (u_{ij}^h) \quad \& \quad f^h = (h^2 f(x_{kl}))$$

actually,  
 $B^h f^h$   
(mass matrix)

## The matrix equation

$$A^h u^h = f^h$$

where

$$u^h = (u_{ij}^h) \quad \& \quad f^h = (h^2 f(x_{kl}))$$

& the matrix is given by the stencil

$$A_{ij}^h = \frac{1}{3} \begin{pmatrix} -1 & -1 & -1 \\ -1 & 8 & -1 \\ -1 & -1 & -1 \end{pmatrix}. \quad \text{stiffness matrix}$$

## Some matrix properties

$$A_{ij}^h = \frac{1}{3} \begin{pmatrix} -1 & -1 & -1 \\ -1 & 8 & -1 \\ -1 & -1 & -1 \end{pmatrix}$$

- Symmetric!  $ij$  "reaches" to  $i\pm 1j\pm 1$  as  $i\pm 1j\pm 1$  to  $ij$ .
- Singular?  $A^h \mathbf{1} = \mathbf{0}$ ?! Depends on boundaries!

Dirichlet west boundary:

$$A_{ij}^h = \frac{1}{3} \begin{pmatrix} 0 & -1 & -1 \\ 1 & 8 & -1 \\ 3 & -1 & -1 \end{pmatrix}$$

- Positive definite!

Diagonally dominant (strictly so @ boundaries).

But now we need to understand the PDE better, starting with choosing our universe of functions...

## A word about Sobolev spaces

- We're mucking about with forms like  $\int_{\Omega} (u_x v_x + u_y v_y) d\Omega$ , so we need to know that derivatives of functions in our universe can be multiplied together & integrated. It's enough to have  $\int_{\Omega} (u_x^2 + u_y^2) d\Omega < \infty$ , so our universe is  $H_0^1(\Omega) = \{u : u, u_x, u_y \in L^2(\Omega), u|_{\partial\Omega} = 0\}$ , where  $L^2(\Omega) = \{u : \int_{\Omega} u^2 d\Omega < \infty\}$ .
- We want what we compute to get close to what we want, so we're concerned about convergence in our universe: we need to know that limits of things that satisfy  $\int_{\Omega} (u_x^2 + u_y^2) d\Omega < \infty$  also satisfy it, that is, stay in our universe (completeness). Even if you start with nice continuous differential functions, you are led to some strange ones in your space.
- Think of  $L^2(\Omega) = \{u : \int_{\Omega} u^2 d\Omega < \infty\}$ . If  $u = 0$  except at finitely many points (say,  $u(i/n, j/n) = 1$  for  $i, j = 1, 2, \dots, n$ ), then  $\int_{\Omega} u^2 d\Omega = 0$ , so  $u = "0"$ ! This is true for any finite  $n$ , so it's true for a countable infinity of nonzeros!



## FE twin towers

$(\nabla u, \nabla v) = (f, v)$  **weak form** vs.  $(\nabla u, \nabla u)/2 - (f, u)$  **functional**

- Model problem  $\Omega \subset \mathbb{R}^2$ :  $Lu = -u_{xx} - u_{yy} = f$  in  $\Omega$ ,  $u = 0$  on  $\partial\Omega$ . Possible because  $L$  is self-adjoint. Duality: Solving  $Lu = f$  is equivalent to minimizing the weak functional  $F(u) = "(Lu, u)"/2 - (f, u) = (\nabla u, \nabla u)/2 - (f, u)$
- Sobolev spaces:  $L^2(\Omega) = \{u : \int_{\Omega} u^2 d\Omega < \infty\}$ ,  $H_0^1(\Omega) = \{u : u, u_x, u_y \in L^2(\Omega), u|_{\partial\Omega} = 0\}$ . Short story: 1<sup>st</sup> derivative test  $\nabla F(u) = Lu - f = 0$ , 2<sup>nd</sup> derivative test  $F''(u) = L > 0$
- $L$  is self-adjoint positive definite (more later):  $"(Lu, v)" = (\nabla u, \nabla v) = \int_{\Omega} (u_x v_x + u_y v_y) d\Omega = "(Lv, u)"$ ,  $"(Lu, u)" > 0$  if  $u \neq 0$ :  $\int_{\Omega} (u_x^2 + u_y^2) d\Omega \Rightarrow u=c$ , but  $u|_{\partial\Omega} = 0$ .

We use  $"(Lu, u)"$  for simplicity, but we really mean  $(\nabla u, \nabla u)$ .

This is formal:  $Lu$  is not defined on all of  $H_0^1(\Omega)$ .

## Long story...

dropping " " from  $(L \cdot, \cdot)$  for simplicity

- Using symmetry & linearity of  $L$  & bilinearity of the inner product:  $F(u+v) = (L(u+v), u+v)/2 - (f, u+v) = (Lu, u)/2 + (Lu, v) + (Lv, v)/2 - (f, u) - (f, v) = F(u) + [(Lu, v) - (f, v)] + (Lv, v)/2$ .
- Suppose  $u$  minimizes  $F$  but  $[(Lu, v) - (f, v)] < 0$  for some  $v$ . (Just flip the sign of  $v$  for the case  $[(Lu, v) - (f, v)] > 0$ .) Now replace  $v$  by  $\epsilon v$ :  $F(u + \epsilon v) = F(u) + \epsilon [(Lu, v) - (f, v)] + \epsilon (Lv, v)/2$ .
- Small enough  $\epsilon > 0$  means  $[(Lu, v) - (f, v)] + \epsilon (Lv, v)/2 < 0$ , which leads us to conclude that  $F(u + \epsilon v) < F(u)$ , a contradiction.
- This contradiction shows that  $(Lu, v) - (f, v)$  must be 0 for all  $v$ . Since this argument can easily be reversed, we thus conclude that  $F(u+v) \geq F(u) \quad \forall v \in H_0^1(\Omega) \Leftrightarrow (Lu, v) = (f, v) \quad \forall v \in H_0^1(\Omega)$ .

# Minimizing F via Rayleigh-Ritz

$$F(u) = (Lu, u)/2 - (f, u)$$

- Discretize by minimizing

$$F(u^h) = (Lu^h, u^h)/2 - (f, u^h)$$

over  $u^h \in H^h$ .

- Same as Galerkin that solves

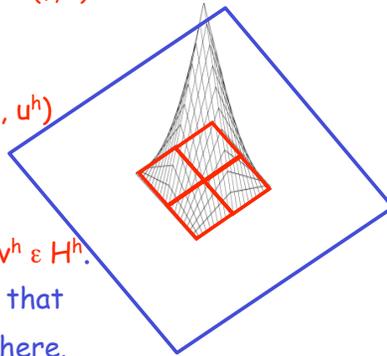
$$(Lu^h, v^h) = (f, v^h) \quad \forall v^h \in H^h.$$

- Basis:  $\epsilon_{(ij)}^h$  is the element of  $H^h$  that equals 1 @ node ij & 0 elsewhere.

- Expansion:  $u^h(x,y) = \sum_{ij} u_{ij}^h \epsilon_{(ij)}^h(x,y)$ .

- Old problem: What is  $Lu^h = -u_{xx}^h - u_{yy}^h$  ???

$$(Lu^h, v^h) = (\nabla u^h, \nabla v^h)$$



# Abstract FE relaxation

- Relaxation involves "local" changes:  $u^h \leftarrow u^h - s \epsilon_{(ij)}^h$  for some scalar  $s$  &  $\epsilon^h = \epsilon_{(ij)}^h$ .

But how do we pick  $s$  ???

# Abstract FE coarsening

again with focus on functions

- Coarsening involves a "global" change:  $u^h \leftarrow u^h + w^{2h}$  for some coarse-grid function  $w^{2h}$ .

But how do we pick  $w^{2h}$  ???

- Use FE principle of minimizing  $F(u^h + w^{2h})$  over  $w^{2h}$ :

$$F(u^h + w^{2h}) = (L(u^h + w^{2h}), u^h + w^{2h})/2 - (f, u^h + w^{2h})$$

$$= F(u^h) + (Lu^h - f, w^{2h}) + (Lw^{2h}, w^{2h})/2.$$

- Let  $w^{2h}$  be the root of the gradient of this quadratic functional w.r.t.  $w^{2h}$ . This is tricky because you need to write the gradient as a function in the subspace  $H^{2h}$ . We go instead from abstract functions to nodal vectors...

# Concrete FE coarsening: $I_{2h}^h$

now with focus on nodal vectors

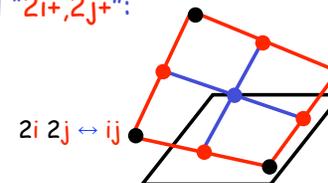
- Adding nodal representations of  $v^h$  &  $v^{2h}$ :

$$v^{2h}(x,y) = \sum_{ij} v_{ij}^{2h} \epsilon_{(ij)}^{2h}(x,y) \quad \text{(sum over } 2h \text{ indices } \leftrightarrow \text{ even } h \text{ indices)}$$

$$= \sum_{ij} v_{ij}^h \epsilon_{(ij)}^h(x,y) \quad \text{(sum over } h \text{ indices)}$$

- We should be able to do this because  $v^{2h} \in H^{2h} \subset H^h$ .

- Cell "2i+2j+":



$$v_{2i+2j}^h = v_{ij}^{2h}$$

$$v_{2i+1,2j}^h = (v_{ij}^{2h} + v_{i+1,j}^{2h})/2$$

$$v_{2i,2j+1}^h = (v_{ij}^{2h} + v_{i,j+1}^{2h})/2$$

$$v_{2i+1,2j+1}^h = (v_{ij}^{2h} + v_{i+1,j}^{2h} + v_{i,j+1}^{2h} + v_{i+1,j+1}^{2h})/4$$

Bilinear interpolation!

# Concrete FE coarsening (cont'd)

$$A^h u^h = f^h$$

$$A_{ij}^h = \frac{1}{3} \begin{pmatrix} -1 & -1 & -1 \\ -1 & 8 & -1 \\ -1 & -1 & -1 \end{pmatrix}$$

- Solving this matrix equation is equivalent to minimizing  $F^h(v^h) \equiv (A^h v^h, v^h)/2 - (f^h, v^h)$  (parens here mean Euclidean norm) over  $v^h \in H^h$ . So how do we now correct  $v^h$  ???

- We minimize  $F^h(v^h + I_{2h}^h v^{2h})$  over  $v^{2h} \in H^{2h}$ :

$$\begin{aligned} & F^h(v^h + I_{2h}^h v^{2h}) \\ &= (A^h (v^h + I_{2h}^h v^{2h}), v^h + I_{2h}^h v^{2h})/2 - (f^h, v^h + I_{2h}^h v^{2h}) \\ &= F^h(v^h) + \underbrace{(A^{2h} v^{2h}, v^{2h})/2 - (f^{2h}, v^{2h})}_{\text{variational conditions}} \quad f^{2h} = I_{2h}^T (f^h - A^h v^h) \\ &\equiv F^h(v^h) + F^{2h}(v^{2h}). \quad A^{2h} = I_{2h}^T A^h I_{2h} \end{aligned}$$

CU-Boulder

393 of 396

## Multigrid rules!

We conclude with a few observations:

- We have barely scratched the surface of the myriad ways that multigrid has been, & can be, employed.
- With diligence & care, multigrid can be made to handle many types of complications in a robust, efficient manner.
- Further extensions to multigrid methodology are being sought by many people working on many different problems.

CU-Boulder

395 of 396

# Outline

## Chapters 1-5:

- Model Problems
- Basic Iterative Methods
  - Convergence tests
  - Analysis
- Elements of Multigrid
  - Relaxation
  - Coarsening
- Implementation
  - Complexity
  - Diagnostics
- Some Theory
  - Spectral vs. algebraic

## Chapters 6-10:

- Nonlinear Problems
  - Full approximation scheme
- Selected Applications
  - Neumann boundaries
  - Anisotropic problems
  - Variable meshes
  - Variable coefficients
- Algebraic Multigrid (AMG)
  - Matrix coarsening
- Multilevel Adaptive Methods
  - FAC
- Finite Elements
  - Variational methodology

Homework Due!  
Presentations Next Week!

CU-Boulder

394 of 396

## Multigrid/multilevel/multiscale an important methodology

- Multigrid has proved successful on a wide variety of problems, especially elliptic PDEs, but has also found application in parabolic & hyperbolic PDEs, integral equations, evolution problems, geodesic problems, ...
- It can be optimal, often  $O(\# \text{ points})$ .
- It can be robust in a practical sense.
- It is of great interest because it is one of the very few scalable algorithms, & it can be parallelized readily & efficiently!
- But multigrid can also be a real pain!!!

CU-Boulder

396 of 396